# Link Prediction and Network Inference

CS224W: Social and Information Network Analysis
Jure Leskovec, Stanford University

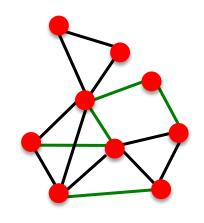
http://cs224w.stanford.edu



#### **Link Prediction in Networks**

#### The link prediction task:

• Given  $G[t_0, t_0]$  a graph on edges up to time  $t_0$  output a ranked list L of links (not in  $G[t_0, t_0]$ ) that are predicted to appear in  $G[t_1, t_1]$ 



#### Evaluation:

- $n=|E_{new}|$ : # new edges that appear during the test period  $[t_1,t_1]$
- lacktriangle Take top n elements of L and count correct edges

# **Link Prediction via Proximity**

#### Predict links in a evolving collaboration network

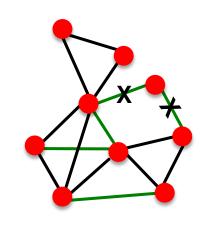
	training period			Core		
	authors	papers	$collaborations^1$	authors	$ E_{old} $	$ E_{new} $
astro-ph	5343	5816	41852	1561	6178	5751
cond-mat	5469	6700	19881	1253	1899	1150
gr-qc	2122	3287	5724	486	519	400
hep-ph	5414	10254	47806	1790	6654	3294
hep-th	5241	9498	15842	1438	2311	1576

- Core: Since network data is very sparse
  - Consider only nodes with in-degree and out-degree of at least 3

# **Link Prediction via Proximity**

#### Methodology:

- For every pair of nodes (x,y) compute proximity c
  - # of common neighbors c(x,y) of x and y
- Sort pairs by the decreasing score
- $E_{new}^* := E_{new} \cap (\mathsf{Core} \times \mathsf{Core})$ 
  - (only condier/predict edges where both endpoints are in the core)
- Predict top n pairs as new links



# Link Prediction via Proximity

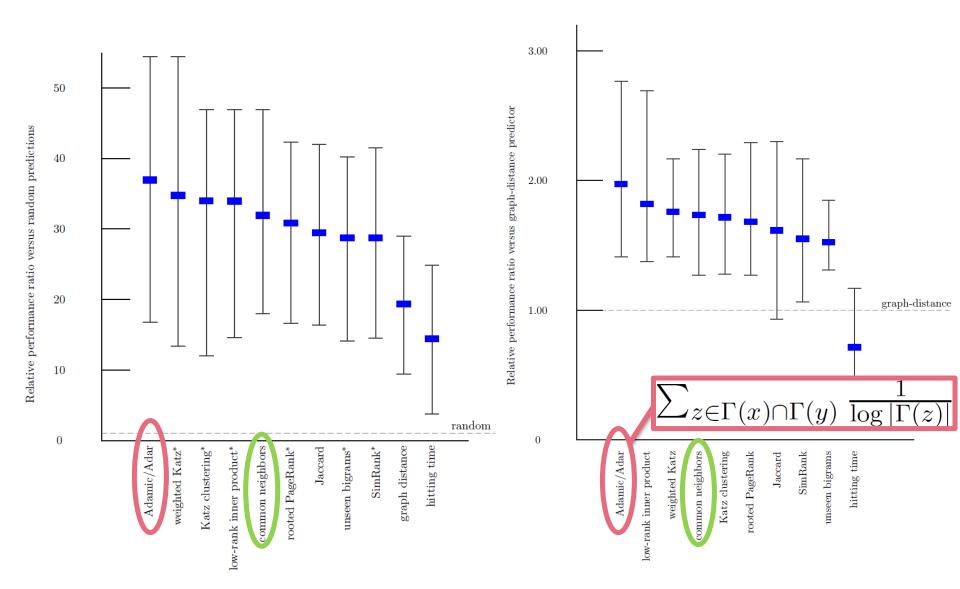
#### For every pair of nodes (x,y) compute:

graph distance	(negated) length of shortest path between $x$ and $y$
common neighbors	$ \Gamma(x) \cap \Gamma(y) $
Jaccard's coefficient	$\frac{ \Gamma(x)\cap\Gamma(y) }{ \Gamma(x)\cup\Gamma(y) }$
Adamic/Adar	$\sum_{z \in \Gamma(x) \cap \Gamma(y)} \frac{1}{\log  \Gamma(z) }$
preferential attachment	$ \Gamma(x)  \cdot  \Gamma(y) $
$\mathrm{Katz}_{eta}$	$\sum_{\ell=1}^\infty eta^\ell \cdot  paths_{x,y}^{\langle\ell angle} $

```
where \mathsf{paths}_{x,y}^{\langle\ell\rangle} := \{ \mathsf{paths} \text{ of length exactly } \ell \text{ from } x \text{ to } y \} weighted: \mathsf{paths}_{x,y}^{\langle 1 \rangle} := \mathsf{number of collaborations between } x, y. unweighted: \mathsf{paths}_{x,y}^{\langle 1 \rangle} := 1 \text{ iff } x \text{ and } y \text{ collaborate.}
```

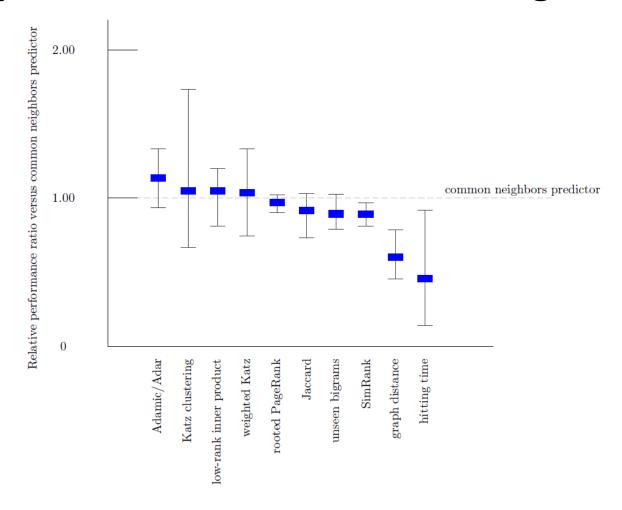
 $\Gamma(x)$  ... degree of node x

### Results: Improvement over Random



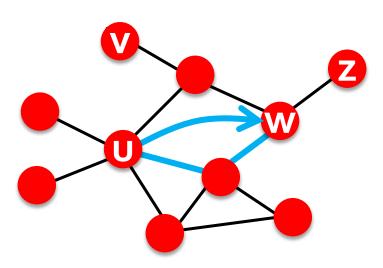
### Results: Common Neighbors

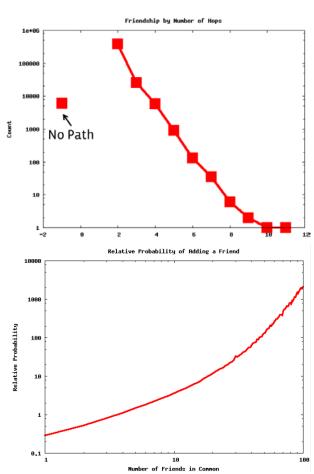
#### Improvement over #common neighbors



# **Supervised Link Prediction**

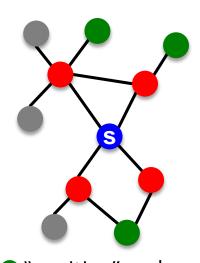
- How to learn to predict new friends?
  - Facebook's People You May Know
  - Let's look at the data:
    - 92% of new friendships on FB are friend-of-a-friend
    - More common friends helps





### **Supervised Link Prediction**

- Recommend a list of possible friends
- Supervised machine learning setting:
  - Training example:
    - For every node s have a list of nodes she will create links to  $\{v_1 \dots v_k\}$ 
      - Use FB network from May 2011 and  $\{v_1...v_k\}$  are the new friendships you created since then
  - Task:
    - For a given node s rank nodes  $\{v_1 \dots v_k\}$  higher than other nodes in the network
- Supervised Random Walks based on work by Agarwal&Chakrabarti



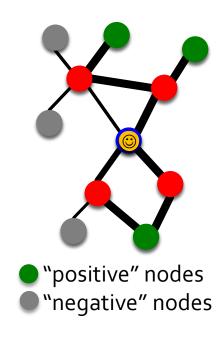
"positive" nodes"negative" nodes

#### **Green nodes**

are the nodes to which s creates links in the future

# **Supervised Link Prediction**

- How to combine node/edge attributes and the network structure?
  - Learn a strength of each edge based on:
    - Profile of user u, profile of user v
    - Interaction history of u and v
  - Do a PageRank-like random walk from s to measure the "proximity" between s and other nodes
  - Rank nodes by their "proximity" (i.e., visiting prob.)

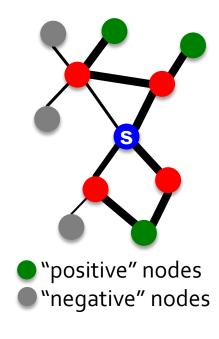


### Supervised Random Walks

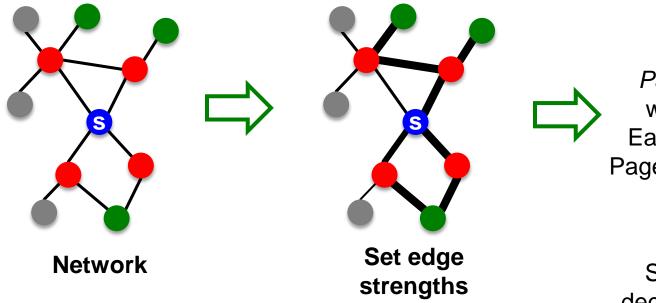
- Let s be the center node
- Let  $f_w(u,v)$  be a function that assigns a strength to each edge:

$$a_{uv} = f_{\beta}(u, v) = exp(-\beta^{\mathrm{T}} \Psi_{uv})$$

- $\Psi_{uv}$  is a feature vector
  - Features of node u
  - Features of node v
  - Features of edge (u,v)
- ( $\beta$  is the parameter vector we want to learn!)
- Do Random Walk with Restarts from s where transitions are according to edge strengths  $a_{uv}$



#### **SRW: Prediction**



 $a_{uv} = f_{\beta}(u,v)$ 

Personalized
PageRank on the
weighted graph.
Each node u gets a
PageRank proximity p<sub>u</sub>



Sort nodes by the decreasing PageRank proximity p<sub>u</sub>



Recommend top k nodes with the highest proximity  $p_{ij}$  to node s

- How to estimate edge strengths?
  - How to set parameters  $\beta$  of  $f_{\beta}(u,v)$ ?

# Personalized PageRank

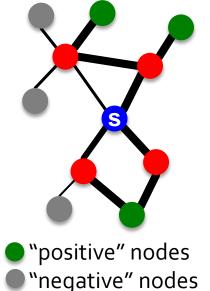
- a<sub>uv</sub> .... Strength of edge (u,v,)
- Random walk transition matrix:

$$Q'_{uv} = \begin{cases} \frac{a_{uv}}{\sum_{w} a_{uw}} & \text{if } (u, v) \in E, \\ 0 & \text{otherwise} \end{cases}$$

PageRank transition matrix:

$$Q_{ij} = (1 - \alpha)Q'_{ij} + \alpha \mathbf{1}(j = s)$$

- with prob.  $\alpha$  jump back to s
- Compute PageRank vector:  $p = p^T Q$
- Rank nodes u by  $p_u$



### **The Optimization Problem**

- Each node u has a score  $p_u$
- Positive nodes  $D = \{d_1, ..., d_k\}$
- Negative nodes L = {the rest}
- What do we want?

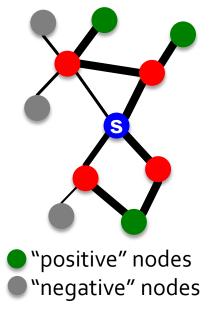
Want to find  $\beta$  such that  $p_l < p_d$ 

$$\min_{\beta} F(\beta) = ||\beta||^2$$

such that

$$\forall d \in D, l \in L: p_l < p_d$$

- The exact solution to the above problem may not exist
- So we make the constrains "soft" (i.e., optional)



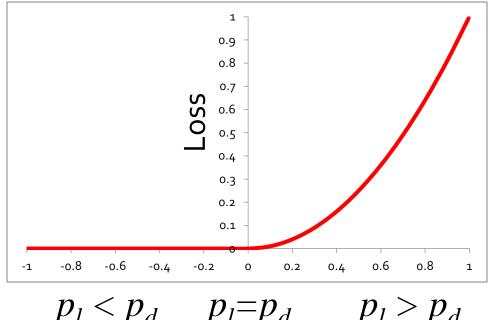
### Making Constraints "Soft"

#### Want to minimize:

$$\min_{\beta} F(\beta) = \sum_{d \in D, l \in L} h(p_l - p_d) + \lambda ||\beta||^2$$

• Loss: h(x)=0 if x<0,  $x^2$  else  $h(x)=\max\{x,0\}^2$ 

$$h(x) = \max\{x, 0\}^2$$

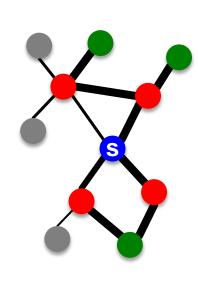


# Solving the Problem: Intuition

How to minimize F?

$$\min_{\beta} F(\beta) = \sum_{d \in D, l \in L} h(p_l - p_d) + \lambda ||\beta||^2$$

- Both  $p_l$  and  $p_d$  depend on  $\beta$ 
  - Given  $\beta$  assign edge weights  $a_{uv} = f_{\beta}(u, v)$
  - Using transition matrix  $Q=[a_{uv}]$ compute PageRank scores  $p_u$
  - Rank nodes by the PageRank score
- Want to find  $\beta$  such that  $p_l < p_d$



#### **Gradient Descent**

#### How to minimize F?

$$\min_{\beta} F(\beta) = \sum_{d \in D, l \in L} h(p_l - p_d) + \lambda ||\beta||^2$$

Take the derivative!

$$\begin{split} \frac{\partial F(\beta)}{\partial \beta} &= \sum_{l,d} \frac{\partial h(p_l - p_d)}{\partial \beta} + 2\lambda\beta \\ &= \sum_{l,d} \frac{\partial h(p_l - p_d)}{\partial (p_l - p_d)} \left(\frac{\partial p_l}{\partial \beta}\right) + \frac{\partial p_d}{\partial \beta} + 2\lambda\beta \end{split} \qquad \begin{array}{c} h(x) = \max\{x,0\}^2 \\ \text{Easy} \end{array}$$

We know:

$$p = p^T Q$$
 i.e.  $p_u = \sum_j p_j Q_{ju}$ 

So:

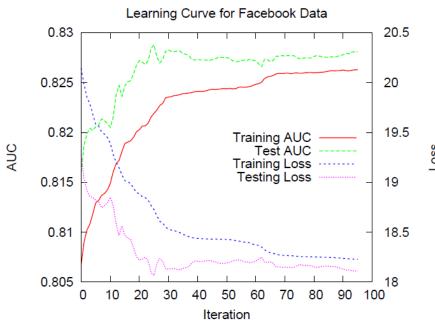
$$\frac{\partial p_u}{\partial \beta} = \sum_{j} Q_{ju} \frac{\partial p_j}{\partial \beta} + p_j \frac{\partial Q_{ju}}{\partial \beta}$$

Looks like the PageRank equation!

$$Q'_{uv} = \begin{cases} \frac{a_{uv}}{\sum_{w} a_{uw}} & \text{if } (u, v) \in E, \\ 0 & \text{otherwise} \end{cases}$$

# **Optimizing F**

- To optimize F, use gradient based method:
  - Pick a random starting point  $\beta_0$
  - Compute the personalized PageRank vector p
  - Compute the gradient with respect to the weight vector  $\beta$
  - Update β
    - Optimize using quasi-Newton method



#### Data: Facebook

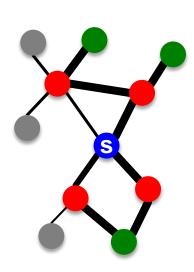
#### Facebook Iceland network

- 174,000 nodes (55% of population)
- Avg. degree 168
- Avg. person added 26 friends/month





- D={ new friendships of s created in Nov '09 }
- Negative examples:
  - L={ other nodes s did not create new links to }
- Limit to friends of friends:
  - on avg. there are 20k FoFs (max 2M)!



### Experimental setting

- Node and Edge features for learning:
  - Node: Age, Gender, Degree
  - Edge: Age of an edge, Communication, Profile visits, Co-tagged photos
- Baselines:
  - Decision trees and logistic regression:
    - Above features + 10 network features (PageRank, common friends)
- Evaluation:
  - AUC and Precision at Top20

#### Results: Facebook Iceland

- Facebook: predict future friends
  - Adamic-Adar already works great
  - Logistic regression also strong
  - SRW gives slight improvement

Learning Method	AUC	Prec@20
Random Walk with Restart	0.81725	6.80
Adamic-Adar	0.81586	7.35
Common Friends	0.80054	7.35
Degree	0.58535	3.25
DT: Node features	0.59248	2.38
DT: Network features	0.76979	5.38
DT: Node+Network	0.76217	5.86
DT: Path features	0.62836	2.46
DT: All features	0.72986	5.34
LR: Node features	0.54134	1.38
LR: Network features	0.80560	7.56
LR: Node+Network	0.80280	7.56
LR: Path features	0.51418	0.74
LR: All features	0.81681	7.52
SRW: one edge type	0.82502	6.87
SRW: multiple edge types	0.82799	7.57

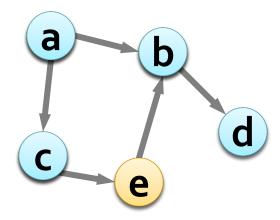
# Network Inference

# Hidden and implicit networks

- Many networks are implicit or hard to observe:
  - Hidden/hard-to-reach populations:
    - Network of needle sharing between drug injection users
  - Implicit connections:
    - Network of information propagation in online news media
- But we can observe results of the processes taking place on such (invisible) networks:
  - Virus propagation:
    - Drug users get sick, and we observe when they see the doctor
  - Information networks:
    - We observe when media sites mention information
- Question: Can we infer the hidden networks?

# Inferring the Diffusion Networks

There is a hidden diffusion network:



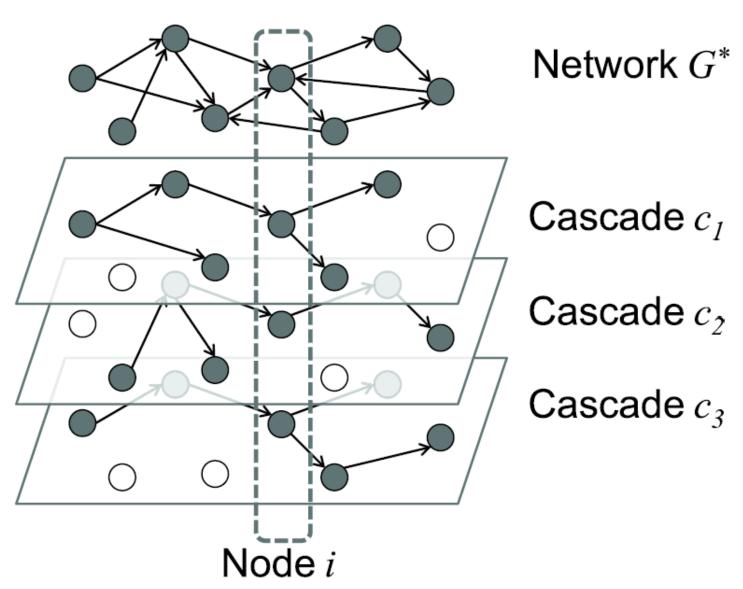
- We only see times when nodes get "infected":
  - Cascade c<sub>1</sub>: (a,1), (c,2), (b,3), (e,4)
  - Cascade c<sub>2</sub>: (c,1), (a,4), (b,5), (d,6)
- Want to infer who-infects-whom network!

# **Examples and Applications**

Word of mouth & Virus propagation **Viral marketing** Viruses propagate Recommendations and **Process** through the network influence propagate We only observe when We only observe when We observe people get sick people buy products But NOT who infected But NOT who influenced It's hidden whom whom

#### Can we infer the underlying network?

# Inferring the Diffusion Network



#### **Network Inference: The Task**

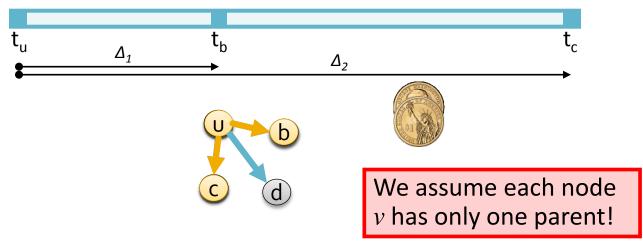
- Goal: Find a graph G that best explains the observed information times
  - Given a graph G, define the likelihood P(C|G):
    - Define a model of information diffusion over a graph
    - $P_c(u,v)$  ... prob. that *u* infects *v* in cascade *c*
    - P(c|T) ... prob. that c spread in particular pattern T
    - P(c|G) ... prob. that cascade c occurred in G
    - P(G|C) ... prob. that a set of cascades C occurred in G
- Questions:
  - How to efficiently compute P(G|C)? (given a single G)
  - How to efficiently find  $G^*$  that maximizes P(G|C)? (over  $O(2^{N*N})$  graphs)

### Cascade Diffusion Model

- Continuous time cascade diffusion model:
  - Cascade c reaches node u at  $t_u$  and spreads to u's neighbors:
    - With probability  $\beta$  cascade propagates along edge (u, v) and we determine the infection time of node v

$$t_v = t_u + \Delta$$

e.g.: △ ~ Exponential or Power-law



#### Cascade Diffusion Model

#### The model for 1 cascade:

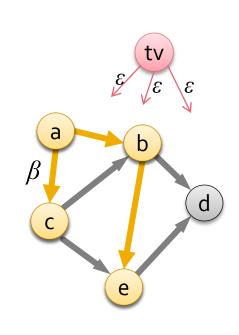
• Cascade reaches node u at time  $t_u$ , and spreads to u's neighbors v:

With prob.  $\beta$  cascade propagates along edge (u,v) and  $t_v = t_u + \Delta$ 



$$P_c(u,v) \propto P(t_v - t_u)$$
 if  $t_v > t_u$  else  $\varepsilon$  e.g.:  $P_c(u,v) \propto e^{-\Delta t}$ 

- ullet  $\varepsilon$  captures influence external to the network
  - At any time a node can get infected from outside with small probability  $\varepsilon$



### **Cascade Probability**

#### Given node infection times and pattern T:

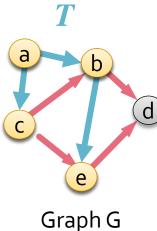
$$c = \{ (a,1), (c,2), (b,3), (e,4) \}$$

$$\blacksquare T = \{ a \rightarrow b, a \rightarrow c, b \rightarrow e \}$$

Prob. that c propagates in pattern T

$$P(c|T) = \prod_{\substack{(u,v) \in E_T \\ \text{Edges that "propagated"}}} \beta P_c(u,v) \prod_{\substack{u \in V_T, (u,x) \in E \backslash E_T \\ \text{Edges that failed to "propagate"}}} (1-\beta)$$

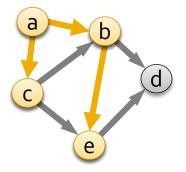
- Approximate it as:  $P(c|T) \approx \prod_{(u,v) \in E_T} P_c(v,u)$ 

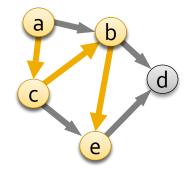


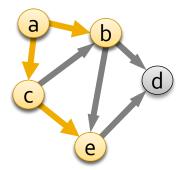
# Complication: Too Many Trees

How likely is cascade c to spread in graph G?

$$c = \{(a,1), (c,2), (b,3), (e,4)\}$$







Need to consider all possible ways for c to spread over G (i.e., all spanning trees T):

$$P(c|G) = \sum_{T \in \mathcal{T}_c(G)} P(c|T) \approx \max_{T \in \mathcal{T}_c(G)} P(c|T)$$

Consider only the most likely propagation tree

### **The Optimization Problem**

Score of a graph G for a set of cascades C:

$$P(C|G) = \prod_{c \in C} P(c|G)$$
$$F_C(G) = \sum_{c \in C} \log P(c|G)$$

Want to find the "best" graph:

$$G^* = \operatorname*{argmax}_{|G| \le k} F_C(G)$$

The problem is NP-hard: MAX-k-COVER [KDD '10]

#### How to Find the Best Tree?

Given a cascade c, what is the most likely propagation tree?

$$\max_{T \in \mathcal{T}_c(G)} P(c|T) = \max_{T \in \mathcal{T}(G)} \sum_{(i,j) \in T} w_c(i,j)$$

- A maximum directed spanning tree
  - Edge (i,j) in G has weight  $w_c(i,j) = log P_c(i,j)$
  - The maximum weight spanning tree on infected nodes: Each node picks an in-edge of

$$^{\text{max weight:}} = \sum_{i \in V} \max_{Par_T(i)} w(Par_T(i), i)$$

Local greedy selection gives optimal tree!

### Great News: Submodularity!

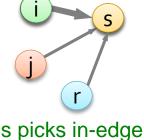
#### Theorem:

F<sub>c</sub>(G) is monotonic, and submodular

#### Proof:

- Single cascade c, some edge e=(r,s) of weight.  $w_{rs}$
- Show  $F_c(G \cup \{e\}) F_c(G) \ge F_c(G' \cup \{e\}) F_c(G')$
- Let w<sub>s</sub> be max weight in-edge of s in G
- Let w'<sub>.s</sub> be max weight in-edge of s in G'
- Since  $\mathbf{G} \subseteq \mathbf{G'}: w_{.s} \leq w'_{.s}$  and  $w_{rs} = w'_{rs}$

■ 
$$F_c(G \cup \{(r,s)\} - F_c(G))$$
  
=  $\max(w_{.s}, w_{rs}) - w_{.s}$   
≥  $\max(w'_{.s}, w_{rs}) - w'_{.s}$   
=  $F_c(G' \cup \{(r,s)\}) - F_c(G')$ 



s picks in-edge of max weight

### **NetInf: The Algorithm**

- The algorithm:
  - Use greedy hill-climbing to maximize  $F_c(G)$ :
  - Start with empty G<sub>0</sub> (G with no edges)
  - Add k edges (k is parameter)
  - At every step i add an edge to  $G_i$  that maximizes the marginal improvement

$$e_i = \underset{e \in G \setminus G_{i-1}}{\operatorname{argmax}} F_C(G_{i-1} \cup \{e\}) - F_C(G_{i-1})$$

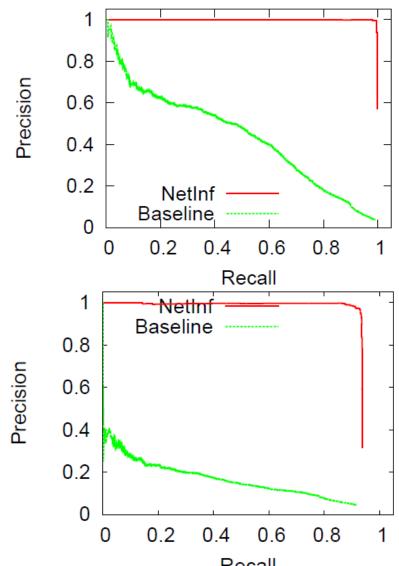
# Experiments: Synthetic data

#### Synthetic data:

- Take a graph G on k edges
- Simulate info. diffusion
- Record node infection times
- Reconstruct G

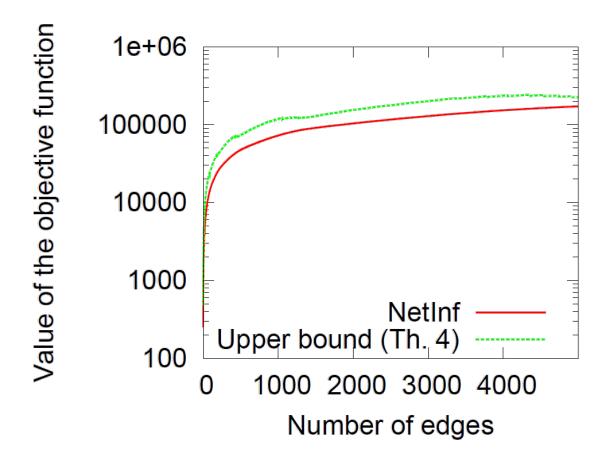
#### Evaluation:

- How many edges of G can NetInf find?
  - Break-even point: 0.95
  - Performance is independent of the structure of G!



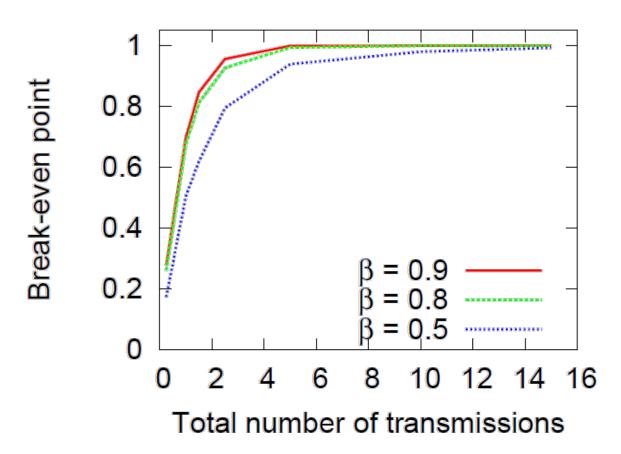
### **How Good is Our Graph?**

We achieve ≈ 90 % of the best possible network!



### How Many Cascades Do We Need?

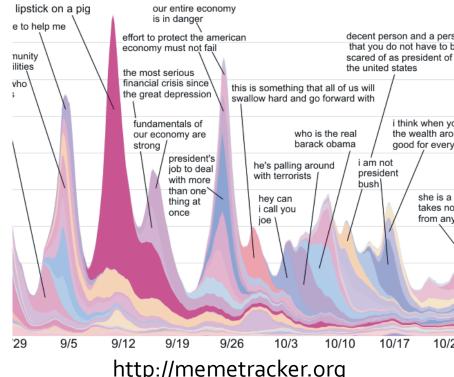
With 2x as many infections as edges, the break-even point is already 0.8 - 0.9!



### Experiments: Real data

#### Memetracker dataset:

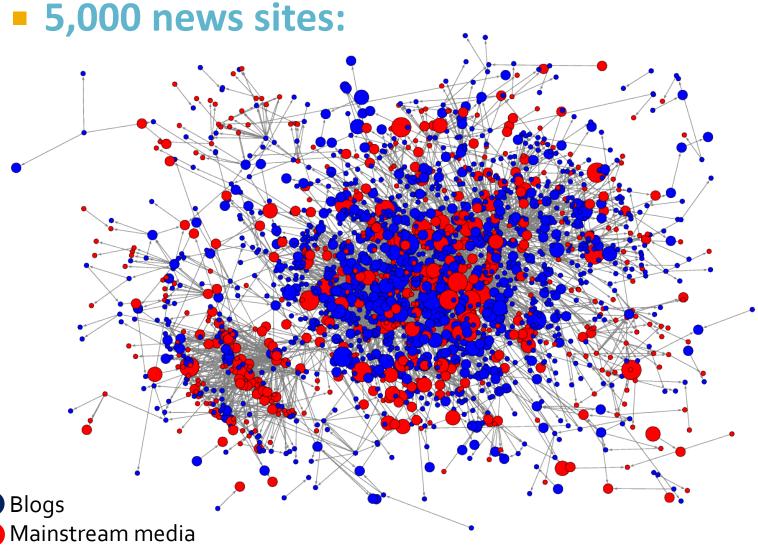
- 172m news articles
- Aug '08 Sept '09
- 343m textual phrases
- Times  $t_c(w)$  when site w mentions phrase c



http://memetracker.org

- Given times when sites mention phrases
- Infer the network of information diffusion:
  - Who tends to copy (repeat after) whom

# **Example: Diffusion Network**



# Diffusion Network (small part)

