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# Optimizing Submodular Functions

CS246: Mining Massive Datasets

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# Announcements

- **HW 4 and Colab 8 due today (3/6 at 11:59 pm)**
- **HW 3, Colab 5, and Colab 6 grades released today (3/6)**
- **Colab 9 released today (3/6)**

# Recommendations: Diversity

- **Redundancy leads to a bad user experience**

**Obama Calls for Broad Action on Guns**

**Obama unveils 23 executive actions,  
calls for assault weapons ban**

**Obama seeks assault weapons ban,  
background checks on all gun sales**

- **Uncertainty around information need => don't put all eggs in one basket**
- **How do we optimize for diversity directly?**



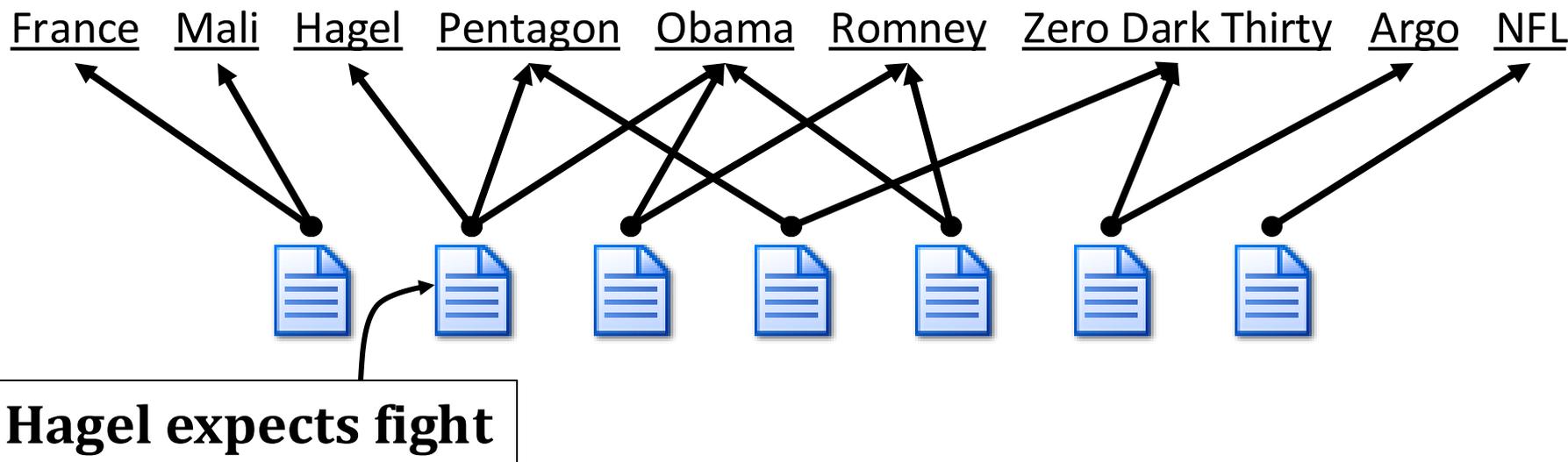




# Diversity as Coverage

# What is being covered?

- **Q: What is being covered?**
- **A: Concepts** (In our case: Named entities)



- **Q: Who is doing the covering?**
- **A: Documents**

# Simple Abstract Model

- **Suppose we are given a set of documents  $D$** 
  - Each document  $d$  covers a set  $X_d$  of words/topics/named entities  $W$

- **For a set of documents  $A \subseteq D$  we define**

$$F(A) = \left| \bigcup_{i \in A} X_i \right|$$

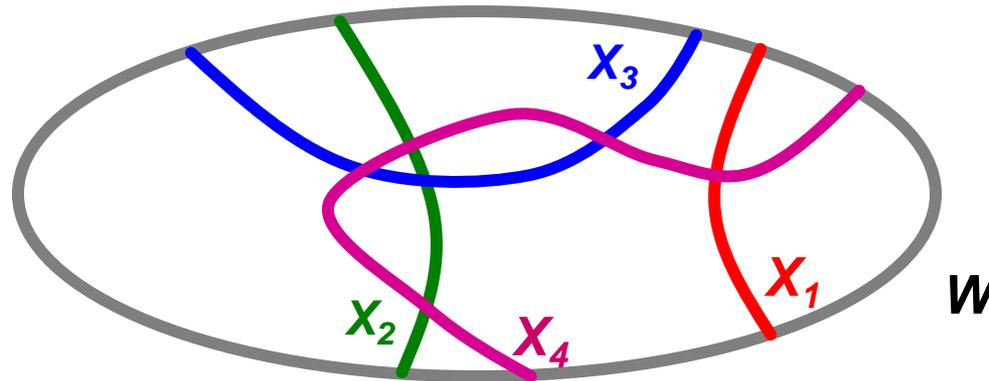
- **Goal: We want to**

$$\max_{|A| \leq k} F(A)$$

- **Note:  $F(A)$  is a set function:  $F(A): \text{Sets} \rightarrow \mathbb{N}$**

# Maximum Coverage Problem

- Given universe of elements  $W = \{w_1, \dots, w_n\}$  and sets  $X_1, \dots, X_m \subseteq W$



- **Goal: Find  $k$  sets  $X_i$  that cover the most of  $W$** 
  - More precisely: Find  $k$  sets  $X_i$  whose size of the union is the largest
  - **Bad news: A known NP-complete problem**

# Simple Greedy Heuristic

## Simple Heuristic: Greedy Algorithm:

- Start with  $A_0 = \{\}$
- For  $i = 1 \dots k$ 
  - Find set  $d$  that  $\max F(A_{i-1} \cup \{d\})$
  - Let  $A_i = A_{i-1} \cup \{d\}$

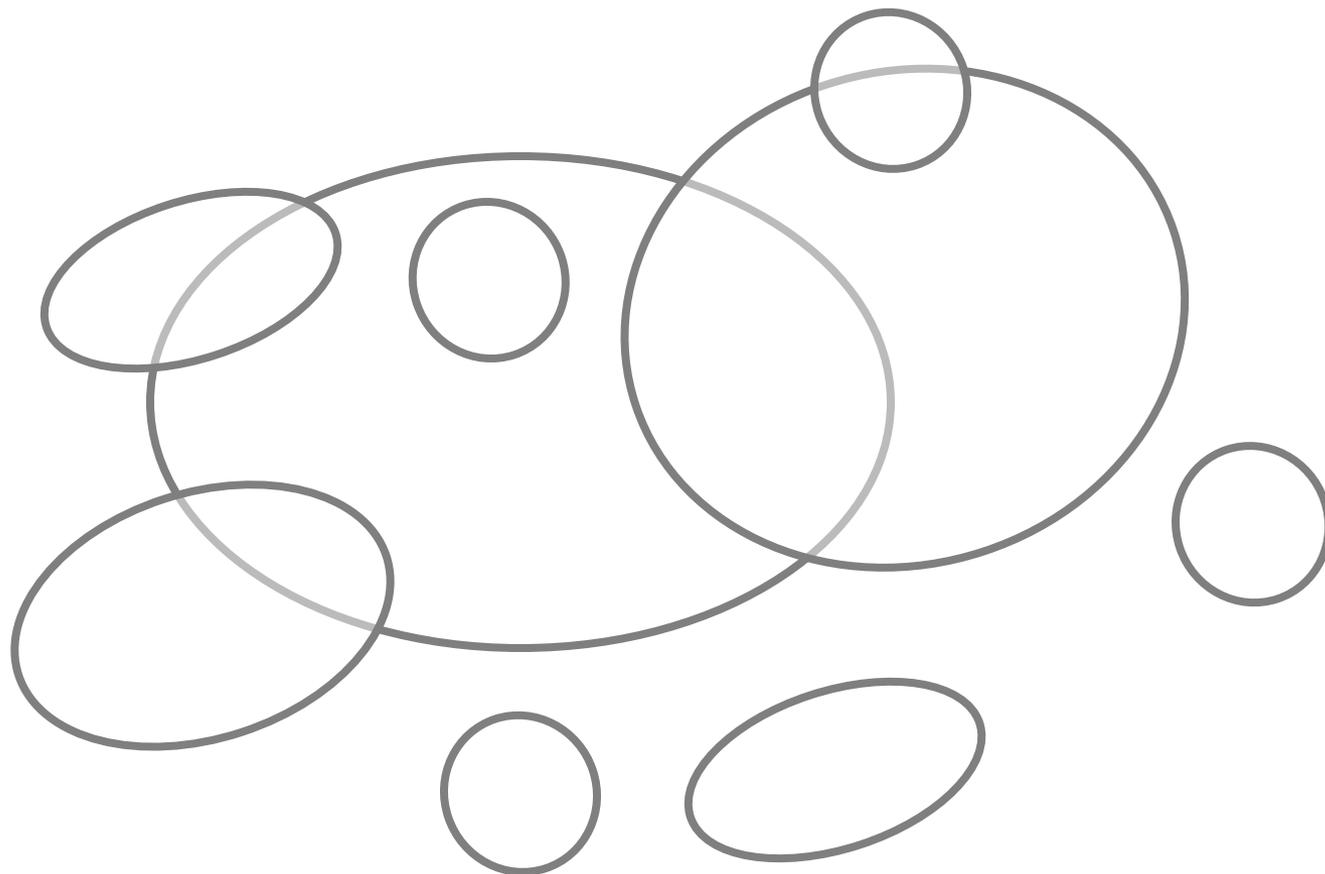
$$F(A) = \left| \bigcup_{d \in A} X_d \right|$$

## ■ Example:

- Eval.  $F(\{d_1\}), \dots, F(\{d_m\})$ , pick best (say  $d_1$ )
- Eval.  $F(\{d_1\} \cup \{d_2\}), \dots, F(\{d_1\} \cup \{d_m\})$ , pick best (say  $d_2$ )
- Eval.  $F(\{d_1, d_2\} \cup \{d_3\}), \dots, F(\{d_1, d_2\} \cup \{d_m\})$ , pick best
- And so on...

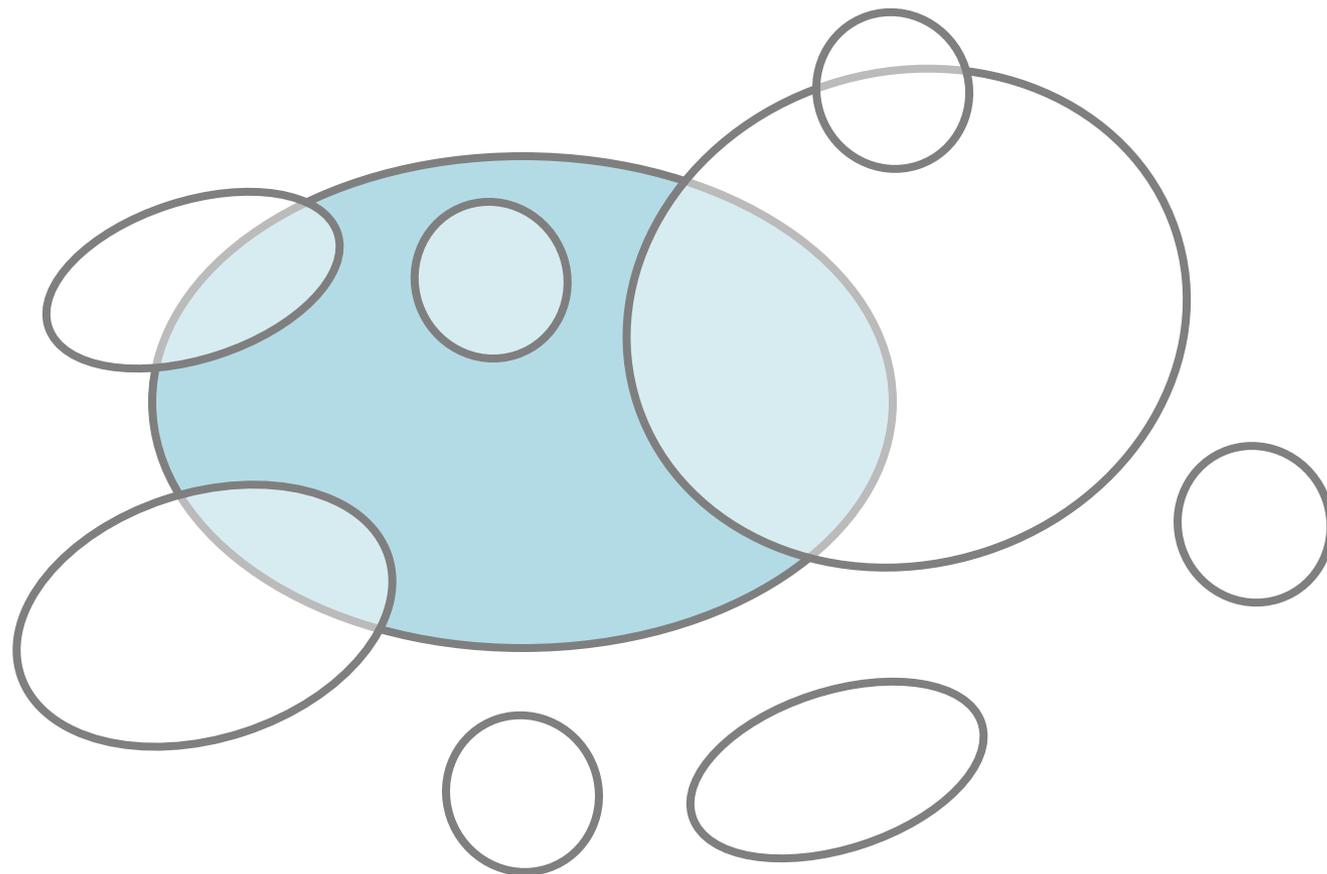
# Simple Greedy Heuristic

- **Goal: Maximize the covered area**



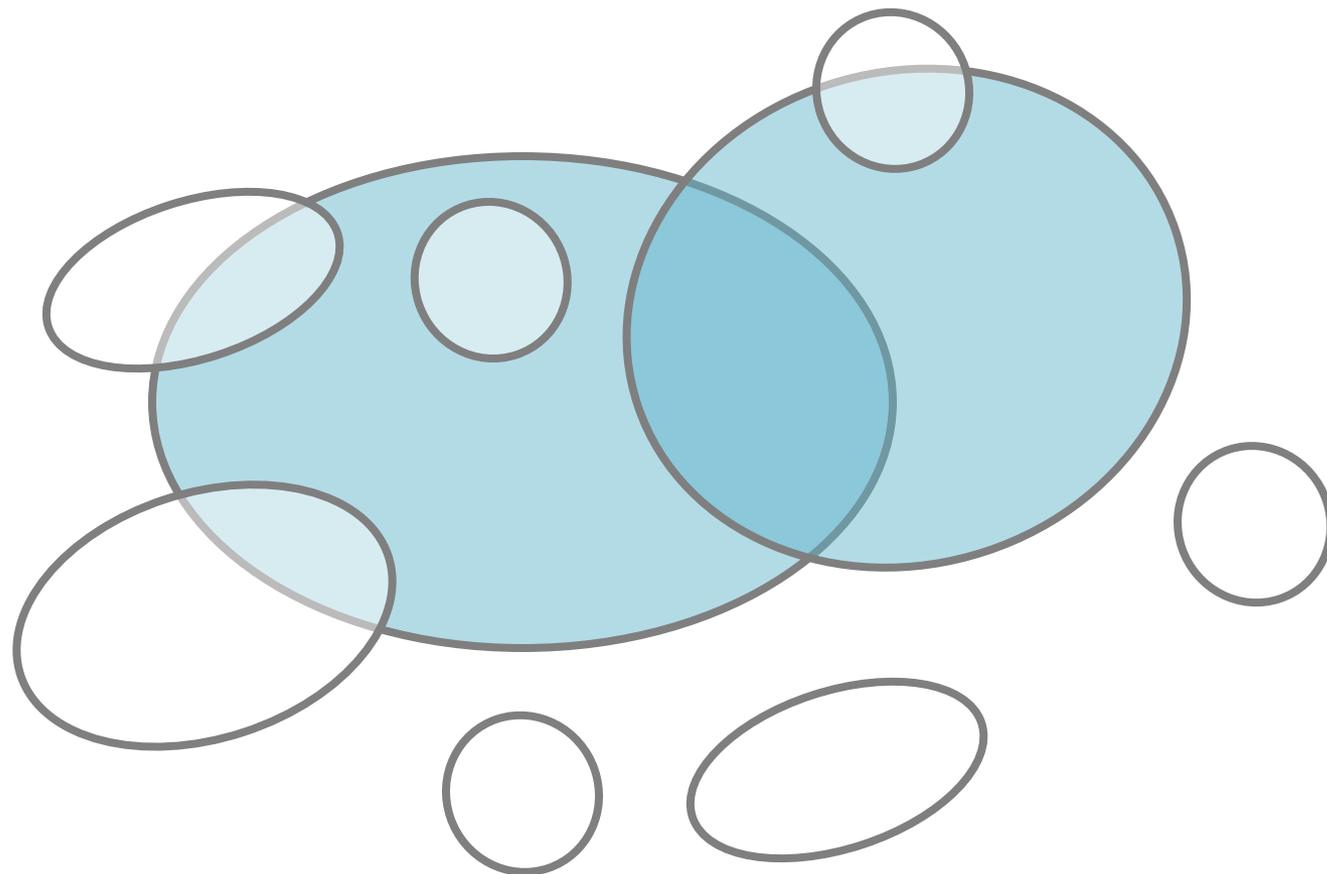
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- **Goal: Maximize the covered area**



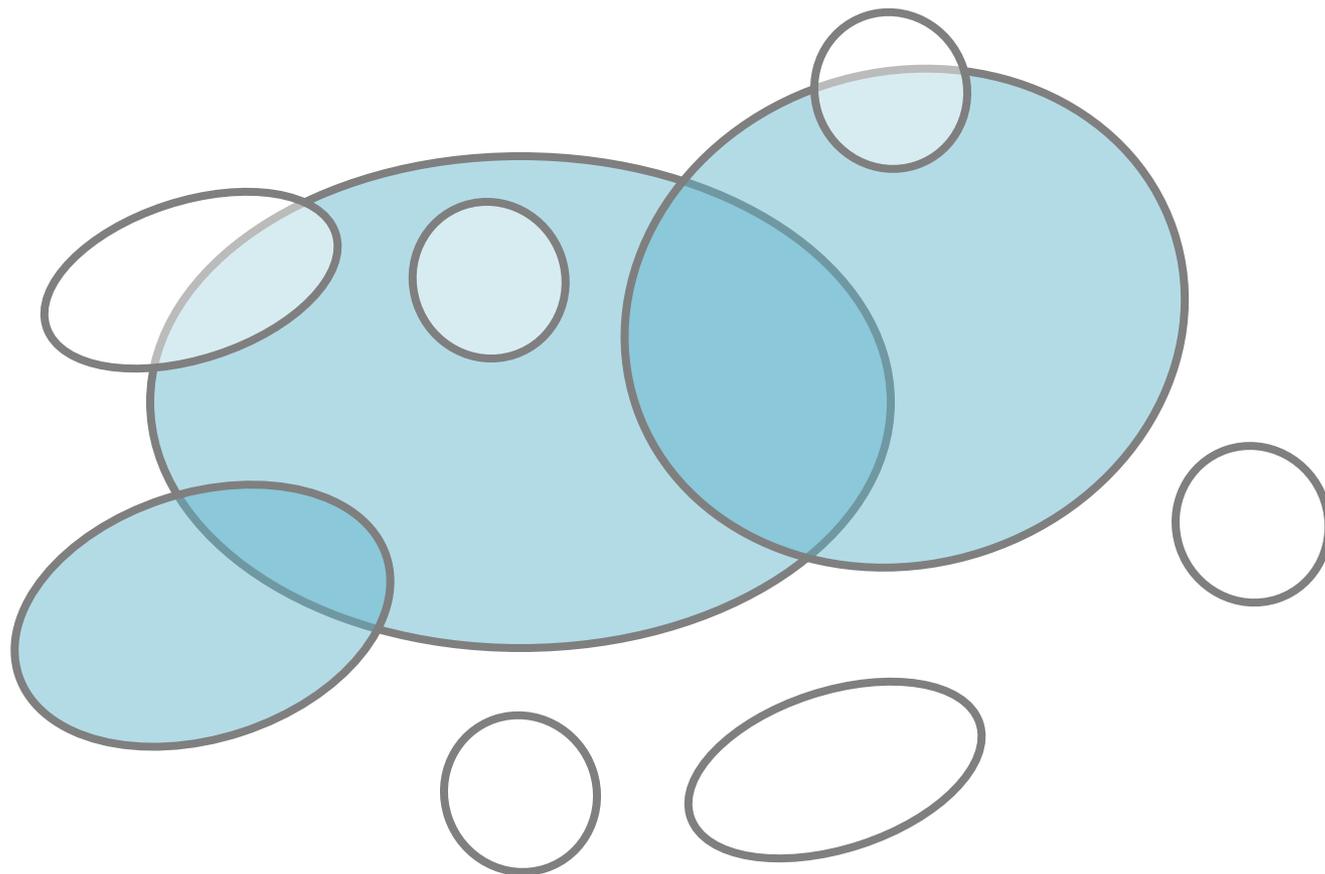
# Simple Greedy Heuristic

- **Goal: Maximize the covered area**



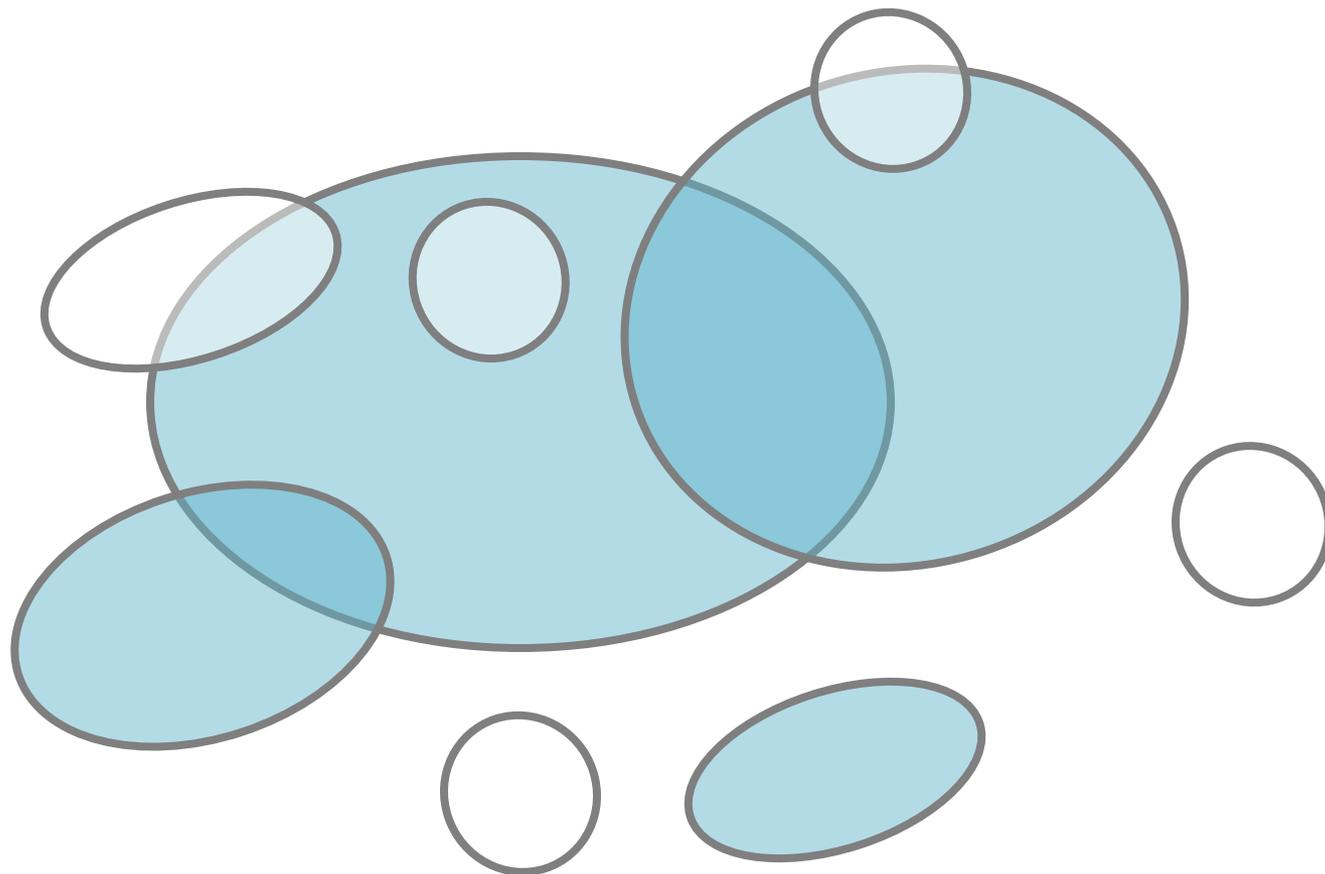
# Simple Greedy Heuristic

- **Goal: Maximize the covered area**

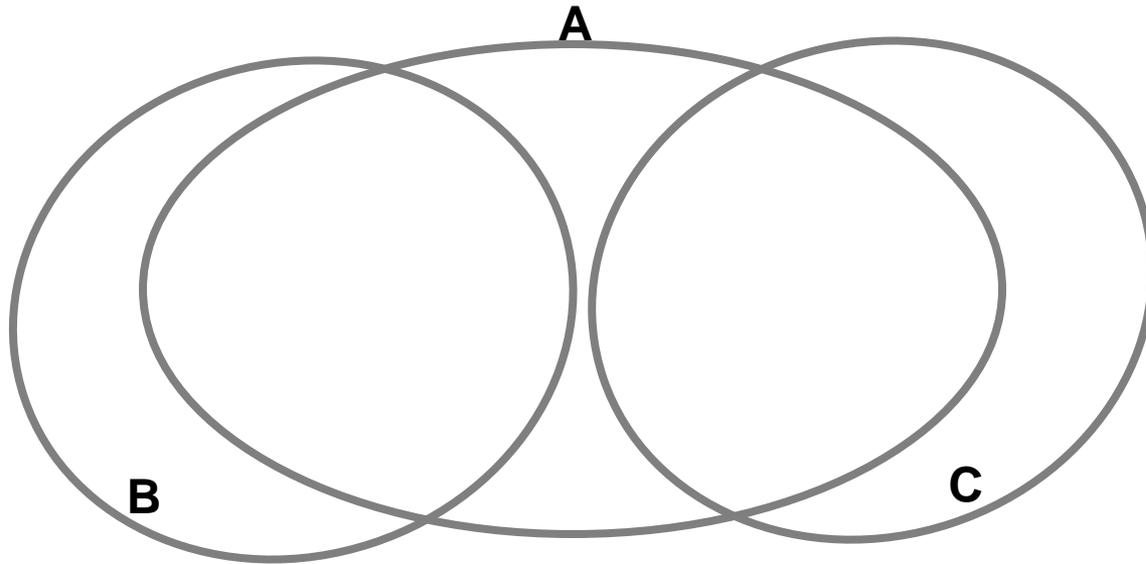


# Simple Greedy Heuristic

- **Goal: Maximize the covered area**



# When Greedy Heuristic Fails?



- **Goal:** Maximize the size of the covered area
- Greedy first picks A and then C
- But the optimal way would be to pick B and C

# Approximation Guarantee

- **Greedy produces a solution  $A$**   
**where:  $F(A) \geq (1-1/e)*OPT$  ( $F(A) \geq 0.63*OPT$ )**  
[Nemhauser, Fisher, Wolsey '78]
- **Claim holds for functions  $F(\cdot)$  with 2 properties:**
  - **$F$  is monotone:** (adding more docs doesn't decrease coverage)  
if  $A \subseteq B$  then  $F(A) \leq F(B)$  and  $F(\{\})=0$
  - **$F$  is submodular:**  
adding an element to a set gives less or equal improvement than adding it to one of its subsets

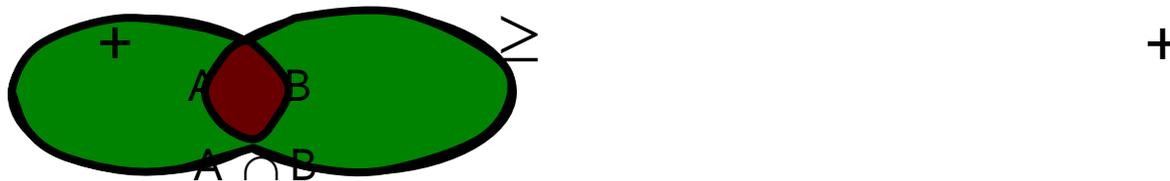
# Submodularity: Definition

## Definition:

- Set function  $F(\cdot)$  is called **submodular** if:

For all  $A, B \subseteq W$ :

$$F(A) + F(B) \geq F(A \cup B) + F(A \cap B)$$



# Submodularity: Or equivalently

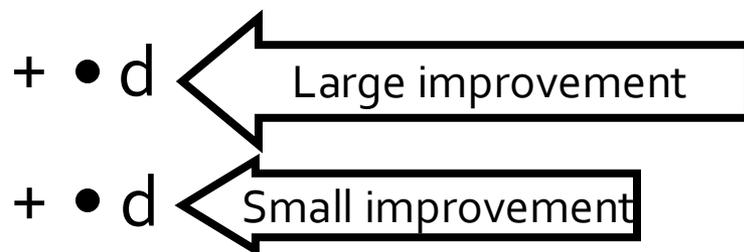
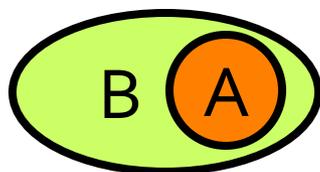
- **Diminishing returns** characterization

**Equivalent definition:**

- Set function  $F(\cdot)$  is called **submodular** if:

For all  $A \subseteq B$ :

$$\underbrace{F(A \cup \{d\}) - F(A)}_{\text{Gain of adding } d \text{ to a small set}} \geq \underbrace{F(B \cup \{d\}) - F(B)}_{\text{Gain of adding } d \text{ to a large set}}$$



# Example: Set Cover

- $F(\cdot)$  is **submodular**:  $A \subseteq B$

$$\underbrace{F(A \cup \{d\}) - F(A)}_{\text{Gain of adding } d \text{ to a small set}} \geq \underbrace{F(B \cup \{d\}) - F(B)}_{\text{Gain of adding } d \text{ to a large set}}$$

Gain of adding  $d$  to a small set

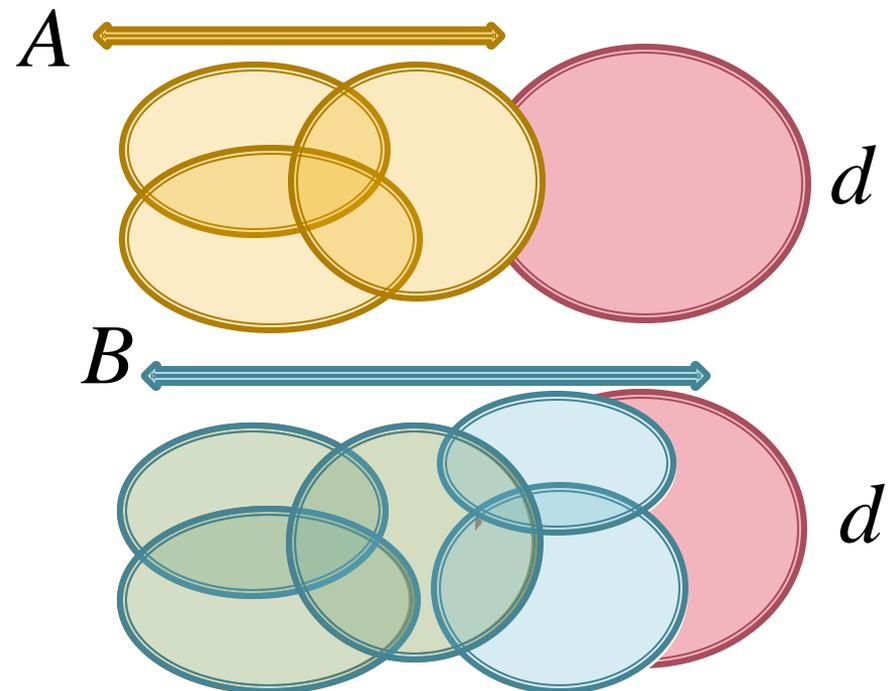
Gain of adding  $d$  to a large set

- **Natural example:**

- Sets  $d_1, \dots, d_m$
- $F(A) = |\cup_{i \in A} d_i|$   
(size of the covered area)

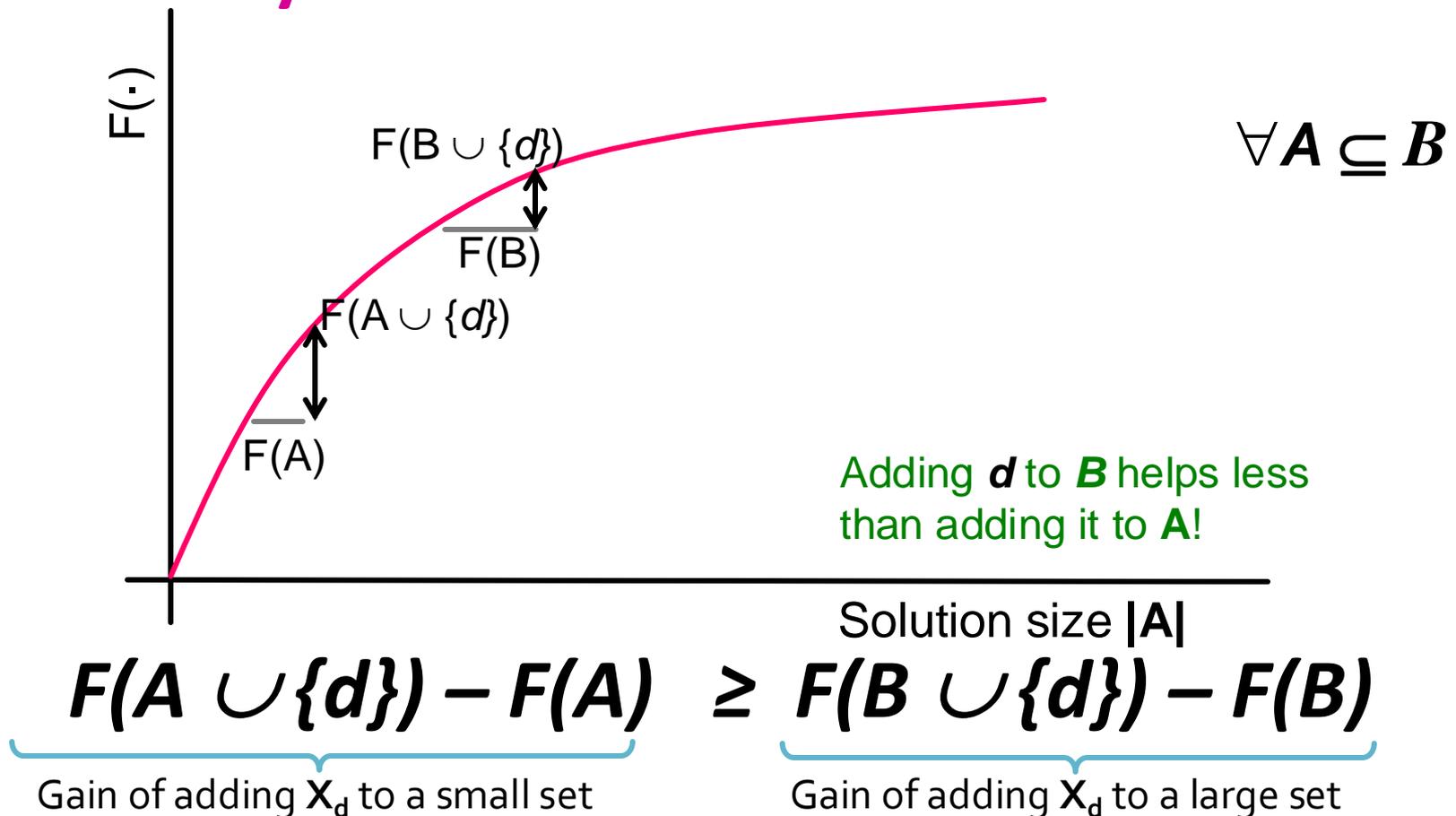
- **Claim:**

$F(A)$  is submodular!



# Submodularity– Diminishing returns

- Submodularity is discrete analogue of concavity



# Submodularity & Concavity

- **Marginal gain:**

$$\Delta_F(d|A) = F(A \cup \{d\}) - F(A)$$

- **Submodular:**

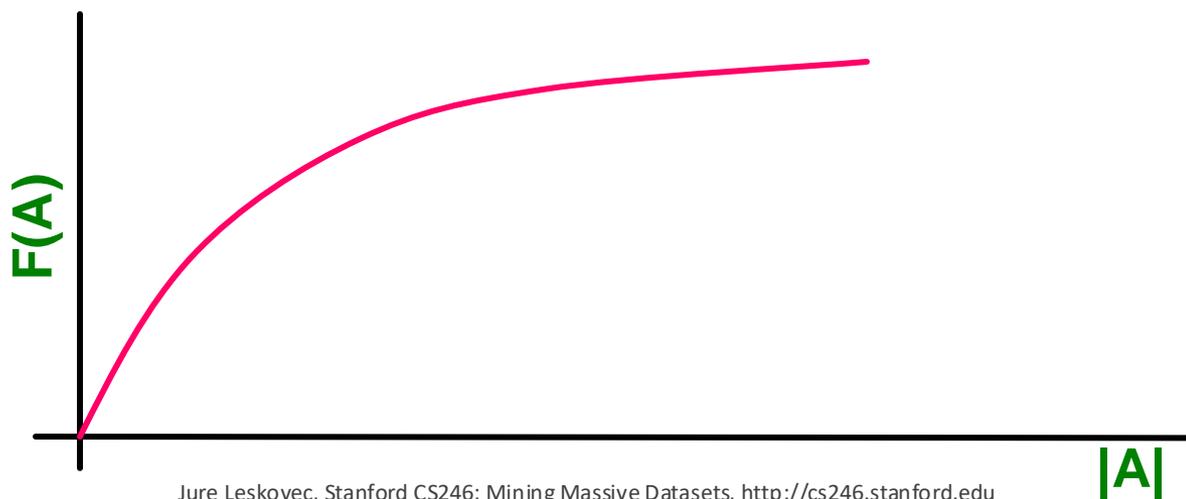
$$A \subseteq B$$

$$F(A \cup \{d\}) - F(A) \geq F(B \cup \{d\}) - F(B)$$

- **Concavity:**

$$a \leq b$$

$$f(a + d) - f(a) \geq f(b + d) - f(b)$$

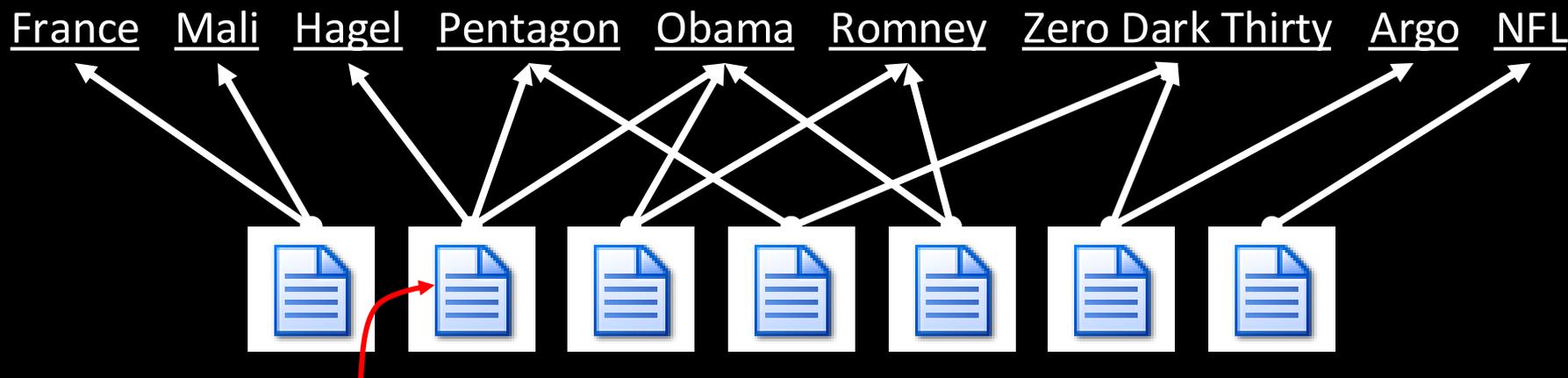


# Submodularity: Useful Fact

- Let  $F_1 \dots F_m$  be **submodular** and  $\lambda_1 \dots \lambda_m > 0$  then  $F(A) = \sum_{i=1}^m \lambda_i F_i(A)$  is **submodular**
  - **Submodularity is closed under non-negative linear combinations!**
- **This is an extremely useful fact:**
  - **Average of submodular functions is submodular:**  
 $F(A) = \sum_i P(i) \cdot F_i(A)$
  - **Multicriterion optimization:**  $F(A) = \sum_i \lambda_i F_i(A)$

# Back to our problem

- **Q: What is being covered?**
- **A: Concepts** (In our case: Named entities)

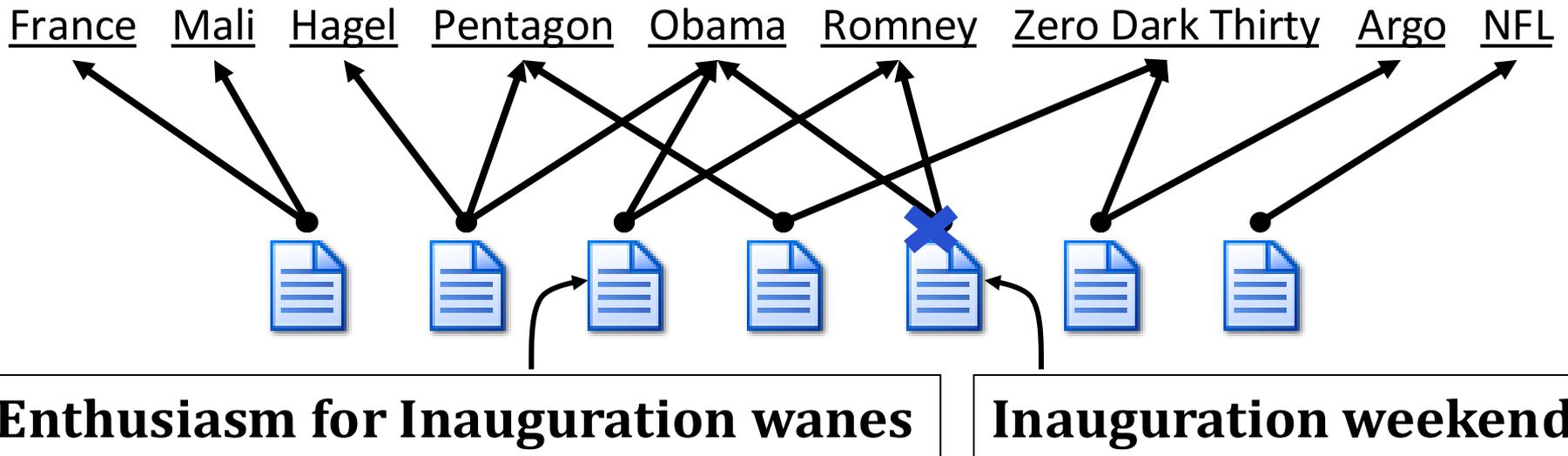


**Hagel expects fight**

- **Q: Who is doing the covering?**
- **A: Documents**

# Back to our Concept Cover Problem

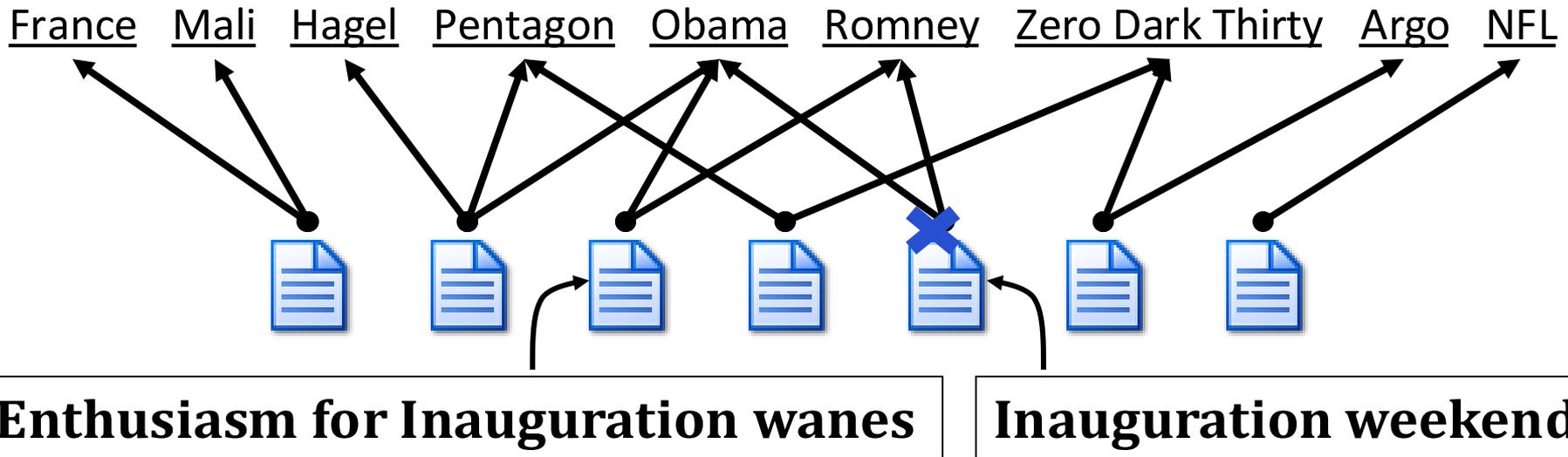
- **Objective:** pick  $k$  docs that cover most concepts



- $F(A)$ : the number of concepts covered by  $A$ 
  - *Elements...concepts, Sets ... concepts in docs*
  - $F(A)$  is submodular and monotone!
  - We can use **greedy algorithm** to optimize  $F$

# The Set Cover Problem

- **Objective:** pick  $k$  docs that cover most concepts



The good:

Penalizes redundancy

Submodular

The bad:

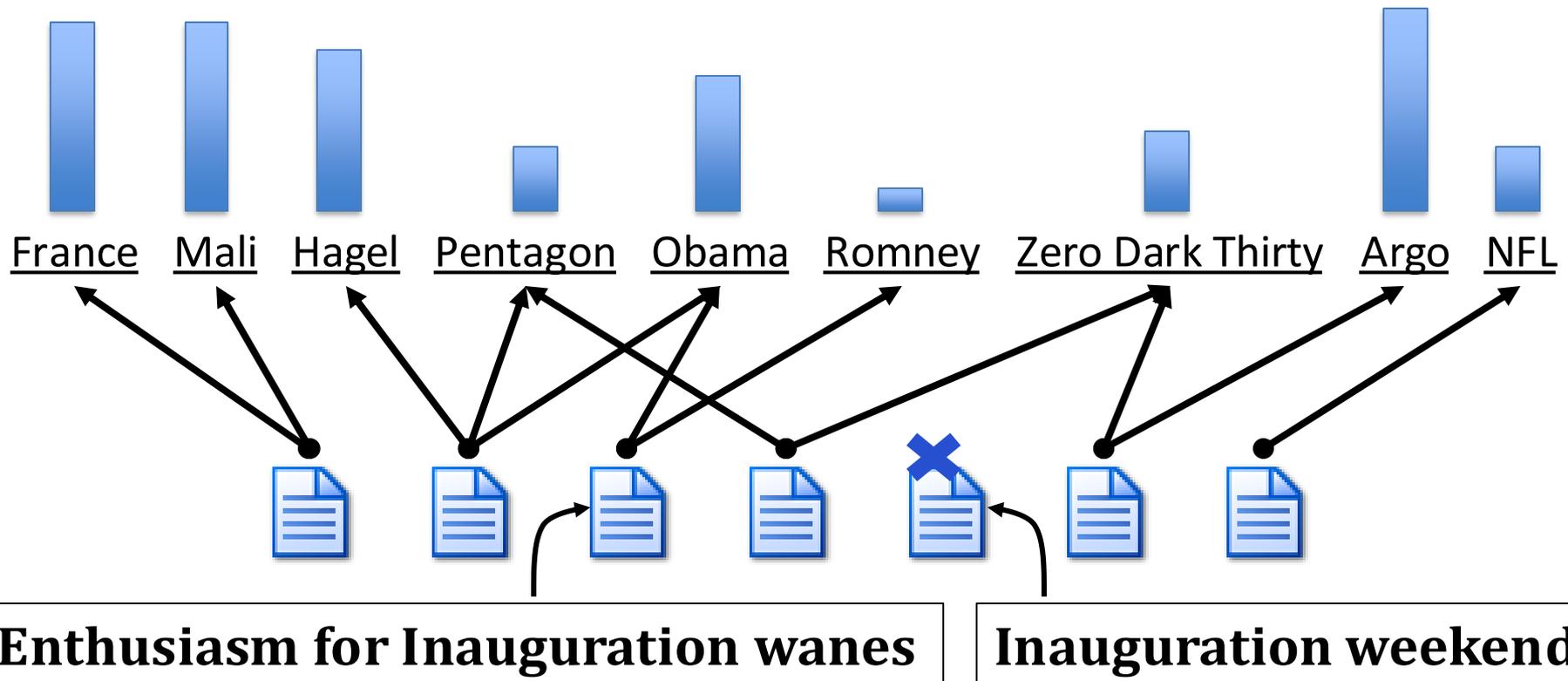
Concept importance?

All-or-nothing too harsh

# Probabilistic Set Cover

# Concept importance?

- **Objective:** pick  $k$  docs that cover most concepts

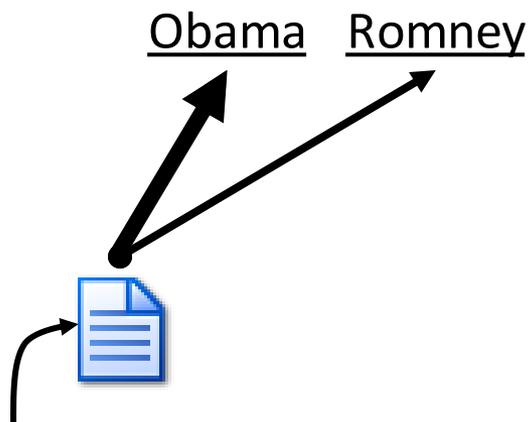


- Each concept  $c$  has importance weight  $w_c$

# All-or-nothing too harsh

- **Document coverage function**

$\text{cover}_d(c) =$  **probability** document **d** covers  
concept **c**  
[e.g., how strongly **d** covers **c**]



**Enthusiasm for Inauguration wanes**

# Probabilistic Set Cover

- **Document coverage function:**

$\text{cover}_d(c) =$  **probability** document **d** covers concept **c**

- $\text{Cover}_d(c)$  can also model how relevant is concept **c** for user **u**

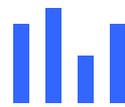
- **Set coverage function:**

$$\text{cover}_{\mathcal{A}}(c) = 1 - \prod_{d \in \mathcal{A}} (1 - \text{cover}_d(c))$$

- Prob. that at least one document in **A** covers **c**

- **Objective:**

$$\max_{\mathcal{A}: |\mathcal{A}| \leq k} F(\mathcal{A}) = \sum_c w_c \text{cover}_{\mathcal{A}}(c)$$

concept weights 

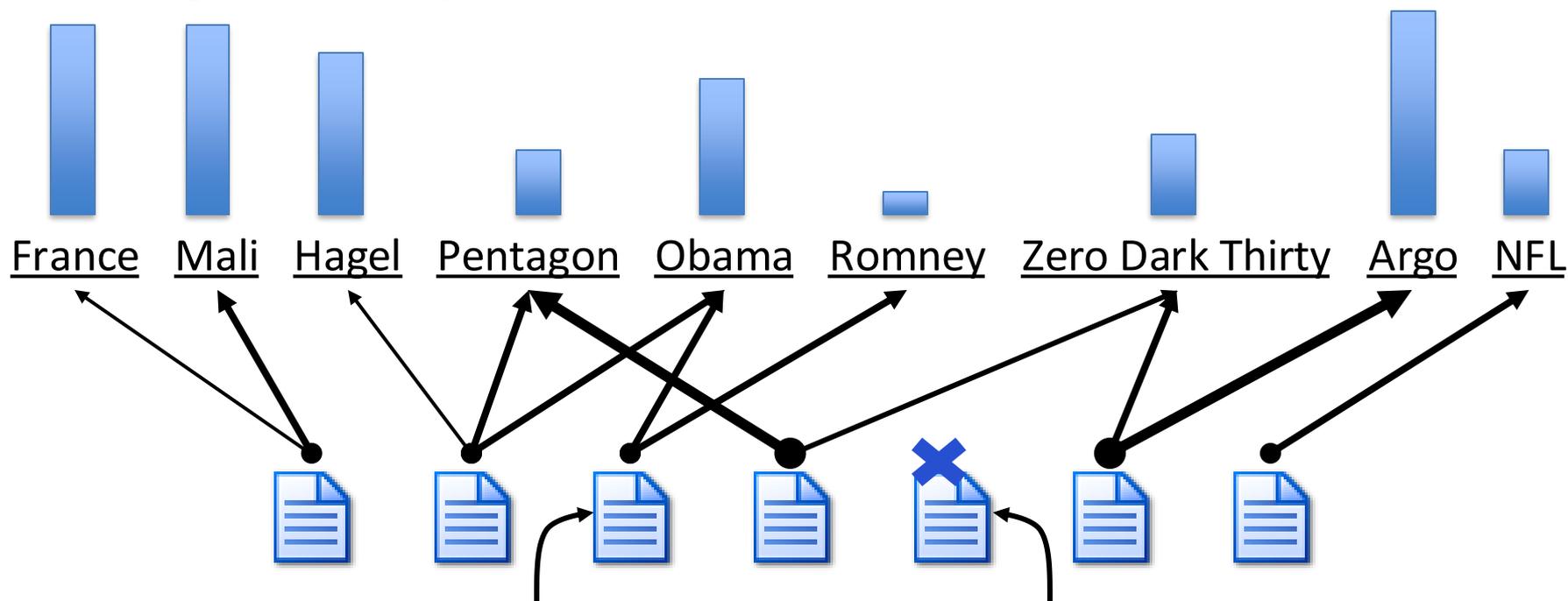
# Optimizing $F(\mathcal{A})$

$$\max_{\mathcal{A}: |\mathcal{A}| \leq k} F(\mathcal{A}) = \sum_c w_c \text{cover}_{\mathcal{A}}(c)$$

- The objective function is also **submodular**
  - Intuitively, it has a **diminishing returns** property
  - Greedy algorithm leads to a  $(1 - 1/e) \sim 63\%$  approximation, i.e., a **near-optimal** solution

# Summary: Probabilistic Set Cover

- **Objective:** pick  $k$  docs that cover most concepts



Enthusiasm for Inauguration wanes

Inauguration weekend

- Each concept  $c$  has importance weight  $w_c$
- Documents partially cover concepts:  $\text{cover}_d(c)$

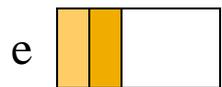
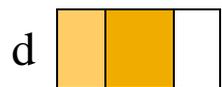
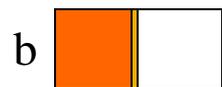
# Lazy Optimization of Submodular Functions

# Submodular Functions

## Greedy

Marginal gain:

$$F(A \cup x) - F(A)$$



Add document with  
highest marginal gain

- **Greedy algorithm is slow!**
  - At each iteration we need to re-evaluate marginal gains of **all remaining documents**
  - Runtime  $O(|D| \cdot K)$  for selecting  $K$  documents out of the set of  $D$  of them

# Speeding up Greedy

- **In round  $i$ :** So far we have  $A_{i-1} = \{d_1, \dots, d_{i-1}\}$ 
  - Now we pick  $d_i = \arg \max_{d \in V} F(A_{i-1} \cup \{d\}) - F(A_{i-1})$ 
    - Greedy algorithm maximizes the “marginal benefit”
 
$$\Delta_i(d) = F(A_{i-1} \cup \{d\}) - F(A_{i-1})$$

- **By submodularity property:**

$$F(A_i \cup \{d\}) - F(A_i) \geq F(A_j \cup \{d\}) - F(A_j) \text{ for } i < j$$

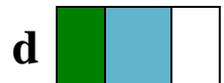
- **Observation: By submodularity:**

For every  $d \in D$

$$\Delta_i(d) \geq \Delta_j(d) \text{ for } i < j \text{ since } A_i \subseteq A_j$$

$$\Delta_i(d) \geq \Delta_j(d)$$

- **Marginal benefits  $\Delta_i(d)$  only shrink!**  
(as  $i$  grows)

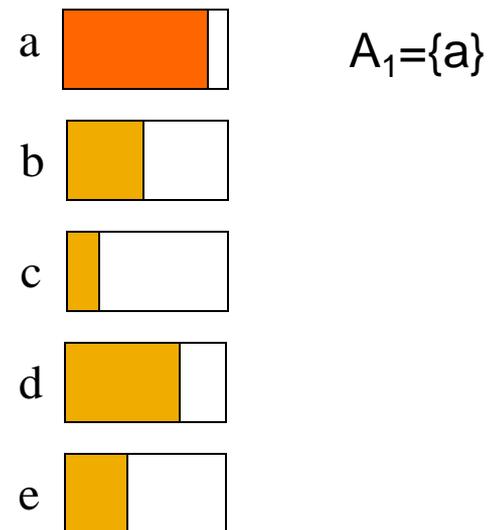


Selecting document  $d$  in step  $i$  covers more words than selecting  $d$  at step  $j$  ( $j > i$ )

# Lazy Greedy

- **Idea:**
  - Use  $\Delta_i$  as upper-bound on  $\Delta_j$  ( $j > i$ )
- **Lazy Greedy:**
  - Keep an ordered list of marginal benefits  $\Delta_j$  from previous iteration
  - Re-evaluate  $\Delta_j$  **only** for top element
  - Re-sort and prune

(Upper bound on)  
Marginal gain  $\Delta_1$

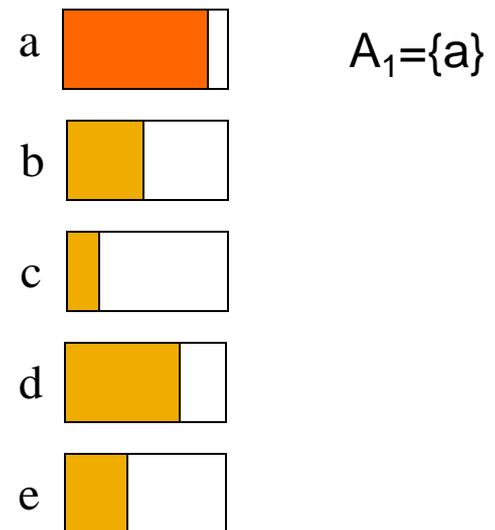


$$F(A \cup \{d\}) - F(A) \geq F(B \cup \{d\}) - F(B) \quad A \subseteq B$$

# Lazy Greedy

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- **Lazy Greedy:**
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Upper bound on  
Marginal gain  $\Delta_2$

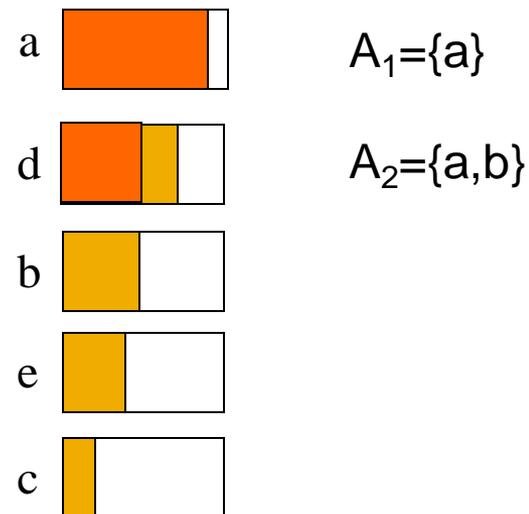


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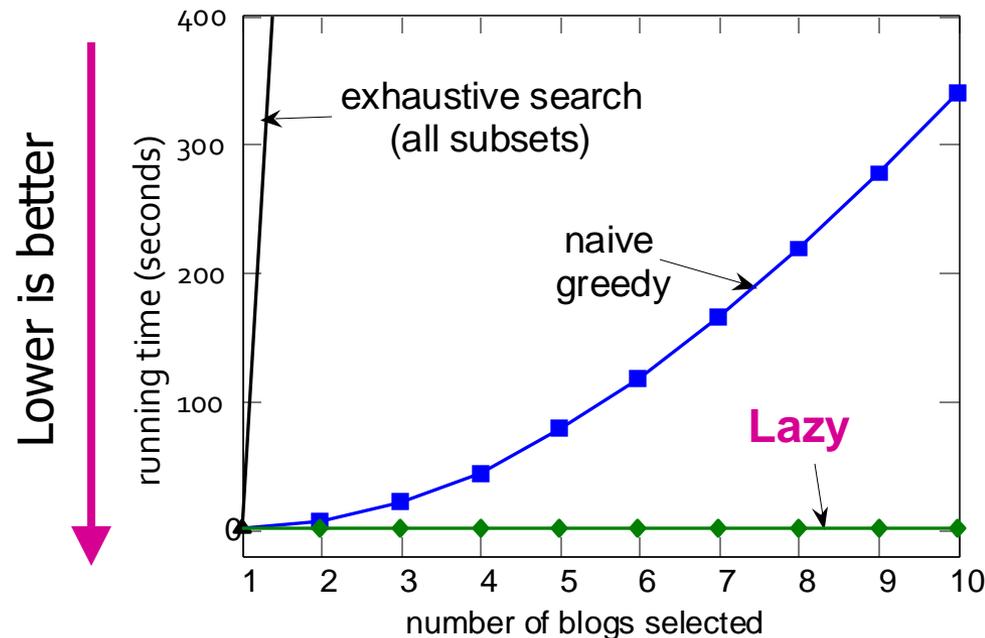
Upper bound on  
Marginal gain  $\Delta_2$



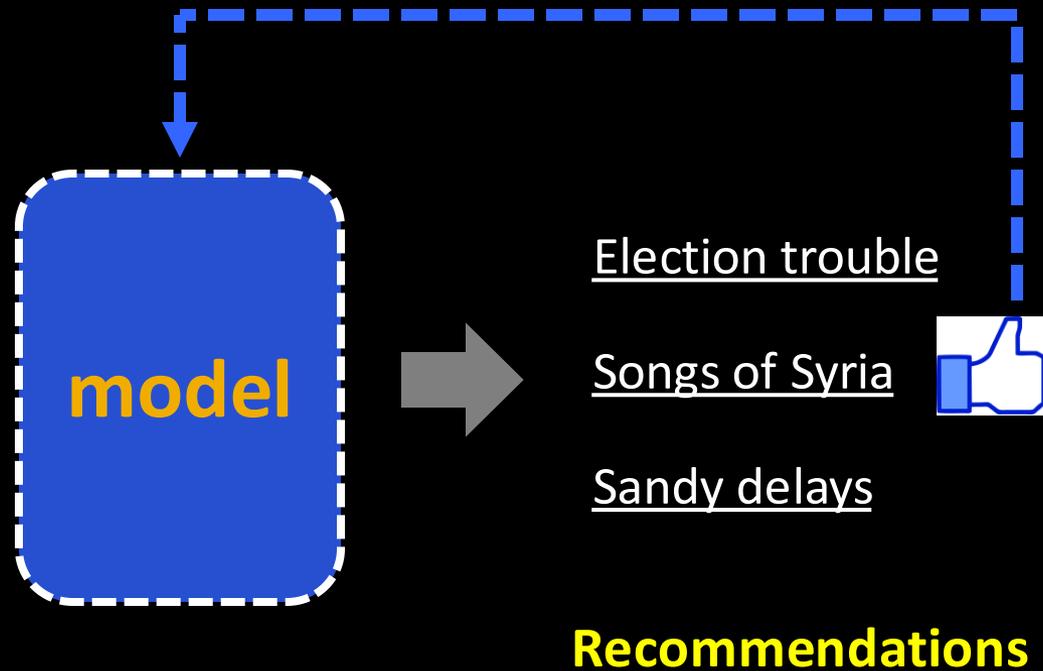
$$F(A \cup \{d\}) - F(A) \geq F(B \cup \{d\}) - F(B) \quad A \subseteq B$$

# Summary so far

- **Summary so far:**
  - Diversity can be formulated as a set cover
  - Set cover is submodular optimization problem
  - Can be (approximately) solved using greedy algorithm
  - Lazy-greedy gives significant speedup

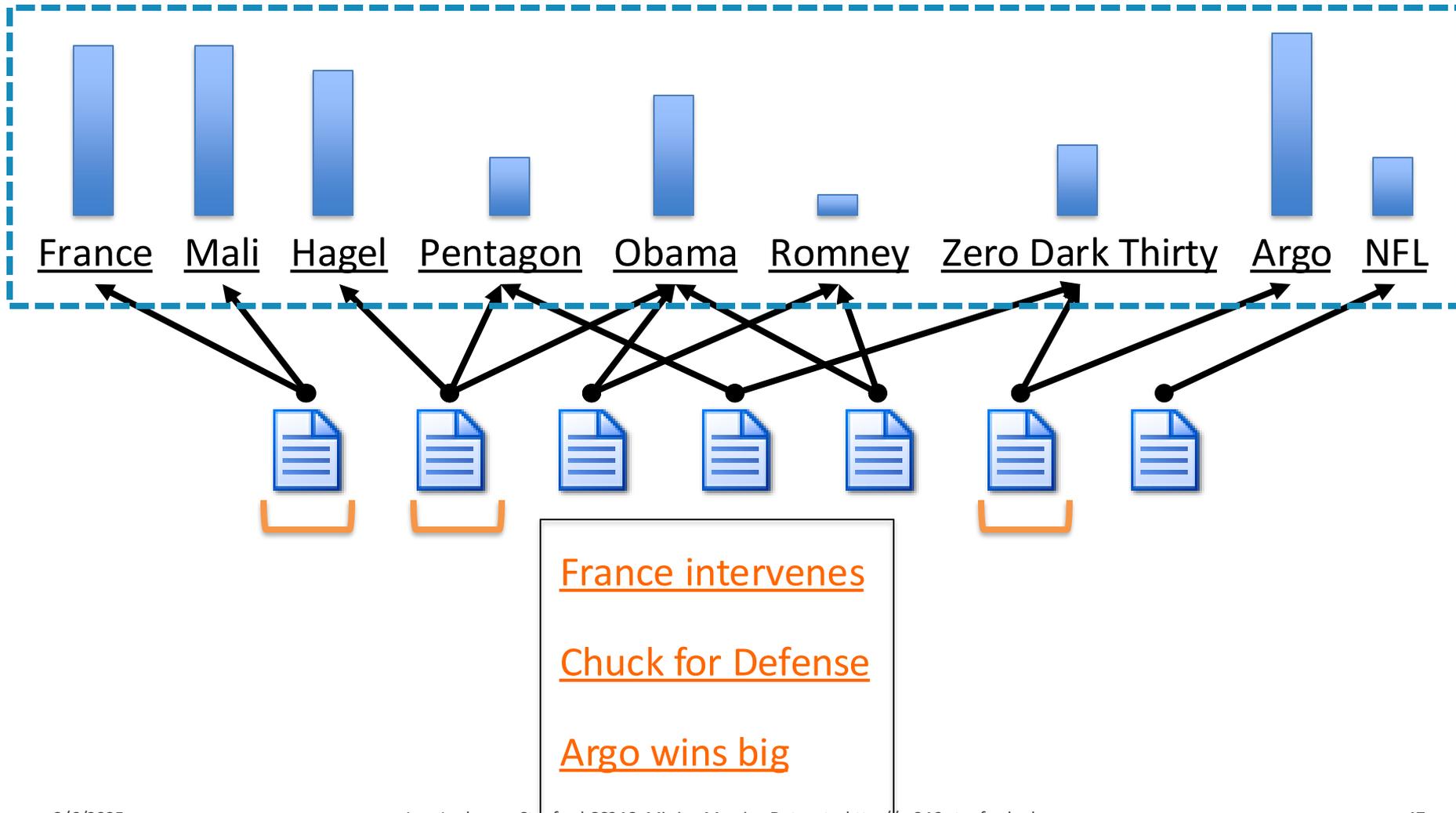


# But what about **personalization?**



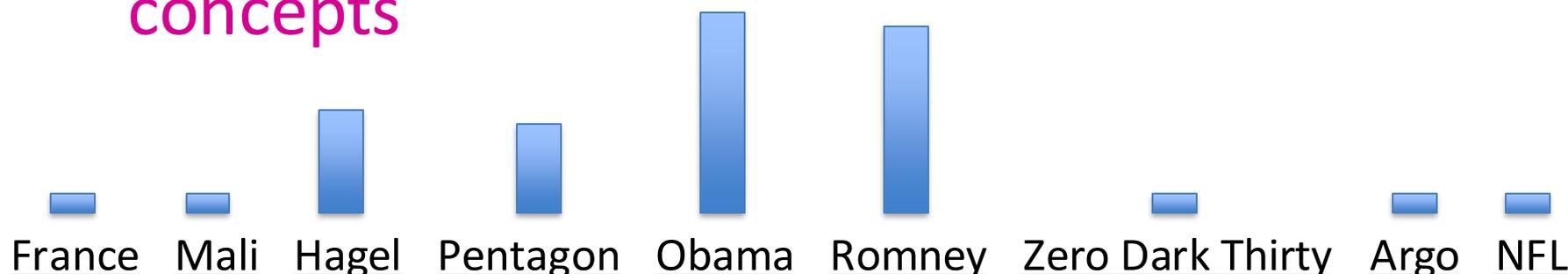
# Concept Coverage

We assumed same concept weighting for all users

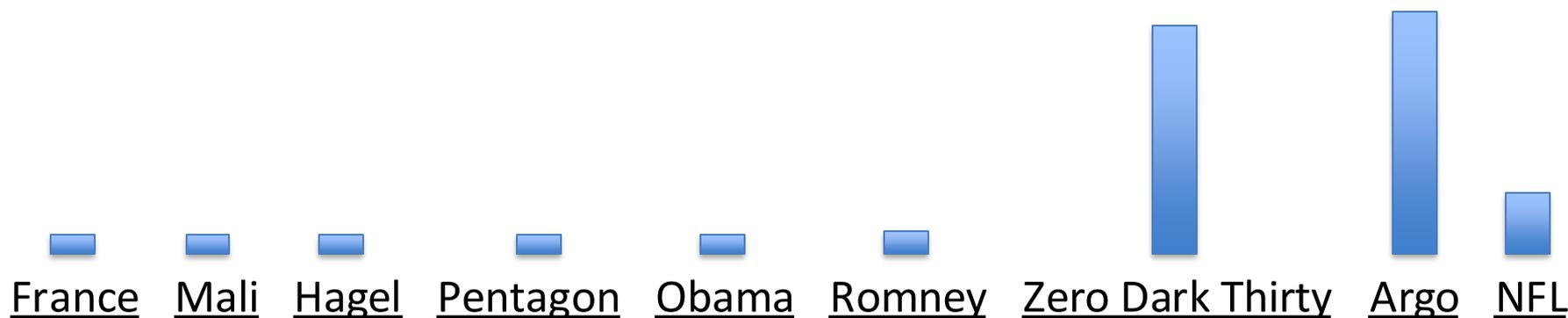


# Personal Concept Weights

- Each user has **different** preferences over concepts



**politico**



**movie buff**

# Personal concept weights

- Assume each user  $u$  has **different** preference vector  $w_c^{(u)}$  over concepts  $c$

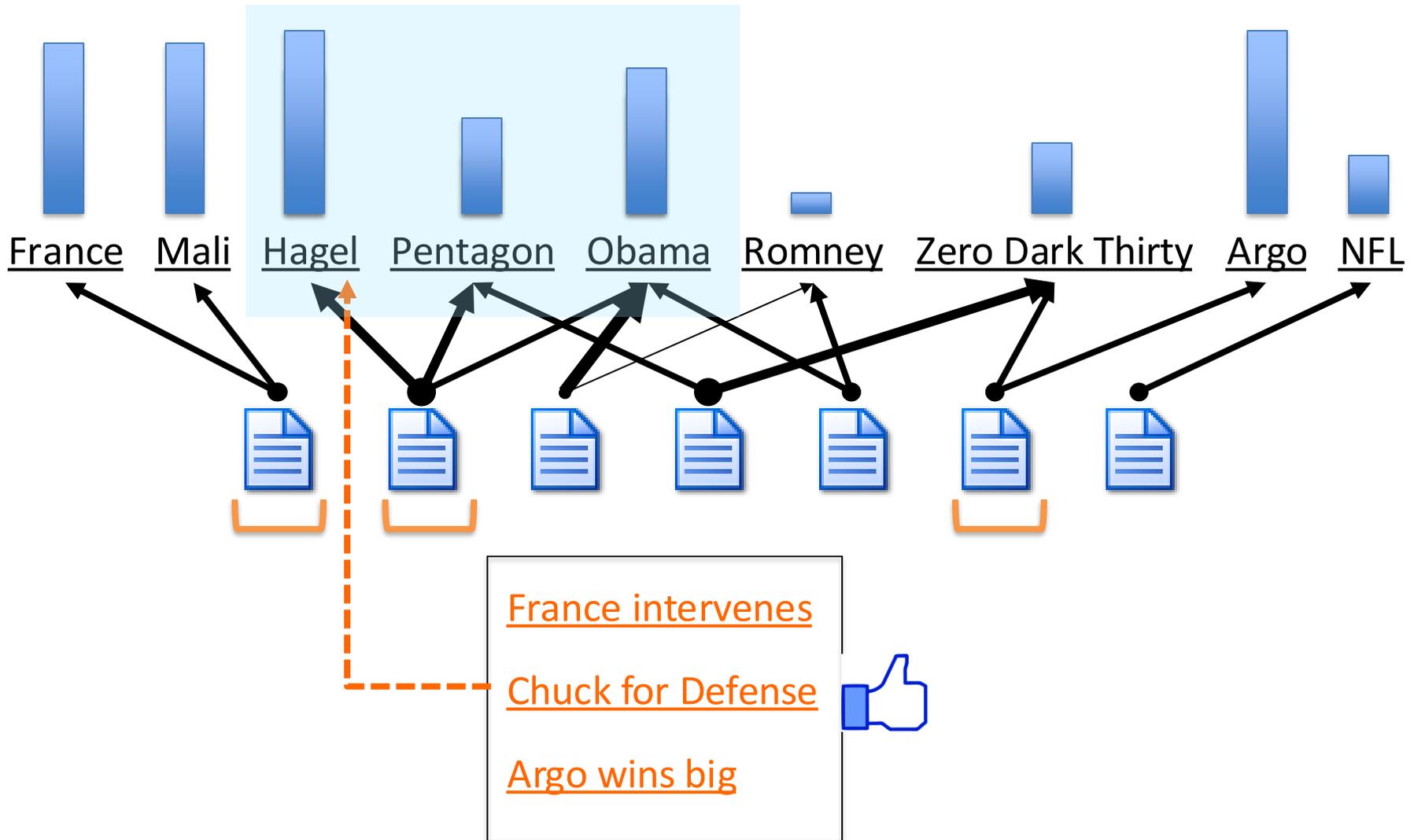
$$\max_{\mathcal{A}: |\mathcal{A}| \leq k} F(\mathcal{A}) = \sum_c w_c \text{cover}_{\mathcal{A}}(c)$$



$$\max_{\mathcal{A}: |\mathcal{A}| \leq k} F(\mathcal{A}) = \sum_c w_c^{(u)} \text{cover}_{\mathcal{A}}(c)$$

- Goal:** Learn personal concept weights from user feedback

# Interactive Concept Coverage



# Multiplicative Weights (MW)

- **Multiplicative Weights algorithm**
  - Assume each concept  $c$  has weight  $w_c$
  - We recommend document  $d$  and receive feedback, say  $r = +1$  or  $-1$
  - **Update the weights:**
    - **For each  $c \in X_d$  set  $w_c = \beta^r w_c$** 
      - If concept  $c$  appears in doc  $d$  and we received positive feedback  $r=+1$  then we increase the weight  $w_c$  by multiplying it by  $\beta$  ( $\beta > 1$ ) otherwise we decrease the weight (divide by  $\beta$ )
    - **Normalize weights so that  $\sum_c w_c = 1$**

# Summary of the Algorithm

- **Steps of the algorithm:**
  1. Identify **items** to recommend from
  2. Identify **concepts** [what makes items redundant?]
  3. **Weigh** concepts by general importance
  4. Define **item-concept coverage function**
  5. **Select** items using probabilistic set cover
  6. Obtain **feedback**, **update** weights