## Network Formation Processes: Power-law degrees and Preferential Attachment

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## Next Time: New Topics

## Observations

Small diameter,
Edge clustering
Patterns of signed edge creation

Viral Marketing, Blogosphere, Memetracking

## Models

## Algorithms

## Decentralized search

Models for predicting edge signs

Influence maximization,
Outbreak detection, LIM

## Scale-Free

Preferential attachment, Copying model

PageRank, Hubs and authorities

Link prediction, Supervised random walks

Community detection: Girvan-Newman, Modularity

## Network Formation Processes

What do we observe that needs explaining?

- Small-world model:
- Diameter
- Clustering coefficient
- Preferential Attachment:

- Node degree distribution
- What fraction of nodes has degree $\boldsymbol{k}$ (as a function of $\boldsymbol{k}$ )?
- Prediction from simple random graph models: $p(k)=$ exponential function of $\boldsymbol{k}$
- Observation: Often a power-law: $\boldsymbol{p}(\boldsymbol{k}) \propto \boldsymbol{k}^{-\boldsymbol{\alpha}}$


## Degree Distributions




## Node Degrees in Networks

- Take a network, plot a histogram of $\boldsymbol{P}(\boldsymbol{k})$ vs. $\boldsymbol{k}$



## Node Degrees in Networks

- Plot the same data on log-log scale:



## Node Degrees: Faloutsos³

- Internet Autonomous Systems
[Faloutsos, Faloutsos and Faloutsos, 1999]




## Node Degrees: Web

## - The World Wide Web [Broder et al., 2000]




## Node Degrees: Barabasi\&Albert

## - Other Networks [Barabasi-Albert, 1999]





Actor collaborations
Web graph
Power-grid

## Exponential vs. Power-Law



- Above a certain $\boldsymbol{x}$ value, the power law is always higher than the exponential!


## Exponential vs. Power-Law

- Power-law vs. Exponential on log-log and semi-log (log-lin) scales



## Exponential vs. Power-Law

Bell Curve



Power Law Distribution


## Power-Law Degree Exponents

- Power-law degree exponent is typically $2<\alpha<3$
- Web graph:
- $\alpha_{\text {in }}=2.1, \alpha_{\text {out }}=2.4$ [Broder et al. 00]
- Autonomous systems:
- $\alpha=2.4$ [Faloutsos ${ }^{3}$, 99]
- Actor-collaborations:
- $\alpha=2.3$ [Barabasi-Albert 00]
- Citations to papers:
- $\alpha \approx 3$ [Redner 98]
- Online social networks:
- $\alpha \approx 2$ [Leskovec et al. 07]




## Scale-Free Networks

- Definition:

Networks with a power-law tail in their degree distribution are called "scale-free networks"

- Where does the name come from?
- Scale invariance: There is no characteristic scale
- Scale invariance is that laws do not change if scales of length, energy, or other variables, are multiplied by a common factor
- Scale-free function: $f(a x)=a^{\lambda} f(x)$
- Power-law function: $\boldsymbol{f}(\boldsymbol{a x})=\boldsymbol{a}^{\lambda} \boldsymbol{x}^{\lambda}=\boldsymbol{a}^{\lambda} \boldsymbol{f}(\boldsymbol{x})$

$$
\begin{aligned}
& \log () \text { or } \operatorname{Exp}() \text { are not scale free! } \\
& f(a x)=\log (a x)=\log (a)+\log (x)=\log (a)+f(x) \\
& f(a x)=\exp (a x)=\exp (x)^{a}=f(x)^{a}
\end{aligned}
$$

## Power-Laws are Everywhere



Many other quantities follow heavy-tailed distributions

## Anatomy of the Long Tail

## ANATOMY OF THE LONG TAIL

Online services carry far more inventory than traditional retailers. Rhapsody, for example, offers 19 times as many songs as War-Mart's stock of 39,000 tunes. The appetite for Rhapsody's more obscure tunes (charted below in yellow) makes up the so-called Long Tail. Mesinwhile, even as consumers flock to mainstream books, music, and films (right), there is real demand for niche fare found only online.



## THE MEU GROUTH MARKET: OBSCURE PRODUCTS YOU CAN'T GET ANYYHERE BUT ONLINE



[^0]
## Not Everyone Likes Power-Laws ©



Mathematics of Power-Laws

## Heavy Tailed Distributions

- Degrees are heavily skewed: Distribution $P(X>x)$ is heavy tailed if:

$$
\lim _{x \rightarrow \infty} \frac{P(X>x)}{e^{-\lambda x}}=\infty
$$

- Note:
- Normal PDF: $p(x)=\frac{1}{\sqrt{2 \pi \sigma}} e^{-\frac{(x-\mu)^{2}}{2 \sigma^{2}}}$
- Exponential PDF: $p(x)=\lambda e^{-\lambda x}$
- then $P(X>x)=1-P(X \leq x)=e^{-\lambda x}$ are not heavy tailed!


## Heavy Tailed Distributions

- Various names, kinds and forms:
- Long tail, Heavy tail, Zipf's law, Pareto's law
- Heavy tailed distributions:
- $\mathrm{P}(\mathrm{x})$ is proportional to:
power law $\mid P(x) \propto x^{-\alpha}$
power law
with cutoff
stretched
exponential
log-normal $\frac{1}{x} \exp \left[-\frac{(\ln x-\mu)^{2}}{2 \sigma^{2}}\right]$

Mathematics of Power-laws

- What is the normalizing constant?

$$
p(x)=Z x^{-\alpha} \quad Z=?
$$

- $\boldsymbol{p}(\boldsymbol{x})$ is a distribution: $\int \boldsymbol{p}(\boldsymbol{x}) \boldsymbol{d} \boldsymbol{x}=\mathbf{1}$

Continuous approximation

- $1=\int_{x_{m}}^{\infty} p(x) d x=Z \int_{x_{m}}^{\infty} x^{-\alpha} d x$

$p(x)$ diverges as $x \rightarrow 0$ so $x_{m}$ is the minimum value of the power-law distribution $x \in\left[x_{m}, \infty\right]$
$-=-\frac{Z}{\alpha-1}\left[x^{-\alpha+1}\right]_{x_{m}}^{\infty}=-\frac{Z}{\alpha-1}\left[\infty^{1-\alpha}-x_{m}^{1-\alpha}\right]$
- $\Rightarrow Z=(\alpha-1) x_{m}^{\alpha-1}$

Need: $\alpha>1$ !

$$
p(x)=\frac{\alpha-1}{x_{m}}\left(\frac{x}{x_{m}}\right)^{-\alpha}
$$

Integral:
$\int(a x)^{n}=\frac{(a x)^{n+1}}{a(n+1)}$

## Mathematics of Power-laws

- What's the expected value of a power-law random variable $X$ ?
- $E[X]=\int_{x_{m}}^{\infty} x p(x) d x=Z \int_{x_{m}}^{\infty} x^{-\alpha+1} d x$
$=\frac{Z}{2-\alpha}\left[x^{2-\alpha}\right]_{x_{m}}^{\infty}=\frac{(\alpha-1) x_{m}^{\alpha-1}}{-(\alpha-2)}\left[\infty^{2-\alpha}-x_{m}^{2-\alpha}\right]$
Need: $\alpha>2$ !

$$
\Rightarrow E[X]=\frac{\alpha-1}{\alpha-2} x_{m}
$$

Power-law density:

$$
p(x)=\frac{\alpha-1}{x_{m}}\left(\frac{x}{x_{m}}\right)^{-\alpha}
$$

$$
Z=\frac{\alpha-1}{x_{m}^{1-\alpha}}
$$

## Mathematics of Power-Laws

- Power-laws have infinite moments!

$$
E[X]=\frac{\alpha-1}{\alpha-2} x_{m}
$$

- If $\alpha \leq 2: E[X]=\infty$
- If $\alpha \leq 3: \operatorname{Var}[X]=\infty$

In real networks $2<\alpha<3$ so:
$E[X]=$ const
$\operatorname{Var}[X]=\infty$

- Average is meaningless, as the variance is too high!
- Consequence: Sample average of $n$ samples from a power-law with exponent $\alpha$




Estimating Power-law
Exponent Alpha

## Estimating Power-Law Exponent $\alpha$

## Estimating $\alpha$ from data:

- (1) Fit a line on log-log axis using least squares:

Solve $\arg \min _{\alpha}(\log (y)-\alpha \log (x)+b)^{2}$


## Estimating Power-Law Exponent $\alpha$

## Estimating $\alpha$ from data:

- Plot Complementary CDF (CCDF) $\boldsymbol{P}(\boldsymbol{X} \geq \boldsymbol{x})$.

Then the estimated $\boldsymbol{\alpha}=\mathbf{1}+\boldsymbol{\alpha}^{\prime}$ where $\boldsymbol{\alpha}^{\prime}$ is the slope of $\boldsymbol{P}(\boldsymbol{X} \geq \boldsymbol{x})$.

- Fact: If $p(x)=P(X=x) \propto x^{-\alpha}$ then $P(X \geq x) \propto x^{-(\alpha-1)}$
- $P(X \geq x)=\sum_{j=x}^{\infty} p(j) \approx \int_{x}^{\infty} Z y^{-\alpha} d y=$
$=\frac{Z}{1-\alpha}\left[y^{1-\alpha}\right]_{x}^{\infty}=\frac{Z}{1-\alpha} x^{-(\alpha-1)}$


## Estimating Power-Law Exponent $\alpha$

## Estimating $\alpha$ from data:

- Use maximum likelihood approach:
- The log-likelihood of observed data $d_{i}$ :
- $L(\alpha)=\ln \left(\prod_{i}^{n} p\left(d_{i}\right)\right)=\sum_{i}^{n} \ln p\left(d_{i}\right)$
$=\sum_{i}^{n}\left(\ln (\alpha-1)-\ln \left(x_{m}\right)-\alpha \ln \left(\frac{d_{i}}{x_{m}}\right)\right)$
- Want to find $\alpha$ that max $L(\alpha)$ : Set $\frac{\mathrm{dL}(\alpha)}{\mathrm{d} \alpha}=0$
$-\frac{\mathrm{dL}(\alpha)}{\mathrm{d} \alpha}=0 \Rightarrow \frac{n}{\alpha-1}-\sum_{i}^{n} \ln \left(\frac{d_{i}}{x_{m}}\right)=0$
$-\Rightarrow \widehat{\boldsymbol{\alpha}}=\mathbf{1}+\boldsymbol{n}\left[\sum_{i}^{n} \ln \left(\frac{\boldsymbol{d}_{\boldsymbol{i}}}{\boldsymbol{x}_{\boldsymbol{m}}}\right)\right]^{\boldsymbol{1}} \quad \begin{aligned} & \text { Power-law density: } \\ & p(x)=\frac{\alpha-1}{x_{m}}\left(\frac{x}{x_{m}}\right)^{-\alpha}\end{aligned}$


## Flickr: Fitting Degree Exponent






## Why are Power-Laws Surprising

## Can not arise from sums of independent events!

- Recall: in $\boldsymbol{G}_{\boldsymbol{n} \boldsymbol{p}}$ each pair of nodes in connected independently with prob. $\boldsymbol{p}$
- $\boldsymbol{X}$... degree of node $\boldsymbol{v}$
- $\boldsymbol{X}_{\boldsymbol{w}} \ldots$ event that $\boldsymbol{w}$ links to $\boldsymbol{v}$
- $\boldsymbol{X}=\sum_{w} \boldsymbol{X}_{\boldsymbol{w}}$
- $E[X]=\sum_{w} E\left[X_{w}\right]=(n-1) p$
- Now, what is $P(X=k)$ ? Central limit theorem!
- $\boldsymbol{X}_{\mathbf{1}}, \ldots, \boldsymbol{X}_{\boldsymbol{n}}$ : random vars with mean $\boldsymbol{\mu}$, variance $\sigma^{2}$
- $\boldsymbol{S}_{\boldsymbol{n}}=\sum_{\boldsymbol{i}} \boldsymbol{X}_{\boldsymbol{i}}: E\left[S_{n}\right]=\boldsymbol{n} \boldsymbol{\mu}, \operatorname{Var}\left[S_{n}\right]=\boldsymbol{n} \boldsymbol{\sigma}^{\mathbf{2}}, \mathrm{SD}\left[S_{n}\right]=\boldsymbol{\sigma} \sqrt{\boldsymbol{n}}$
- $P\left(S_{n}=E\left[S_{n}\right]+x \cdot \operatorname{SD}\left[S_{n}\right]\right) \sim \frac{1}{2 \pi} \mathrm{e}^{-\frac{\mathrm{x}^{2}}{2}}$


## Random vs. Scale-free network



Random network
(Erdos-Renyi random graph)



Scale-free (power-law) network
Degree distribution is Power-law

Consequence of
Power-Law Degrees

## Consequence: Network Resilience

- How does network connectivity change as nodes get removed? [Albert et al. 00; Palmer et al. 01]
- Nodes can be removed:
- Random failure:

- Remove nodes uniformly at random
- Targeted attack:
- Remove nodes in order of decreasing degree
- This is important for robustness of the internet as well as epidemiology


## Network Resilience

- Networks with equal number of nodes and edges:
- ER random graph
- Scale-free network
- Study the properties of the network as an increasing fraction of nodes are removed
- Node selection:
- Random (this corresponds to random failures)
- Nodes with largest degrees (corresponds to targeted attacks)
- Measures:
- Fraction of nodes in the largest connected component
- Average shortest path length between nodes in the largest component


## Network Resilience: Connectivity

- Scale-free graphs are resilient to random attacks, but sensitive to targeted attacks:



## Network Resilience: Connectivity

What proportion of random nodes must be removed in order for the size ( S ) of the giant component to drop to 0 ?

$\gamma .$. degree exponent
K... maximum degree

- Infinite scale-free networks with $\gamma<3$ never break down under random node failures


## Network Resilience: Path Length



- Real networks are resilient to random failures
- $\mathrm{G}_{\mathrm{np}}$ has better resilience to targeted attacks
- E.g., we need to remove all pages of degree >5 to disconnect the Web. But this is a very small fraction of all web pages!


## Resilience in Real Networks



Source: Error and attack tolerance of complex networks. Réka Albert, Hawoong Jeong and Albert-László Barabási

## Zoom-in

The first few \% of nodes removed:

- $E: G_{n p}$
- SF: Scale-free
- Notice how targeted attacks very quickly disconnect the network


## Preferential Attachment Model

## Model: Preferential attachment

- Preferential attachment
[Price '65, Albert-Barabasi '99, Mitzenmacher '03]
- Nodes arrive in order 1,2,...,n
- At step $\boldsymbol{j}$, let $\boldsymbol{d}_{\boldsymbol{i}}$ be the degree of node $\boldsymbol{i}<\boldsymbol{j}$
- A new node $\boldsymbol{j}$ arrives and creates $\boldsymbol{m}$ out-links
- Prob. of $\boldsymbol{j}$ linking to a previous node $\boldsymbol{i}$ is proportional to degree $\boldsymbol{d}_{\boldsymbol{i}}$ of node $\boldsymbol{i}$

$$
P(j \rightarrow i)=\frac{d_{i}}{\sum_{k} d_{k}}
$$

## Rich Get Richer

- New nodes are more likely to link to nodes that already have high degree
- Herbert Simon's result:
- Power-laws arise from "Rich get richer" (cumulative advantage)
- Examples
- Citations [de Solla Price '65]: New citations to a paper are proportional to the number it already has
- Herding: If a lot of people cite a paper, then it must be good, and therefore I should cite it too
- Sociology: Matthew effect, http://en.wikipedia.org/wiki/Matthew effect
- "For whoever has will be given more, and they will have an abundance. Whoever does not have, even what they have will be taken from them."
- Eminent scientists often get more credit than a comparatively unknown researcher, even if their work is similar


## The Exact Model

We will analyze the following model:

- Nodes arrive in order $1,2,3, \ldots, n$
- When node $\boldsymbol{j}$ is created it makes a single out-link to an earlier node $\boldsymbol{i}$ chosen:
- 1) With prob. $p, j$ links to $i$ chosen uniformly at random (from among all earlier nodes)
- 2) With prob. $\mathbf{1}$ - $\boldsymbol{p}$, node $\boldsymbol{j}$ chooses $\boldsymbol{i}$ uniformly at random \& links to a random node $l$ that $i$ points to
- This is same as saying: With prob. $\mathbf{1 - p}$, node $\boldsymbol{j}$ links to node $\boldsymbol{l}$ with prob. proportional to $\boldsymbol{d}_{\boldsymbol{l}}$ (the in-degree of $\boldsymbol{l}$ )
- Our graph is directed: Every node has out-degree 1


## The Model Givens Power-Laws

- Claim: The described model generates networks where the fraction of nodes with in-degree $\boldsymbol{k}$ scales as:

$$
-\left(1+\frac{1}{\infty}\right)
$$

$P\left(d_{i}=k\right) \propto k$
$q$
where $q=1-p$

So we get power-law degree distribution with exponent:

$$
\alpha=1+\frac{1}{1-p}
$$

## Continuous Approximation

- Consider deterministic and continuous approximation to the degree of node $\boldsymbol{i}$ as a function of time $\boldsymbol{t}$
- $\boldsymbol{t}$ is the number of nodes that have arrived so far
- In-Degree $d_{i}(t)$ of node $\boldsymbol{i}(i=1,2, \ldots, n)$ is a continuous quantity and it grows deterministically as a function of time $\boldsymbol{t}$
- Plan: Analyze $d_{i}(t)$ - continuous in-degree of node $\boldsymbol{i}$ at time $\boldsymbol{t}(\boldsymbol{t}>\boldsymbol{i})$
- Note: Node $\boldsymbol{i}$ arrives to the graph at time $\boldsymbol{i}$


## Continuous Degree: What We Know

- Initial condition:
$\boldsymbol{d}_{\boldsymbol{i}}(\boldsymbol{t})=\mathbf{0}$, when $\boldsymbol{t}=\boldsymbol{i} \quad$ (node $i$ just arrived)
- Expected change of $d_{i}(t)$ over time:
- Node $\boldsymbol{i}$ gains an in-link at step $\boldsymbol{t}+\mathbf{1}$ only if a link from a newly created node $t+1$ points to it
- What's the probability of this event?
- With prob. $\boldsymbol{p}$ node $\boldsymbol{t}+\mathbf{1}$ links randomly:
- Links to our node $\boldsymbol{i}$ with prob. $\mathbf{1 / t}$
- With prob. $\mathbf{1}-\boldsymbol{p}$ node $\boldsymbol{t}+\mathbf{1}$ links preferentially:
- Links to our node $\boldsymbol{i}$ with prob. $\boldsymbol{d}_{\boldsymbol{i}}(\boldsymbol{t}) / \boldsymbol{t}$
- Prob. node $t+1$ links to $i$ is: $p \frac{1}{t}+(1-p) \frac{d_{i}(t)}{t}$ Note: each node creates exactly 1 edge. So after $t$ nodes/steps there are $t$ edges in total.


## Continuous Degree

- At $t=4$ node $i=4$ comes. It has out-degree of 1 to deterministically share with other nodes:

| Node $\boldsymbol{i}$ | $\boldsymbol{d}_{i}(\boldsymbol{t})$ | $\boldsymbol{d}_{i}(\boldsymbol{t}+\mathbf{1})$ |
| :---: | :---: | :---: |
| 0 | 0 | $=0+p \frac{1}{4}+(1-p) \frac{0}{4}$ |
| 1 | 2 | $=2+p \frac{1}{4}+(1-p) \frac{2}{4}$ |
| 2 | 0 | $=0+p \frac{1}{4}+(1-p) \frac{1}{4}$ |
| 3 | 1 | $=1+p \frac{1}{4}+(1-p) \frac{1}{4}$ |
| $\mathbf{4}$ | 1 | 0 |



- $d_{i}(t)-d_{i}(t-1)=\frac{\mathrm{d} d_{i}(t)}{\mathrm{d} t}=\mathrm{p} \frac{1}{t}+(1-p) \frac{d_{i}(t)}{t}$
- How does $d_{i}(t)$ evolve as $t \rightarrow \infty$ ?


## What is the rate of growth of $d_{i}$ ?

Expected change of $d_{i}(t)$ :

$$
\begin{array}{ll}
=\underbrace{\boldsymbol{d}_{i}(\boldsymbol{t}+\mathbf{1})-\boldsymbol{d}_{i}(\boldsymbol{t})}=\boldsymbol{p} \frac{1}{t}+(\mathbf{1}-\boldsymbol{p}) \frac{d_{i}(t)}{t} & \\
=\frac{\mathrm{d} d_{i}(t)}{\mathrm{d} t}=p \frac{1}{t}+(1-p) \frac{d_{i}(t)}{t}=\frac{p+q d_{i}(t)}{t} & \begin{array}{l}
\text { Divide by } \\
p+q d_{i}(t)
\end{array} \\
=\frac{1}{p+q d_{i}(t)} \mathrm{d} d_{i}(t)=\frac{1}{t} \mathrm{~d} t & \begin{array}{l}
\text { integrate }
\end{array} \\
=\int \frac{1}{p+q d_{i}(t)} \mathrm{d} d_{i}(t)=\int \frac{1}{t} \mathrm{~d} t & \begin{array}{l}
\text { Exponentiate } \\
\text { and let } A=e^{c}
\end{array} \\
=\frac{1}{q} \ln \left(p+q d_{i}(t)\right)=\ln t+c & \\
=p+q d_{i}(t)=e^{q c} t^{q} \Rightarrow \boldsymbol{d}_{\boldsymbol{i}}(\boldsymbol{t})=\frac{\mathbf{1}}{\boldsymbol{q}}\left((\boldsymbol{A} \boldsymbol{t})^{\boldsymbol{q}}-\boldsymbol{p}\right) & \mathrm{A}=?
\end{array}
$$

## What is the constant A?

## What is the value of constant A?

$$
d_{i}(t)=\frac{1}{q}\left(A t^{q}-p\right)
$$

- We know: $d_{i}(i)=0$
- So: $d_{i}(i)=\frac{1}{q}\left((A i)^{q}-p\right)=0$
$\Rightarrow A=\frac{p}{i q}$
- And so $\Rightarrow d_{i}(t)=\frac{p}{q}\left(\left(\frac{t}{i}\right)^{q}-1\right)$

Observation: Old nodes (small $i$ values) have higher in-degrees $d_{i}(t)$


[^0]:    

