

Macroscopic Evolution of Networks

CS 322: (Social and Information) Network Analysis
Jure Leskovec
Stanford University



Macroscopic Evolution

- How do networks evolve at the macroscopic level?
- Are there global phenomena of network evolution?
- What are analogs of the “small world” and “power-law degrees”

Three questions

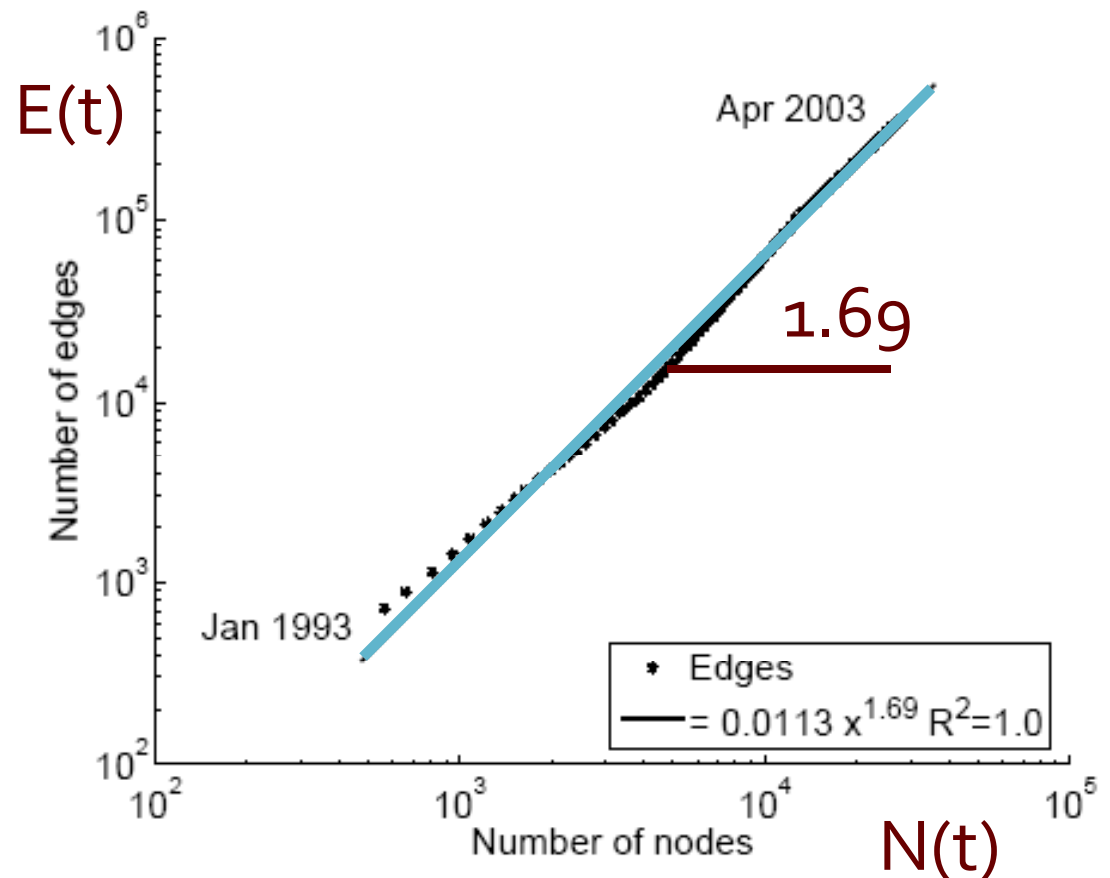
- What is the relation between the number of nodes $n(t)$ and number of edges $e(t)$ over time t ?
- How does diameter change as the network grows?
- How does degree distribution evolve as the network grows?

Network Evolution

- $N(t)$... nodes at time t
- $E(t)$... edges at time t
- Suppose that
$$N(t+1) = 2 * N(t)$$
- Q: what is
$$E(t+1) = 2 E(t)$$
- A: over-doubled!
 - But obeying the Densification Power Law

Densification – Physics Citations

- Citations among physics papers
- 1992:
 - 1,293 papers, 2,717 citations
- 2003:
 - 29,555 papers, 352,807 citations
- For each month M , create a graph of all citations up to month M



Densification Power Law

- Densification Power Law

- the number of edges grows faster than the number of nodes – average degree is increasing

$$E(t) \propto N(t)^a \quad \text{or equivalently} \quad \frac{\log(E(t))}{\log(N(t))} = \text{const}$$

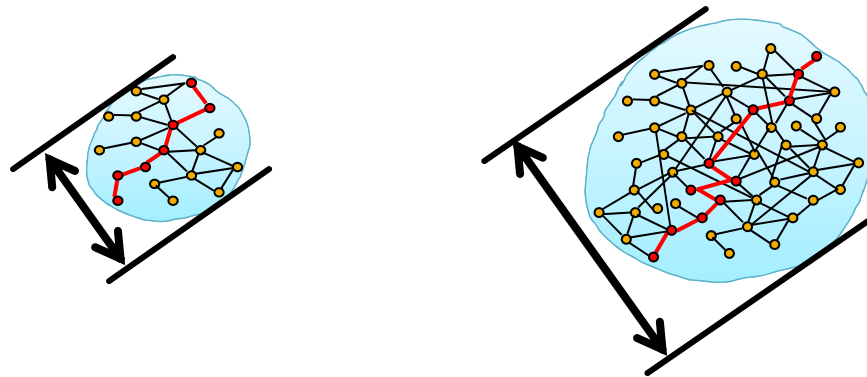
a ... densification exponent: $1 \leq a \leq 2$:

- **a=1**: linear growth – constant out-degree (traditionally assumed)
 - **a=2**: quadratic growth – clique

a = 1.67

Evolution of the Diameter

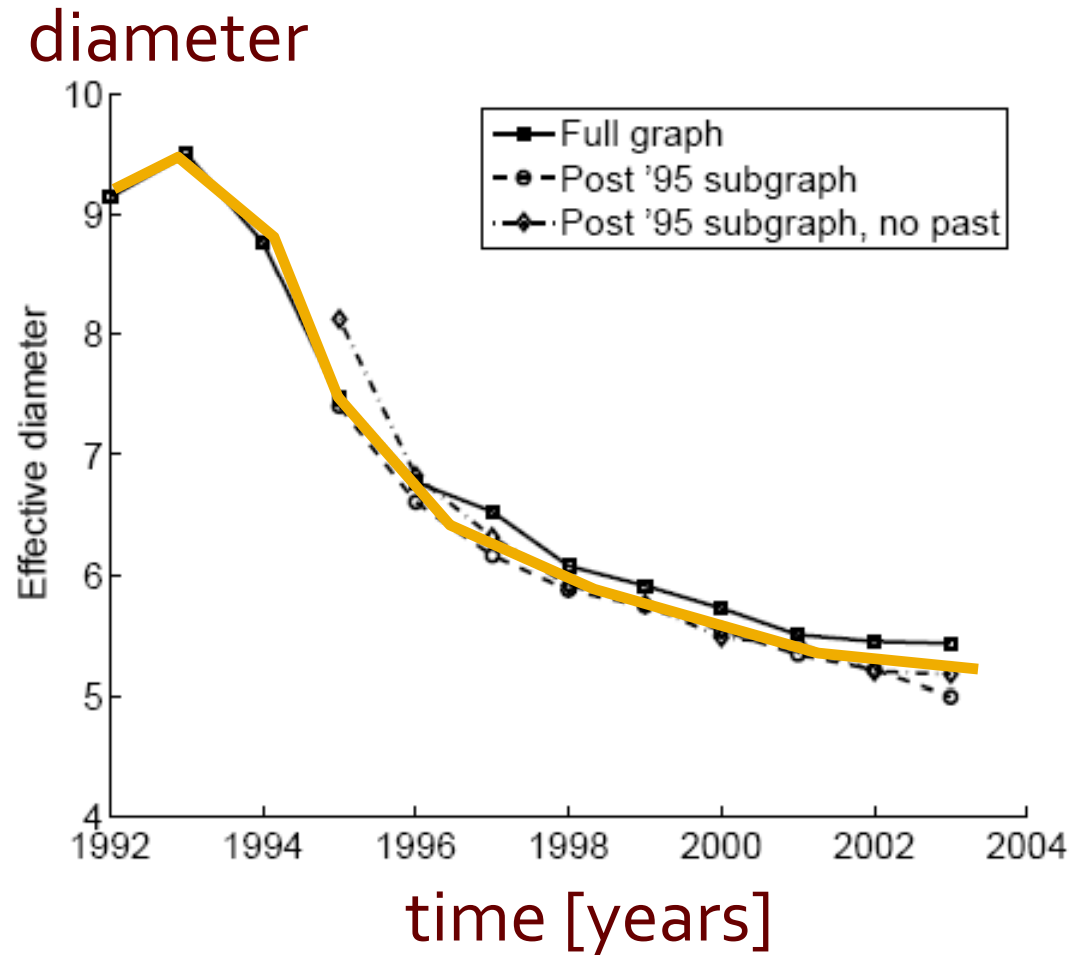
- Prior work on Power Law graphs hints at **Slowly growing diameter**:
 - diameter $\sim O(\log N)$
 - diameter $\sim O(\log \log N)$



- What is happening in real data?

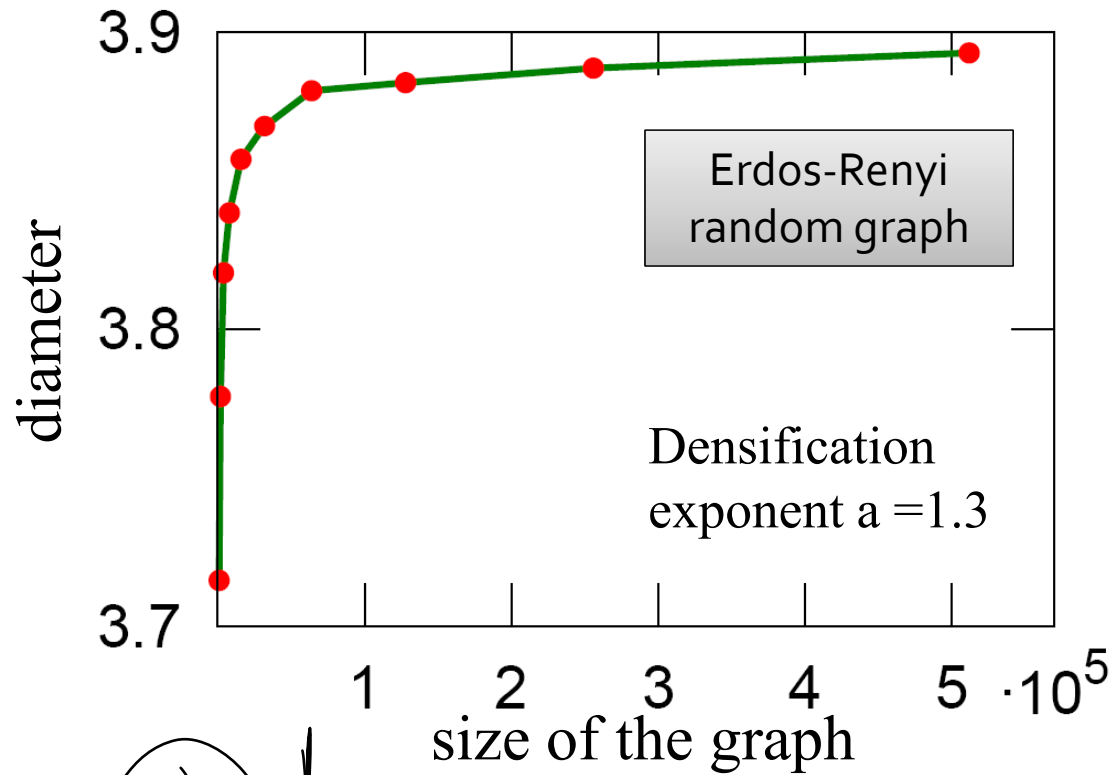
Diameter – ArXiv citation graph

- Citations among physics papers
- 1992 – 2003
- One graph per year



Diameter of a densifying G_{np}

Is shrinking diameter just a consequence of densification?



Densifying random graph has increasing diameter \Rightarrow There is more to shrinking diameter than just densification

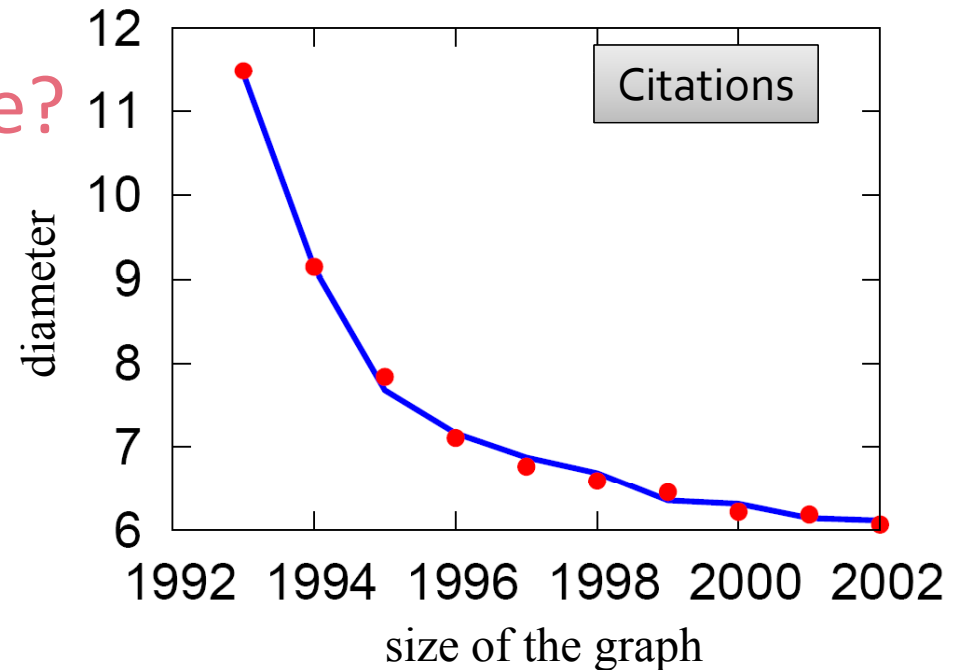
Handwritten notes: $\log E(t)$, $a = 1.3$, $\log N(t)$

Diameter of a rewired network

Is it the degree sequence?

Compare diameter of a:

- True network (red)
- Random network with the same degree distribution (blue)



Densification + degree sequence
give shrinking diameter

Connecting Degrees & Densification

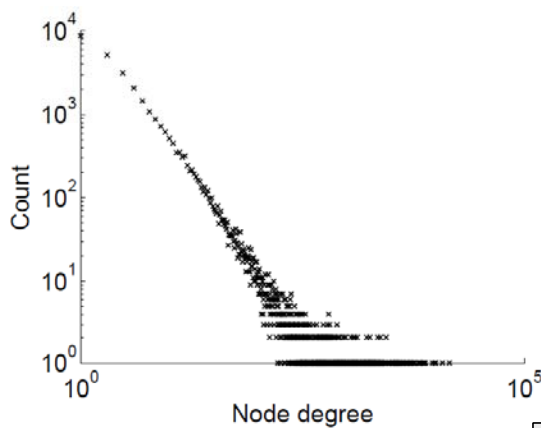
- How does degree distribution evolve to allow for densification?

- 2 Options:

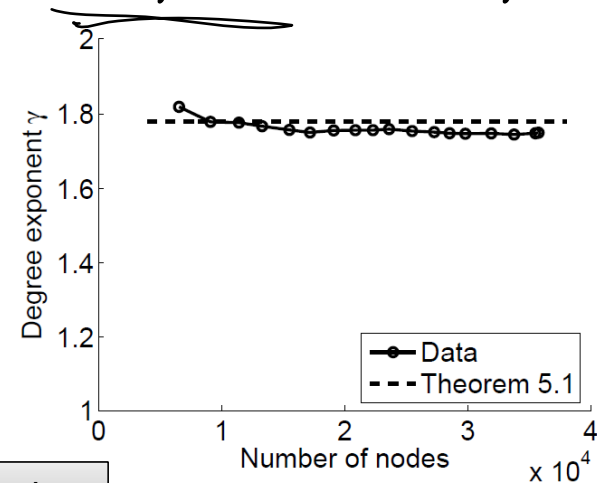
- 1) Degree exponent γ is constant:

- Fact 1: For degree exponent $1 < \gamma < 2$: $a = 2/\gamma$

$\gamma < 2$ REIN
 k -th moment
 won't exist



Email network



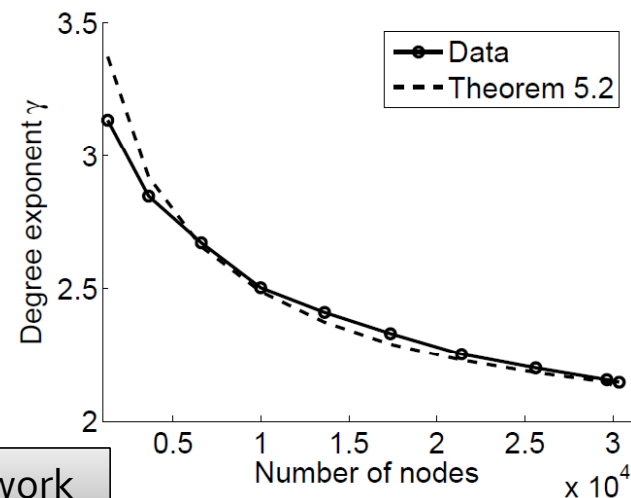
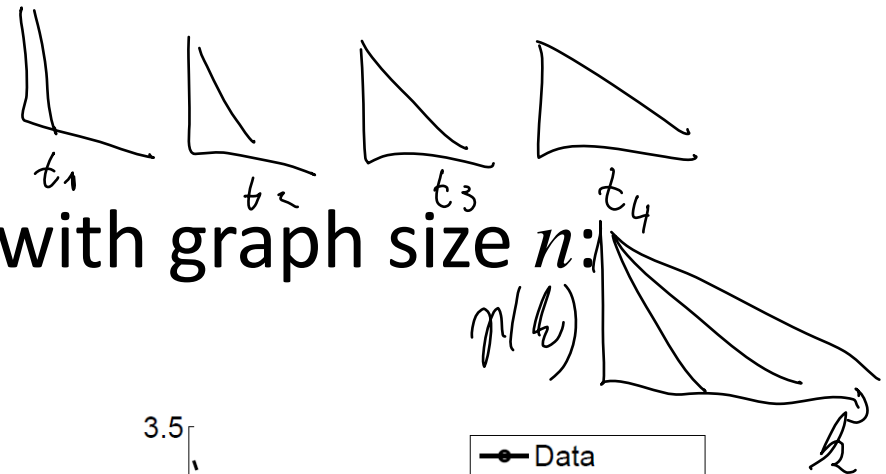
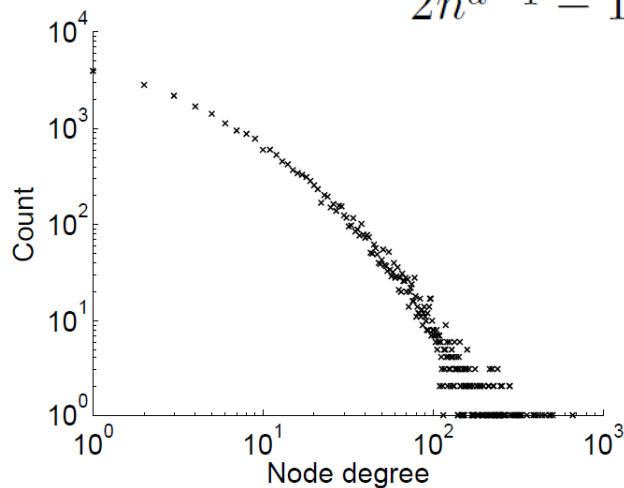
Connecting Degrees & Densification

- How does degree distribution evolve to allow for densification?

- 2 Options:

- 2) Exponent γ_n evolves with graph size n :

Fact 2:
$$\gamma_n = \frac{4n^{a-1} - 1}{2n^{a-1} - 1}$$



Patterns Hold in Many Graphs

- Densification and Shrinking diameter can be observed in many real life graphs:

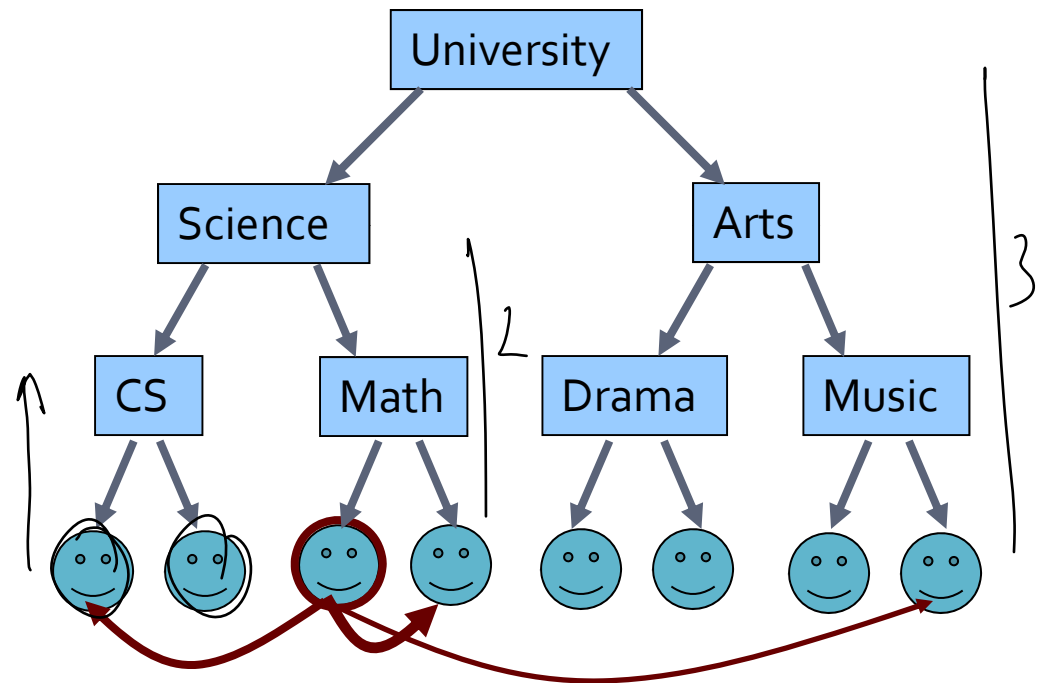
DATASET	NODES	EDGES	TIME	DPL EXPONENT
Arxiv HEP-PH	30,501	347,268	124 months	1.56
Arxiv HEP-TH	29,555	352,807	124 months	1.68
Patents	3,923,922	16,522,438	37 years	1.66
AS	6,474	26,467	785 days	1.18
Affiliation ASTRO-PH	57,381	133,179	10 years	1.15
Affiliation COND-MAT	62,085	108,182	10 years	1.10
Affiliation GR-QC	19,309	26,169	10 years	1.08
Affiliation HEP-PH	51,037	89,163	10 years	1.08
Affiliation HEP-TH	45,280	68,695	10 years	1.08
Email	35,756	123,254	18 months	1.12
IMDB	1,230,276	3,790,667	114 years	1.11
Recommendations	3,943,084	15,656,121	710 days	1.26

Densification – Possible Explanation

- What explains the Densification Power Law and Shrinking diameters
- Can we find a simple model of local behavior, which naturally leads to the observed phenomena?
- Yes! Community Guided Attachment

Community structure

- Let's assume the **community structure**
- One expects many within-group friendships and fewer cross-group ones
- How hard is it to **cross communities?**



Self-similar university community structure

Main Assumption

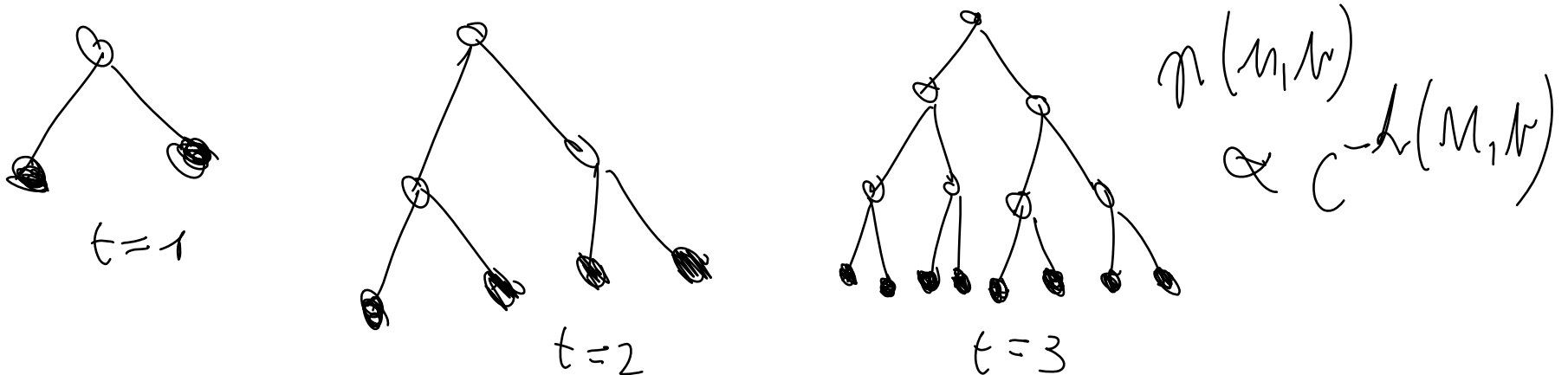
- Assume the cross-community linking probability of nodes at tree-distance h is scale-free
- Then the cross-community linking probability is:

$$f(h) = c^{-h}$$

where: $c \geq 1$... the *Difficulty constant*
 h ... tree-distance

Community guided attachment

- $n = 2^k$ nodes reside in the leaves of the b -way community hierarchy (assume $b=2$)



- Each node then independently creates edges based the community hierarchy, $f(h) = c^{-h}$
- How many edges m are in a graph of n nodes?

Densification Power Law

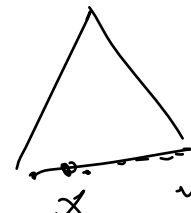
- Claim: Community Guided Attachment random graph model, the expected out-degree of a node is proportional to

1. $n^{1-\log_b(c)}$ if $1 \leq c < b$
2. $\log_b(n)$ if $c = b$
3. *constant* if $c > b$

Proof

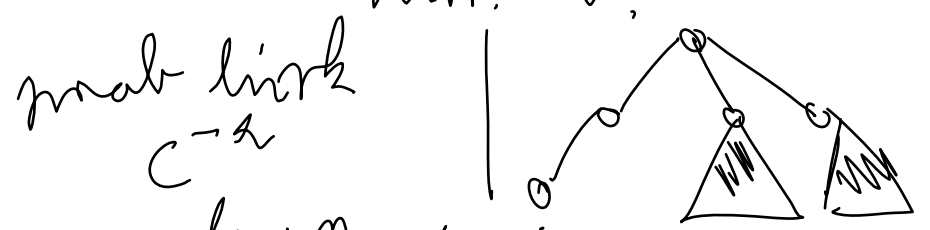
- What is expected out-degree of a node x ?

$$\bar{d}_x = \sum_{v \neq x} p(h(v, x)) = \sum_h \# \text{nodes at } h \cdot \text{prob of link to } h$$



How many nodes are at dist. h ?

$$= \sum_{h=1}^{\log_b n} (b-1) b^{h-1} \cdot c^{-h} = \frac{b-1}{c} \sum_{h=1}^{\log_b n} \left(\frac{b}{c}\right)^{h-1}$$



Proof

IF $c > b$ then $\sum_k^{\log n} \left(\frac{b}{c}\right)^k$ converges

$c < b$ then $-||- \propto$ largest element

$$\propto \left(\frac{b}{c}\right)^{\log n}$$

FACT:
 $X^{\log Y} = Y^{\log X}$

$$\bar{n} \propto n^{\log \frac{b}{c}} = n^{1 - \log b/c}$$

DPL EXP:

$$= n$$

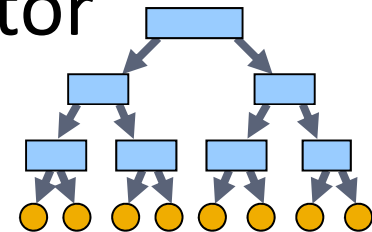
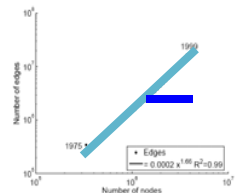
$$E = n \cdot n^{1 - \log b/c} = n^{2 - \log b/c}$$

Densification Power Law (2)

- Claim: The Community Guided Attachment leads to Densification Power Law with exponent

$$a = 2 - \log_b(c)$$

- a ... densification exponent $E(t) \propto N(t)^a$
- b ... community tree branching factor
- c ... difficulty constant, $1 \leq c \leq b$



Difficulty Constant

- DPL:

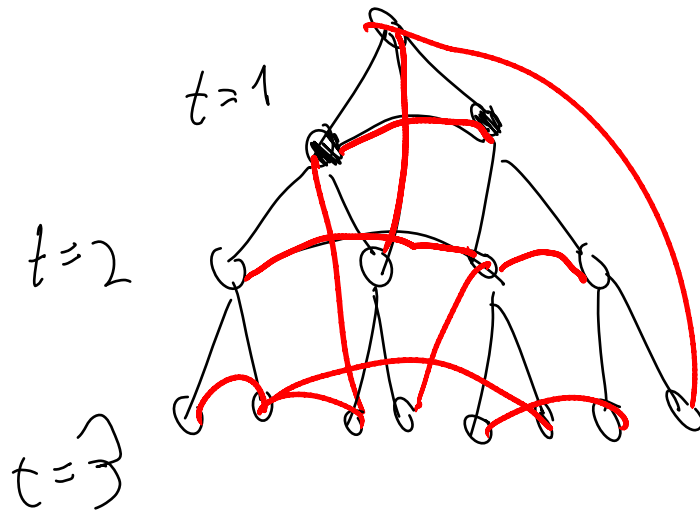
$$a = 2 - \log_b(c)$$

- Gives any non-integer Densification exponent
- If $c = 1$: easy to cross communities
 - Then: $a=2$, quadratic growth of edges – near clique
- If $c = b$: hard to cross communities
 - Then: $a=1$, linear growth of edges – constant out-degree

$$r(v, N_v) = r^{-h}$$

Extension of the model

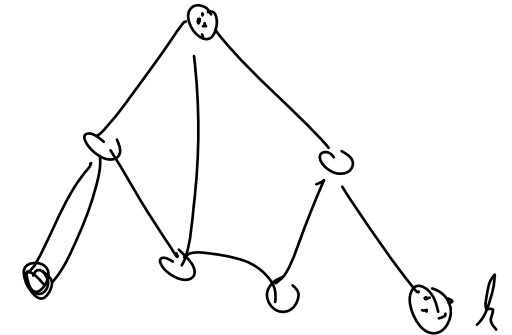
- Community tree evolves by a complete new level of nodes being added in each time step
 - Now nodes also reside in non-leaves of the tree



- Can show:
 - Densification + power-law degree distribution

B-way Branching Tree

- Topics of web pages:
 - Each node represents k pages
 - Nodes everywhere in the tree
 - Link up h levels with prob. $\sim \gamma^{-h}$



- Simplification: only bottom level creates links

Power-laws vs. Log-normal

- Normal $N(x) = \frac{1}{\sqrt{2\pi\sigma}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$
 - Limit of avg. of indep. draws from a random var from any distribution
- X is lognormally distributed if $Y = \ln X$ is normally distributed
- Log-normal density
 $f(x) = \frac{1}{\sqrt{2\pi\sigma}} e^{-\frac{(\ln x - \mu)^2}{2\sigma^2}} d(\ln x) = \frac{1}{\sqrt{2\pi\sigma}} \frac{1}{x} e^{-\frac{(\ln x - \mu)^2}{2\sigma^2}}$

Lognormal Distribution

$$f(x) = \frac{1}{\sqrt{2\pi\sigma}} \frac{1}{x} e^{-\frac{(\ln x - \mu)^2}{2\sigma^2}}$$

- Properties:
 - Finite mean/variance.
 - Skewed: mean > median > mode
 - Multiplicative:
 - X_1 lognormal, X_2 lognormal implies $X_1 X_2$ lognormal.

Power-law – Lognormal: similarities

- Looks similar to power-law on log-log plot:
- Power-law has linear log-density

$$\ln f(x) = -\alpha \ln x + \ln c$$

- For large σ , lognormal has nearly linear log-density:

$$\ln f(x) = -\ln x - \ln \sqrt{2\pi} \sigma - \frac{(\ln x - \mu)^2}{2\sigma^2}$$

- Similarly, both have near linear log-ccdfs.
- Question: how to differentiate them empirically?

Lognormal vs. Power Law

- Question: Is this distribution lognormal or a power law?
 - Reasonable follow-up: Does it matter?
- Generative model:
 - Size-independent + growth
 - Size (e.g., popularity, degree) start at one and changes in each step by a random multiplicative factor F

Generative Models: Lognormal

- Start with an organism of size X_0 .
- At each time step, size changes by a random multiplicative factor. $X_t = F_{t-1}X_{t-1}$
- Claim: If F_t is taken from a lognormal distribution, each X_t is lognormal.

Generative Models: Lognormal

- Claim: If F_t is taken from a lognormal distribution, each X_t is lognormal.
- If F_t are independent, identically distributed then (by CLT) X_t converges to lognormal distribution.
$$X_t = X_{t-1} F_{t-1} = \prod_{i=1}^t F_i = \sum_{i=1}^t \log F_i$$

$$X_0 = 1$$
- $\ln X_t = \text{size at time } t = \ln F_1 + \ln F_2 + \dots + \ln F_t$
- $\ln X_t = \sum \ln F_j \sim \text{normal dist} \rightarrow X_t \text{ is lognormal}$

BUT!

- If there exists a lower bound:

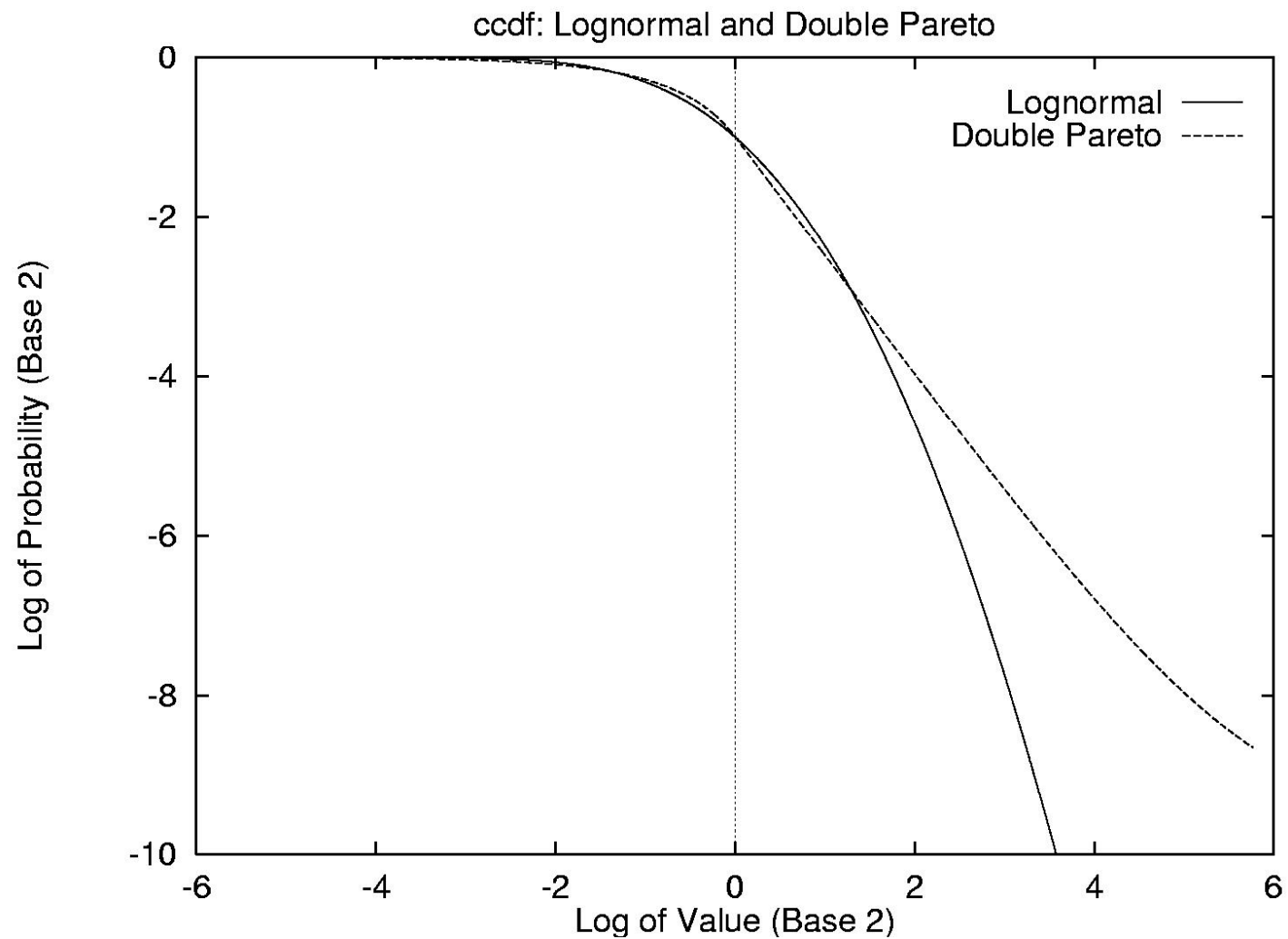
$$X_t = \max(\varepsilon, F_{t-1} X_{t-1})$$

- then X_t converges to a power law distribution
[Champernowne, 1953]
- Lognormal model easily pushed to a power law model

Double Pareto Distributions

- Consider continuous version of lognormal generative model:
 - At time t , $\log X_t$ is normal with mean μt and variance $\sigma^2 t$
- Suppose observation time is distributed exponentially
 - E.g., When Web size doubles every year.
- Resulting distribution is Double Pareto.
 - Between lognormal and Pareto.
 - Linear tail on a log-log chart, but a lognormal body.

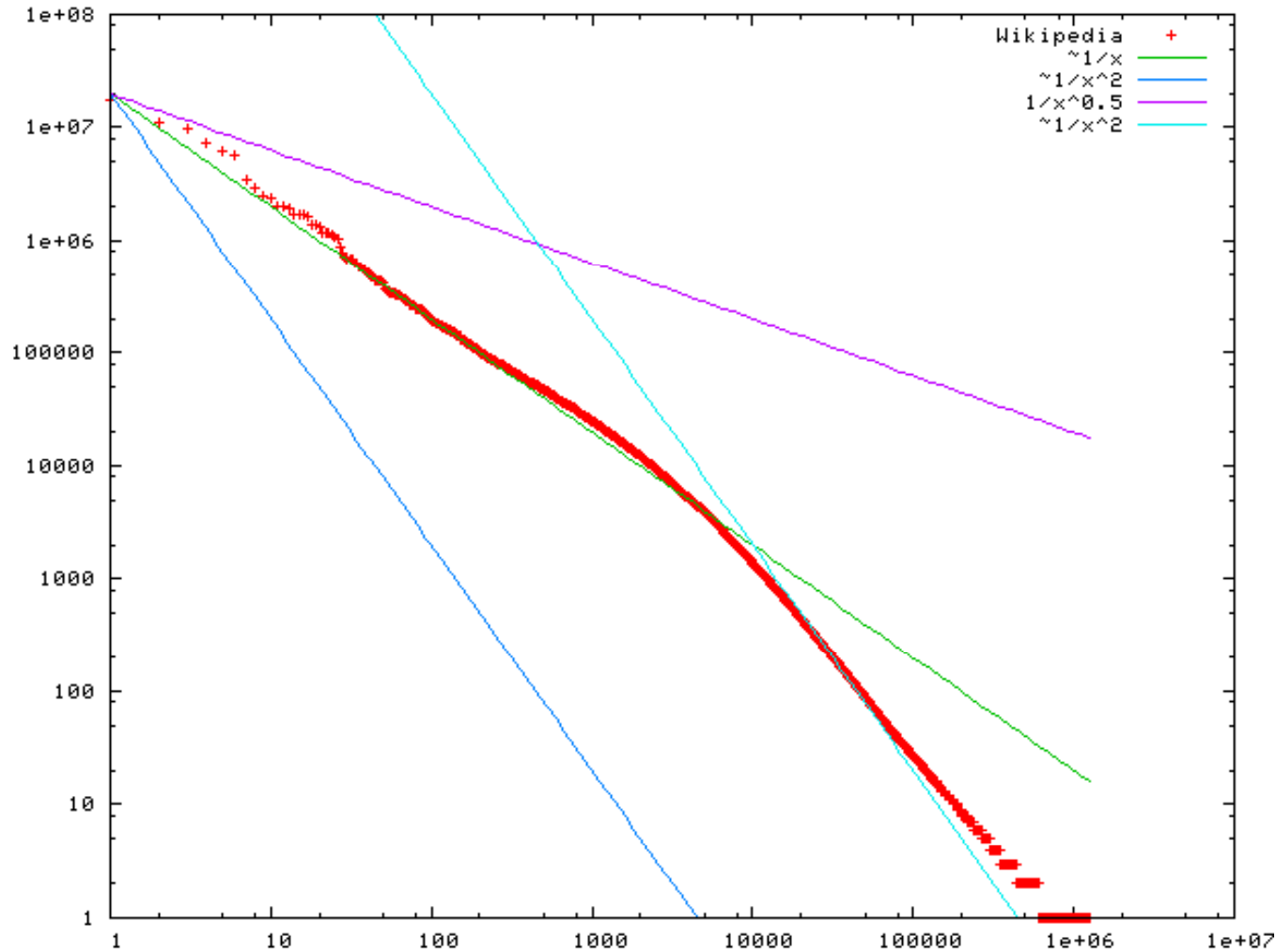
Lognormal vs. Double Pareto



More of power-laws

- Hierarchical model that produces power-laws
- Zipf's law – how does the frequency of a word depend on its rank?
- Empirically: $\text{rank} \sim 1/\text{freq}$

Rank vs. freq of words on Wikipedia



Optimization Model: Power Law

- Mandelbrot experiment: design a language over a d -ary alphabet to optimize information per character.

- Probability of j^{th} most frequently used word is p_j .
- Length of j^{th} most frequently used word is c_j .

- Average information per word:

$$H = -\sum_j p_j \log_2 p_j$$

- Average characters per word:

$$C = \sum_j p_j c_j$$

- Optimization leads to power law.

Explanation: Miller '57

- Alphabet of c letters plus spacebar
- Monkey sits at a typewriter and randomly pushes keys
 - Uniformly random sequence of $c+1$ symbols – spaces mark word boundaries
- Word ranks
 - $1, 2, \dots, c$ 1-letter words
 - $c+1, \dots, c^2+c+1$ 2-letter words
 - $c^2+c+1+1, \dots$ 3-letter words

Explanation

- Each j -letter word has frequency $(c+1)^{-(j+1)}$
- Pick representative point c^j for j -letter word
- Rank c^j is a j -letter word with freq $(c+1)^{-(j+1)}$

Explanation

- Write $(c+1)^{-(j+1)}$ as a function of rank