

# Community Detection: Overlapping Communities

CS224W: Social and Information Network Analysis

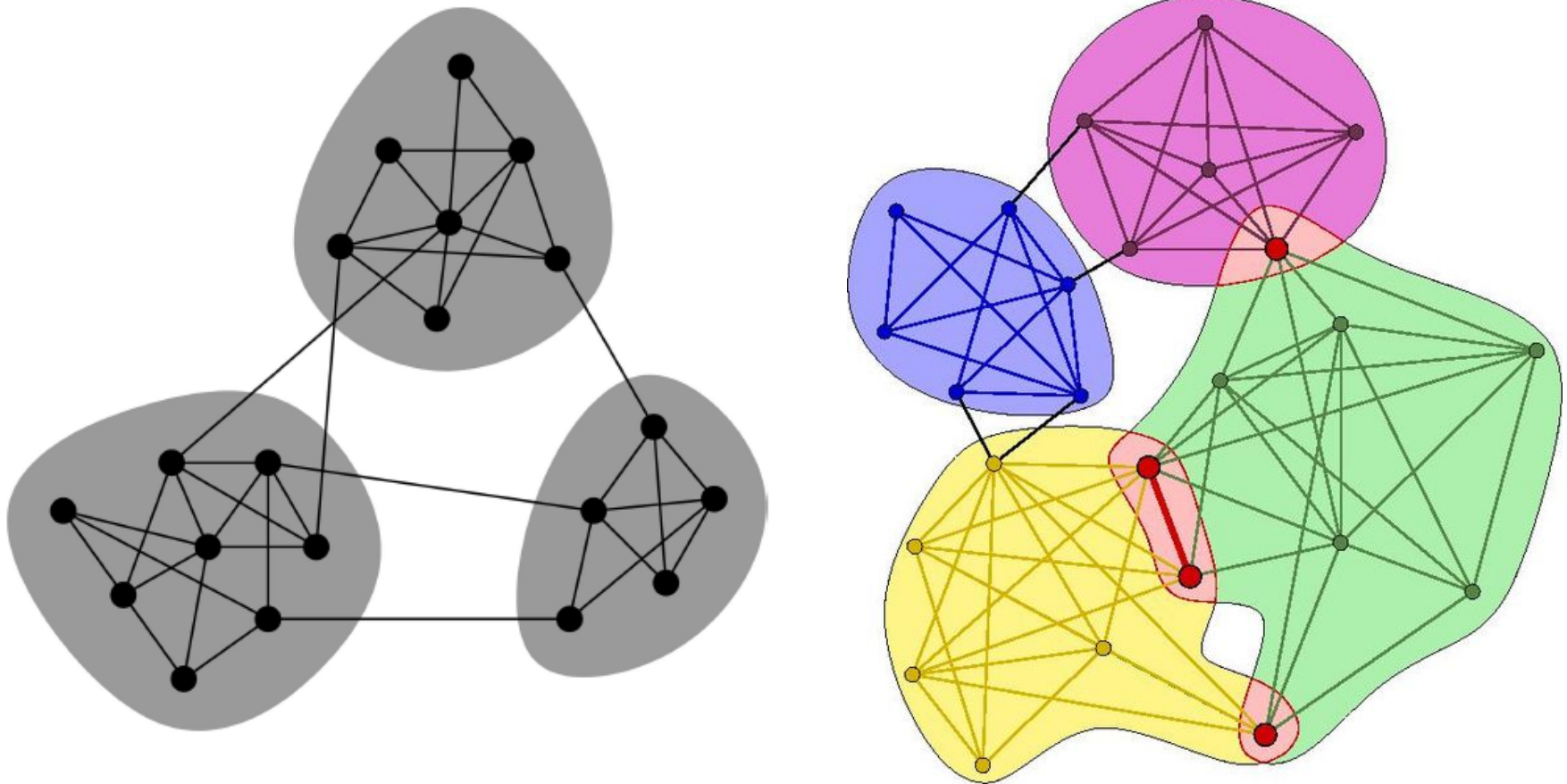
Jure Leskovec, Stanford University

<http://cs224w.stanford.edu>



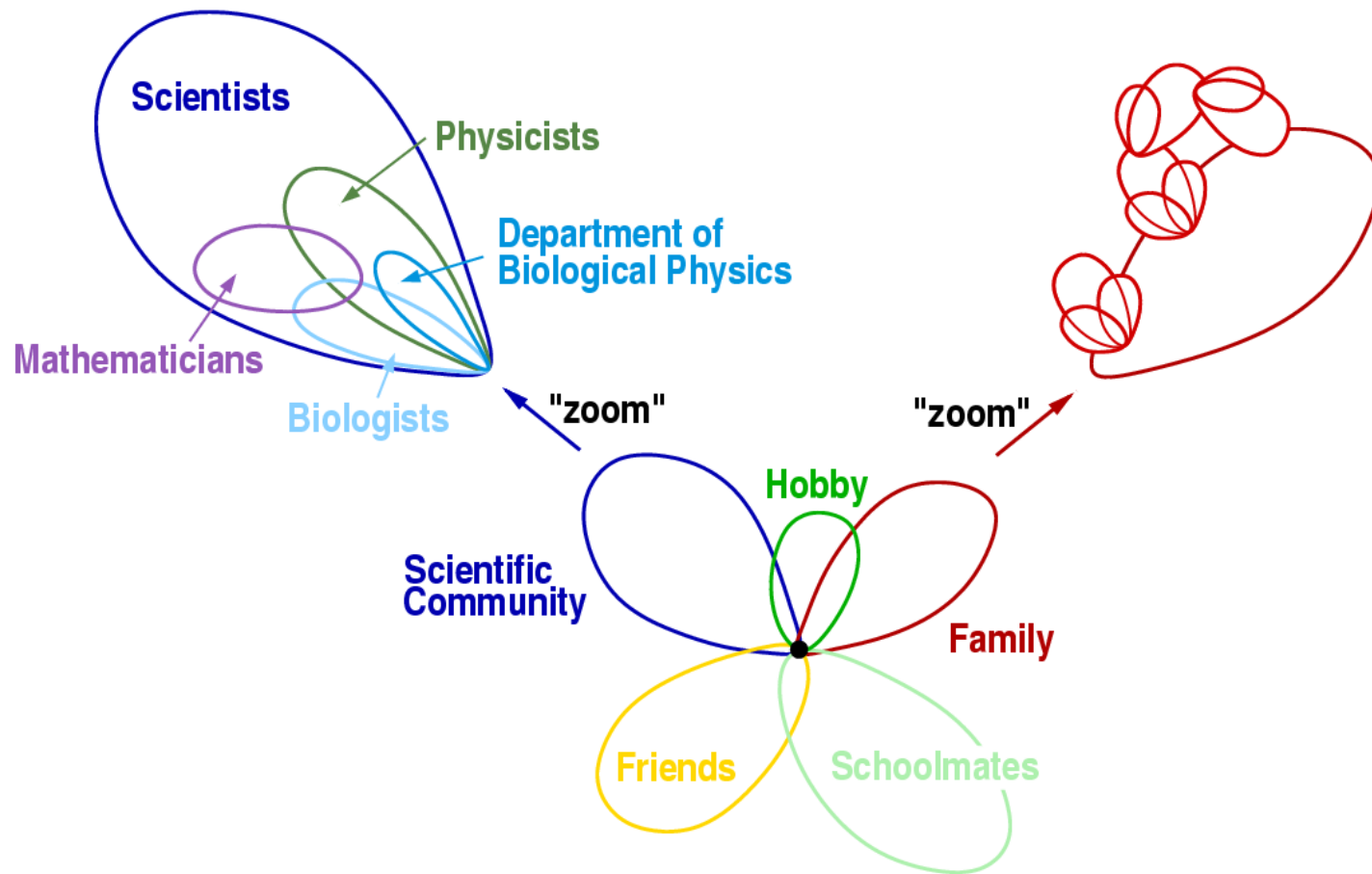
# Overlapping Communities

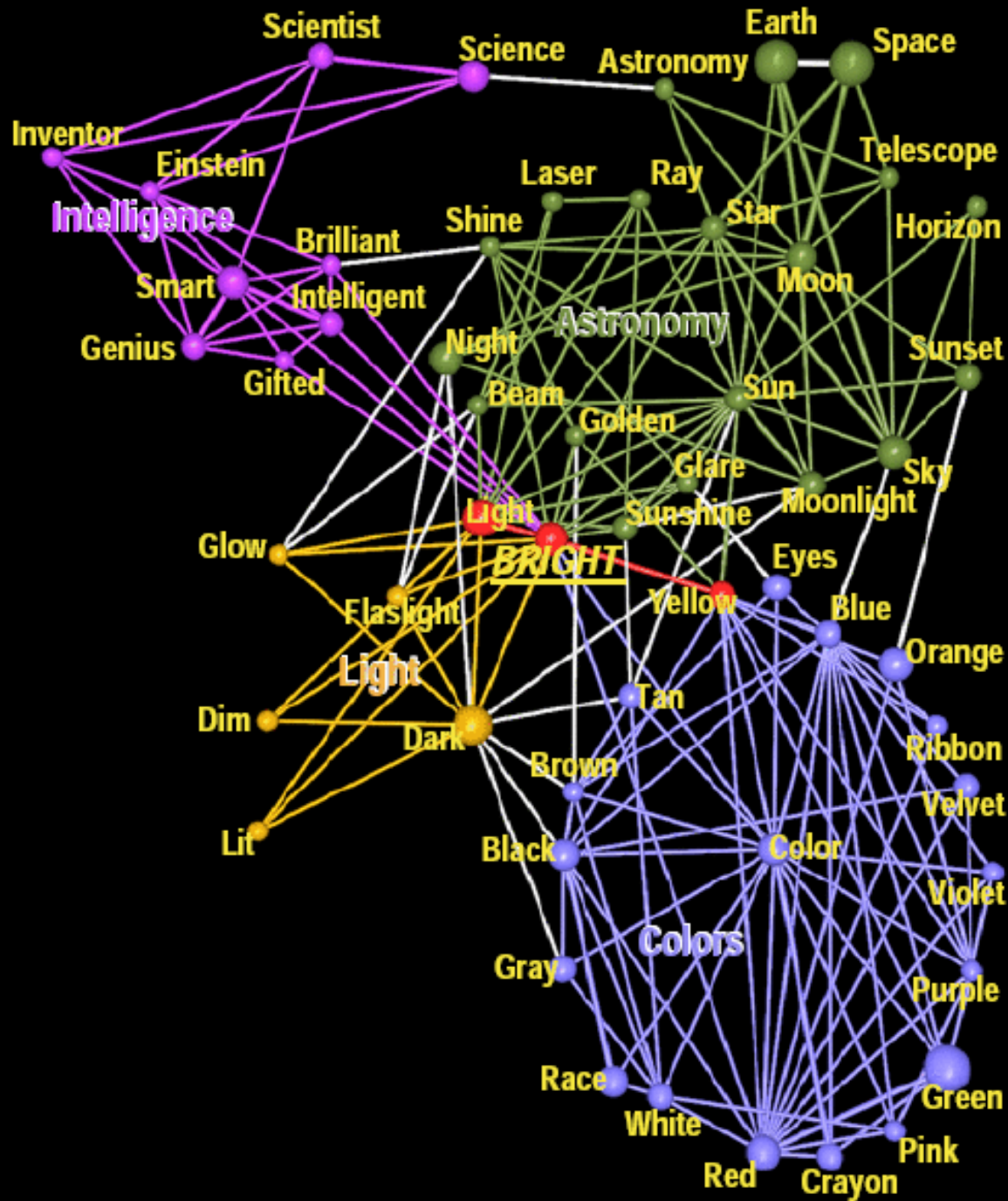
- Non-overlapping vs. overlapping communities



# Overlaps of Social Circles

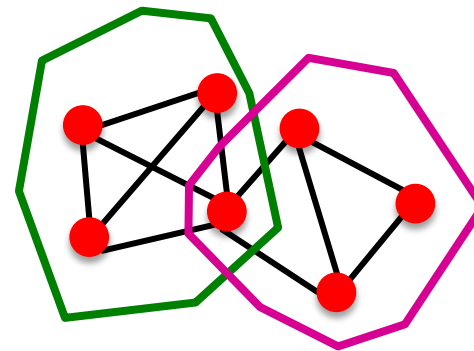
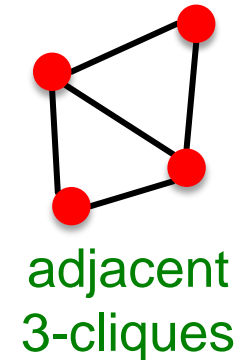
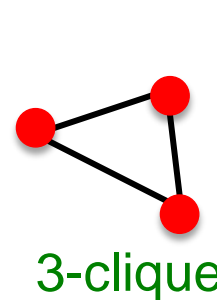
- A node belongs to many social circles





# Clique Percolation Method (CPM)

- Two nodes belong to the same community if they can be connected through adjacent  $k$ -cliques:
  - $k$ -clique:
    - Fully connected graph on  $k$  nodes
  - Adjacent  $k$ -cliques:
    - overlap in  $k-1$  nodes
- $k$ -clique community
  - Set of nodes that can be reached through a sequence of adjacent  $k$ -cliques

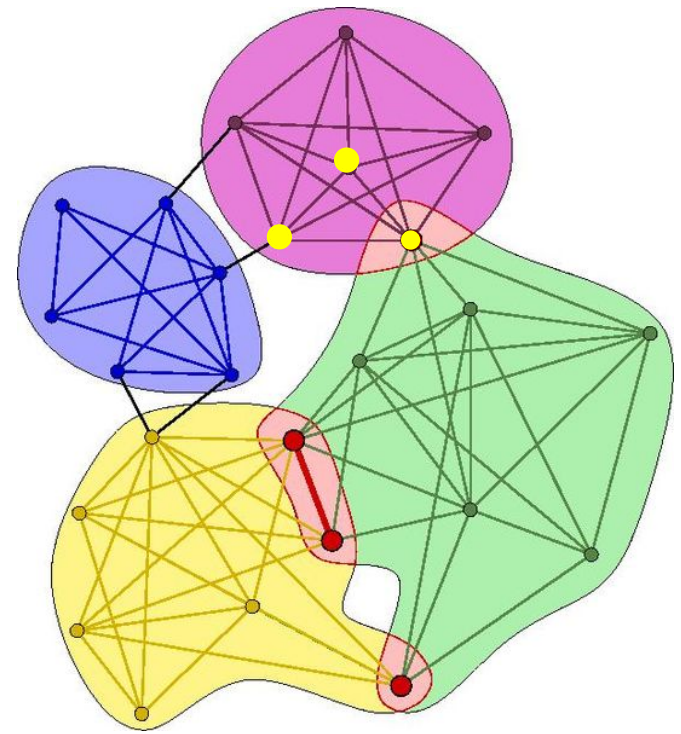




# Clique Percolation Method (CPM)

- Two nodes belong to the same community if they can be connected through adjacent  $k$ -cliques:

4-clique



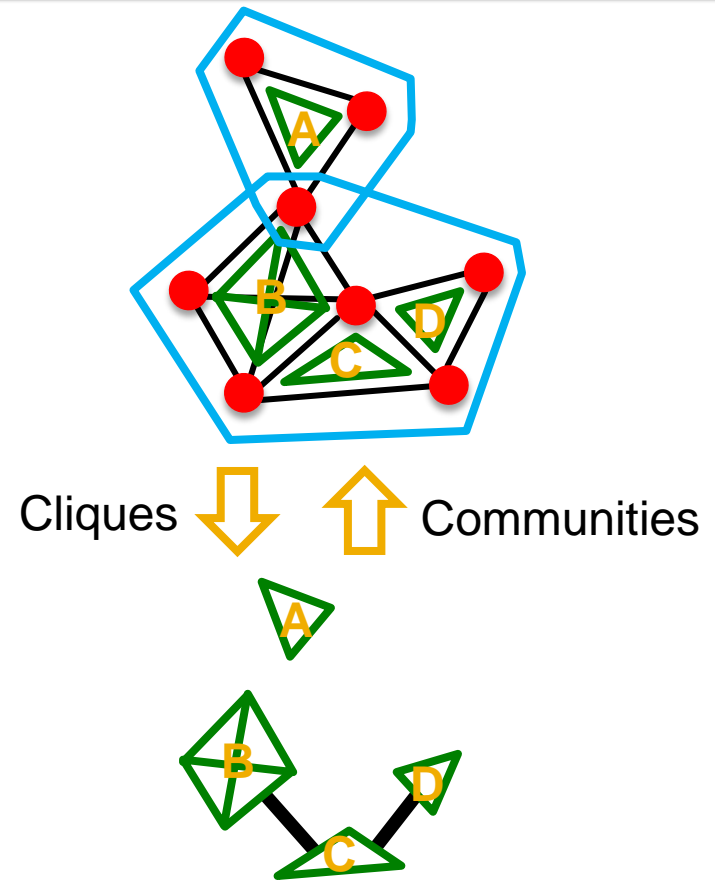
# CPM: Steps

## ■ Clique Percolation Method:

- Find maximal-cliques (not  $k$ -cliques!)
- Clique overlap graph:
  - Each clique is a node
  - Connect two cliques if they overlap in at least  $k-1$  nodes
- Communities:
  - Connected components of the clique overlap matrix

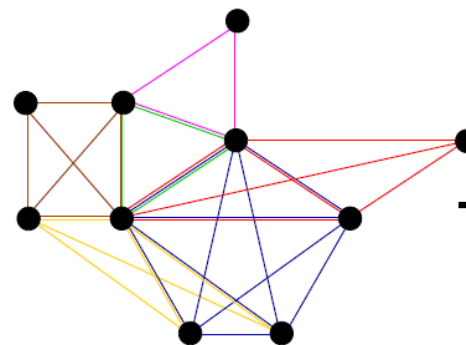
## ■ How to set $k$ ?

- Set  $k$  so that we get the “richest” (most widely distributed cluster sizes) community structure



# CPM method: Example

- Start with graph
- Find maximal cliques
- Create clique overlap matrix
- Threshold the matrix at value  $k-1$ 
  - If  $a_{ij} < k-1$  set 0
- Communities are the connected components of the thresholded matrix



(1) Graph

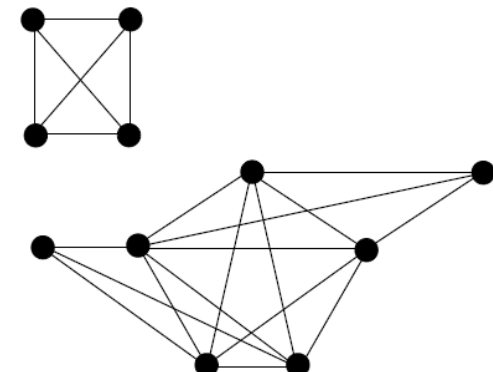
	Blue	Red	Green	Magenta	Yellow	Brown
Blue	5	3	2	1	3	1
Red	3	4	2	1	1	1
Green	2	2	3	2	1	2
Magenta	1	1	2	3	0	1
Yellow	3	1	1	0	4	2
Brown	1	1	2	1	2	4

(2) Clique overlap matrix

$k=4$

	Blue	Red	Green	Magenta	Yellow	Brown
Blue	1	1	0	0	1	0
Red	1	1	0	0	0	0
Green	0	0	0	0	0	0
Magenta	0	0	0	0	0	0
Yellow	1	0	0	0	1	0
Brown	0	0	0	0	0	1

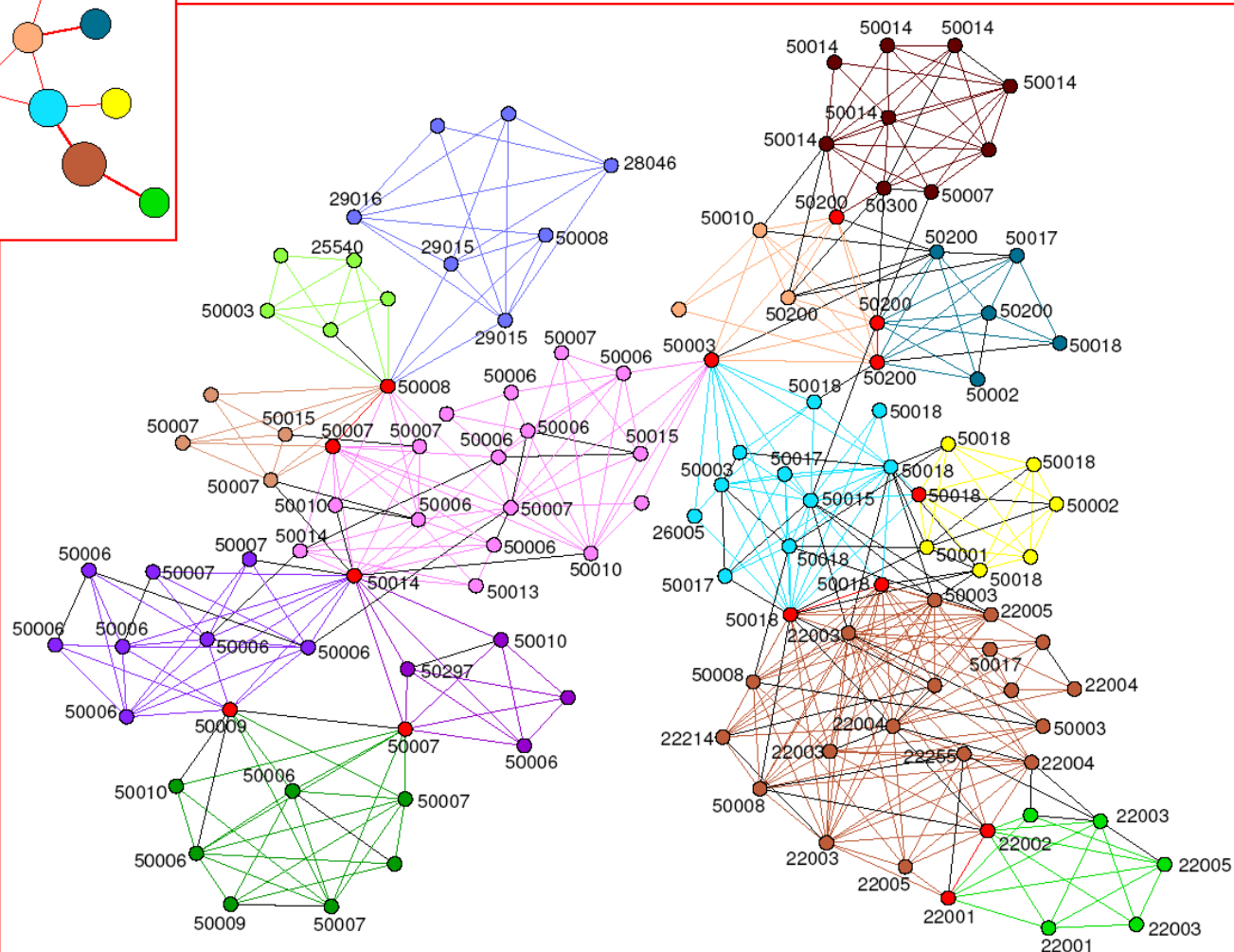
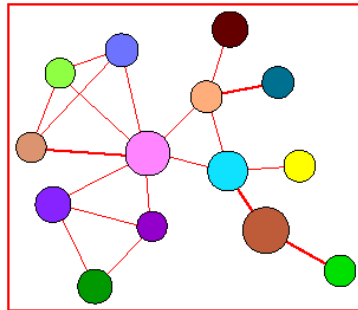
(3) Thresholded matrix at 3



(4) Communities (connected components)

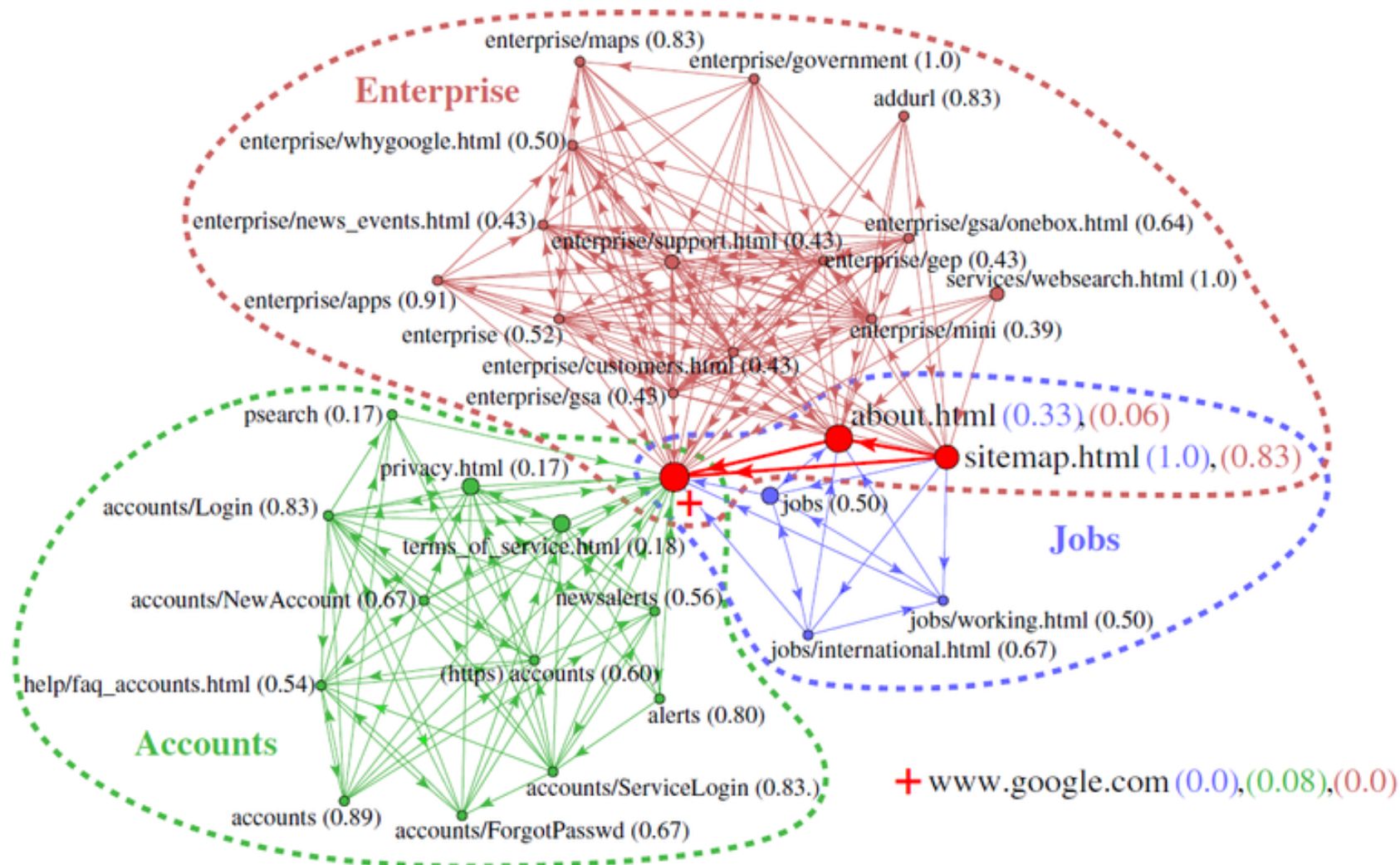


# Example: Phone-Call Network



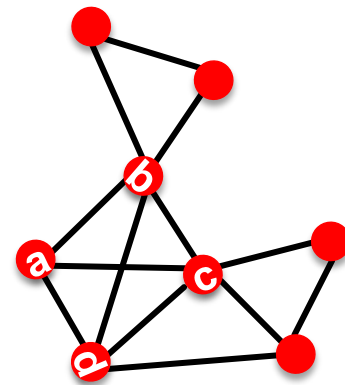
Communities in a  
“tiny” part of a phone  
call network of 4  
million users  
[Palla et al., '07]

# Example: Website



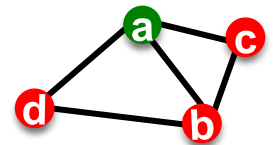
# How to Find Maximal Cliques?

- No nice way, NP-hard combinatorial problem
- **Maximal clique:** clique that can't be extended
  - $\{a,b,c\}$  is a clique but not maximal clique
  - $\{a,b,c,d\}$  is maximal clique
- **Algorithm: Sketch**
  - Start with a seed node
  - Expand the clique around the seed
  - Once the clique cannot be further expanded we found the maximal clique
  - **Note:**
    - This will generate the same clique multiple times



# How to Find Maximal Cliques?

- Start with a seed vertex “a”
- **Goal:** Find the maximal clique Q “a” belongs to
  - **Observation:**
    - If some “x” belongs to Q then it is a member of “a”
      - **Why?** If  $a, x \in Q$  but not  $a-x$ , then Q is not a clique!
- **Recursive algorithm:**
  - Q ... current clique
  - R ... candidate vertices to expand the clique to
- **Example:** Start with “a” and expand around it



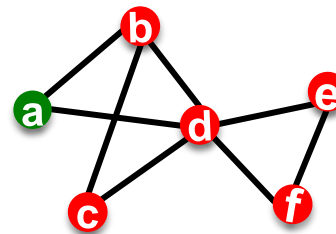
Q=	{a}	{a,b}	{a,b,c}	<b>bktrack</b>	{a,b,d}
R=	{ <u>b</u> ,c,d}	{b,c,d}	{d} ∩ Γ(c) = {}		{c} ∩ Γ(d) = {}
		∩ Γ(b) = { <u>c</u> ,d}			

Steps of the recursive algorithm

Γ(u)...neighbor set of u

# How to Find Maximal Cliques?

- Q ... current clique
- R ... candidate vertices
- **Expand(R, Q)**
  - **while**  $R \neq \{\}$ 
    - $p = \text{vertex in } R$
    - $Q_p = Q \cup \{p\}$
    - $R_p = R \cap \Gamma(p)$
    - **if**  $R_p \neq \{\}$ : **Expand**( $R_p, Q_p$ )
    - **else**: **output**  $Q_p$
    - $R = R - \{p\}$



**Start: Expand(V, {})**

$R = \{a, \dots, f\}, Q = \{\}$

$p = \{a\}$

$Q_p = \{a\}$

$R_p = \{b, d\}$

**Expand**( $R_p, Q$ ):

$R = \{b, d\}, Q = \{a\}$

$p = \{b\}$

$Q_p = \{a, b\}$

$R_p = \{d\}$

**Expand**( $R_p, Q$ ):

$R = \{d\}, Q = \{a, b\}$

$p = \{d\}$

$Q_p = \{a, b, d\}$

$R_p = \{\}$  : **output**  $\{a, b, d\}$

$p = \{d\}$

$Q_p = \{a, d\}$

$R_p = \{b\}$

**Expand**( $R_p, Q$ ):

$R = \{b\}, Q = \{a, d\}$

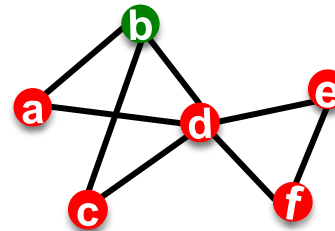
$p = \{b\}$

$Q_p = \{a, d, b\}$

$R_p = \{\}$  : **output**  $\{a, d, b\}$

# How to Find Maximal Cliques?

- Q ... current clique
- R ... candidate vertices
- **Expand(R, Q)**
  - **while**  $R \neq \{\}$ 
    - $p = \text{vertex in } R$
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    - $R = R - \{p\}$



**Start: Expand(V, {})**

$R = \{a, \dots, f\}, Q = \{\}$

$p = \{b\}$

$Q_p = \{b\}$

$R_p = \{a, c, d\}$

**Expand( $R_p, Q$ ):**

$R = \{a, c, d\}, Q = \{b\}$

$p = \{a\}$

$Q_p = \{b, a\}$

$R_p = \{d\}$

**Expand( $R_p, Q$ ):**

$R = \{d\}, Q = \{b, a\}$

$p = \{d\}$

$Q_p = \{b, a, d\}$

$R_p = \{\}$  : **output**  $\{b, a, d\}$

$p = \{c\}$

$Q_p = \{b, c\}$

$R_p = \{d\}$

**Expand( $R_p, Q$ ):**

$R = \{d\}, Q = \{b, c\}$

$p = \{d\}$

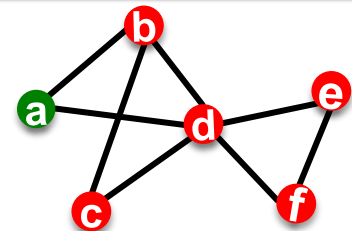
$Q_p = \{b, c, d\}$

$R_p = \{\}$  : **output**  $\{b, c, d\}$



# How to Find Maximal Cliques?

- How to prevent maximal cliques to be generated multiple times?
  - Only output cliques that are lexicographically minimum
    - $\{a,b,c\} < \{b,a,c\}$
  - **Even better:** Only expand to the nodes higher in the lexicographical order



Start:  $\text{Expand}(V, \{\})$

$R = \{a, \dots, f\}, Q = \{\}$

$p = \{a\}$

$Q_p = \{a\}$

$R_p = \{b, d\}$

**Expand( $R_p, Q$ ):**

$R = \{b, d\}, Q = \{a\}$

$p = \{b\}$

$Q_p = \{a, b\}$

$R_p = \{d\}$

**Expand( $R_p, Q$ ):**

$R = \{d\}, Q = \{a, b\}$

$p = \{d\}$

$Q_p = \{a, b, d\}$

$R_p = \{\}$  : **output  $\{a, b, d\}$**

$p = \{d\}$

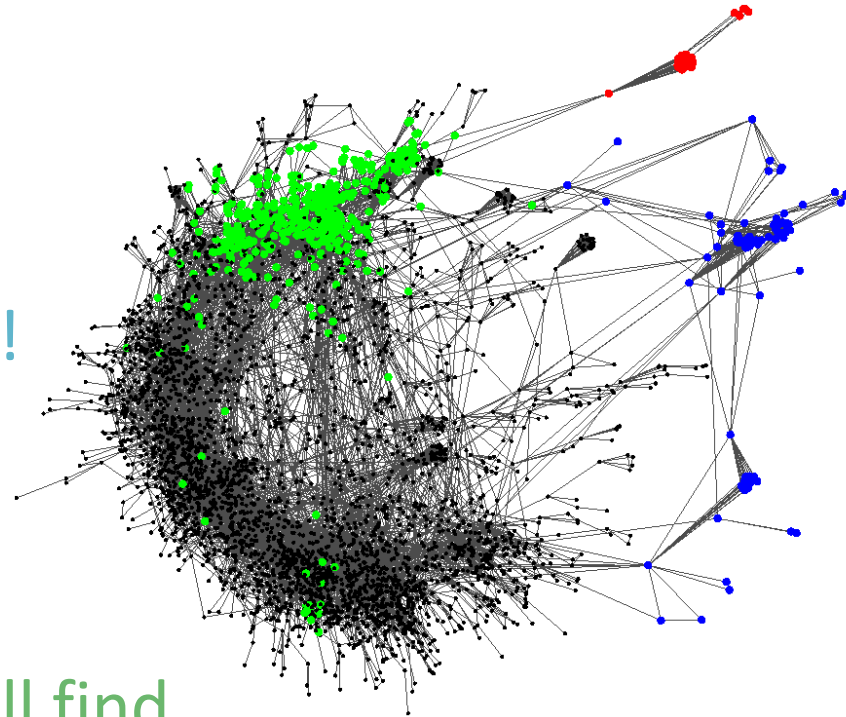
$Q_p = \{a, d\}$  **Don't expand**

$R_p = \{b\}$   **$b < d$**

# How to Model Networks with Communities?

# Reflections: Finding Communities

- Let's rethink what we are doing...
  - Given a network
  - Want to find communities!
- **Need to:**
  - Formalize the notion of a community
  - Need an algorithm that will find sets of nodes that are “good” communities
- **More generally:**
  - How to think about clusters in large networks?



# Clustering Objective Functions

What is a good cluster?

- Many edges internally
- Few pointing outside

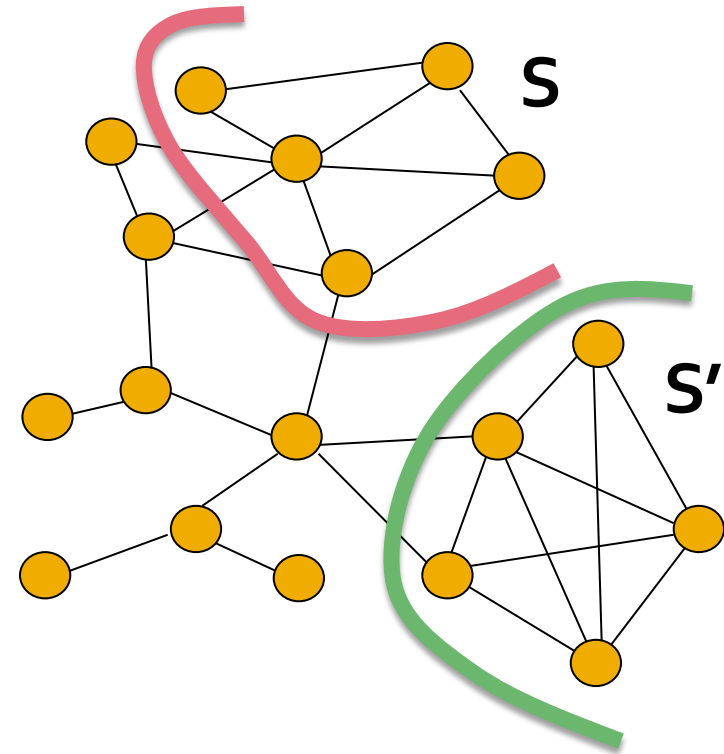
Formally, conductance:

$$\phi(S) = \frac{\sum_{i \in S, j \notin S} A_{ij}}{\min\{A(S), A(\bar{S})\}}$$

Where:  $A(S)$ ...volume

$$A(S) = \sum_{i \in S} \sum_{j \in V} A_{ij}$$

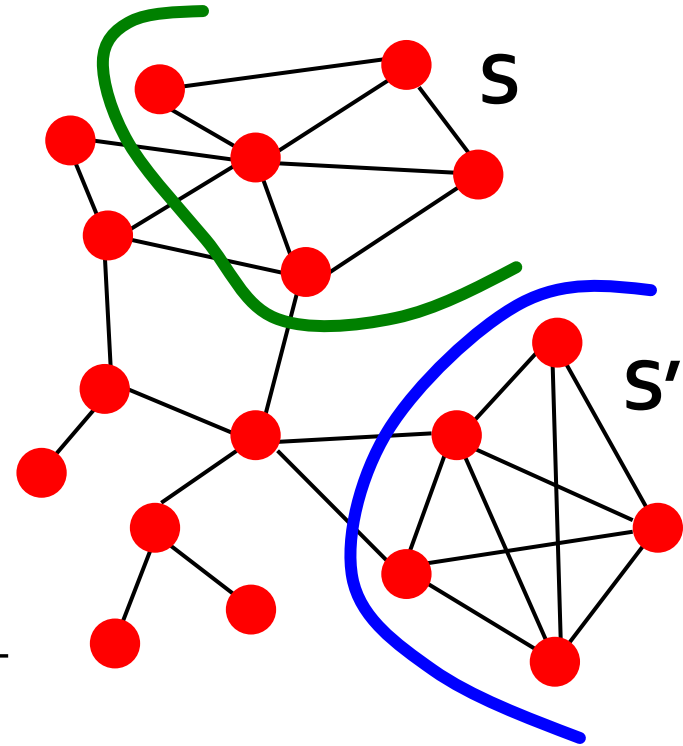
Small  $\Phi(S)$  corresponds to good clusters



# Community Score

- How community like is a set of nodes?
- A good cluster  $S$  has
  - Many edges internally
  - Few edges pointing outside
- Simplest objective function:  
**Conductance**

$$\phi(S) = \frac{|\{(i, j) \in E; i \in S, j \notin S\}|}{\sum_{s \in S} d_s}$$



Small **conductance** corresponds to good clusters

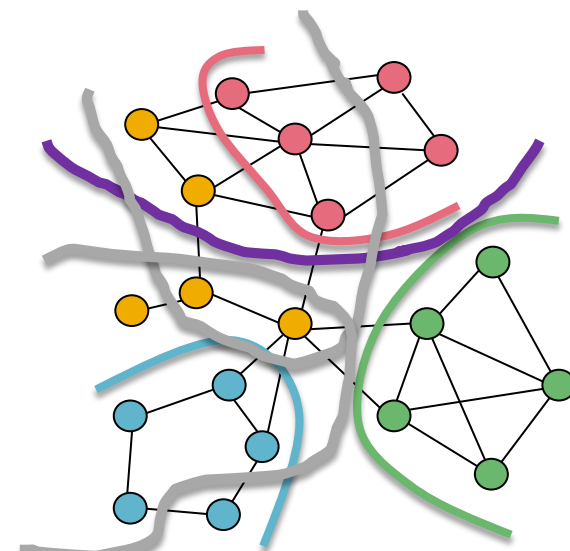
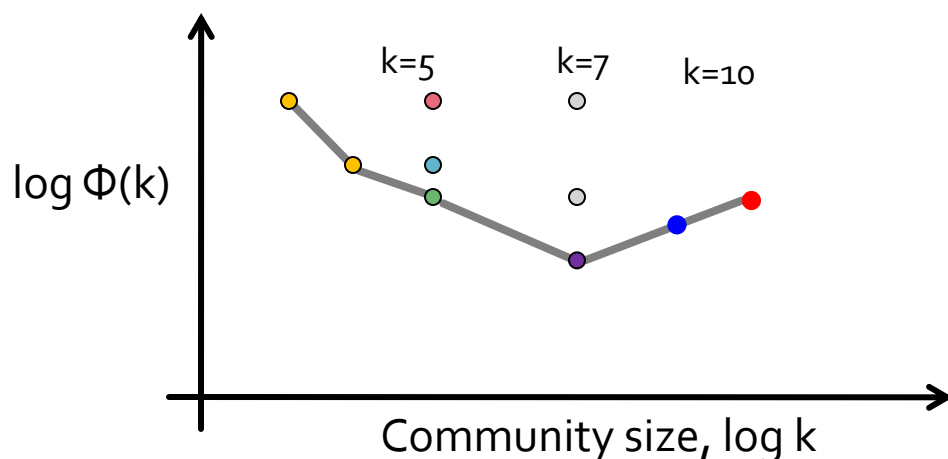
# Network Community Profile Plot

- Define:

Network community profile (**NCP**) plot

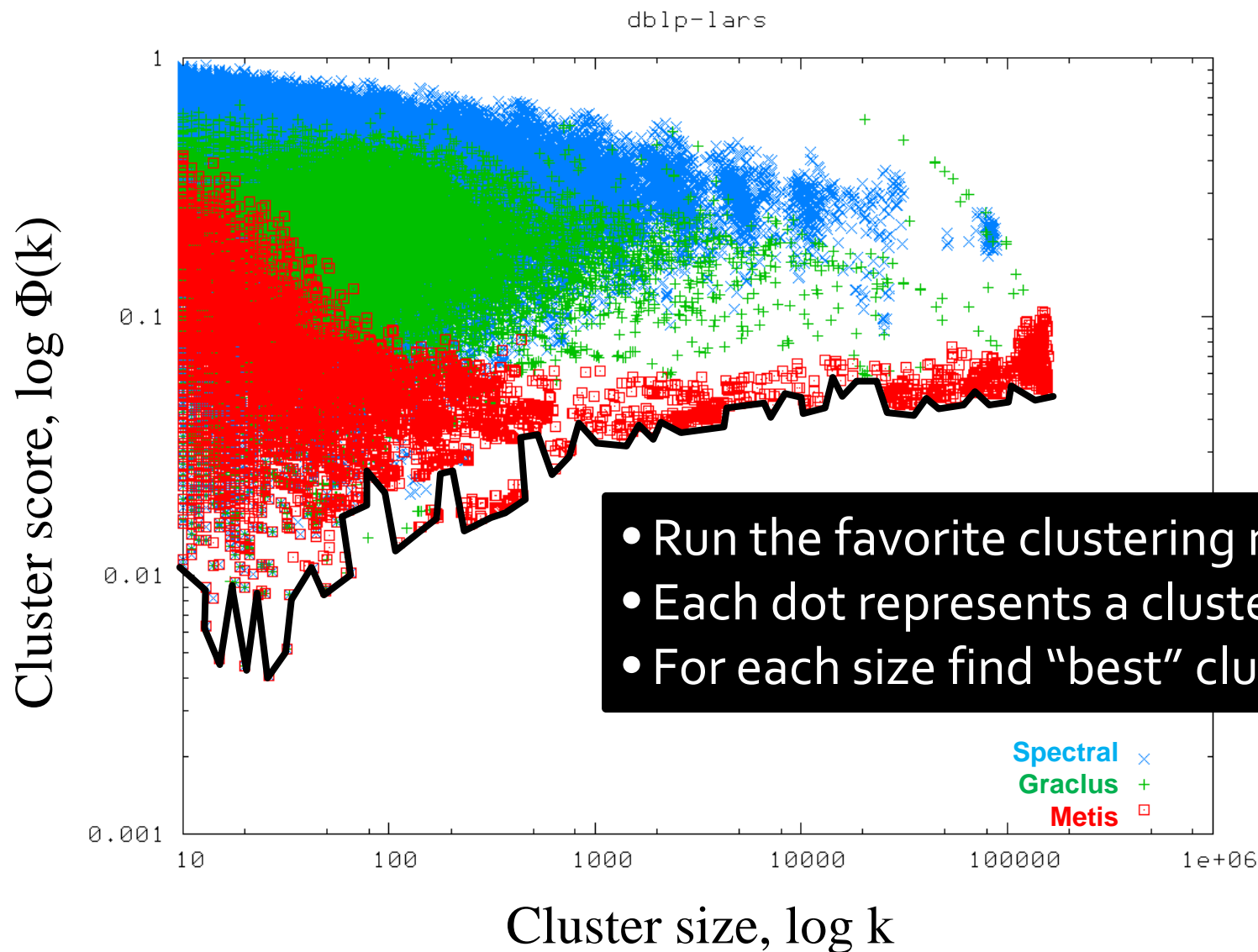
Plot the score of **best** community of size  $k$

$$\Phi(k) = \min_{S \subset V, |S|=k} \phi(S)$$



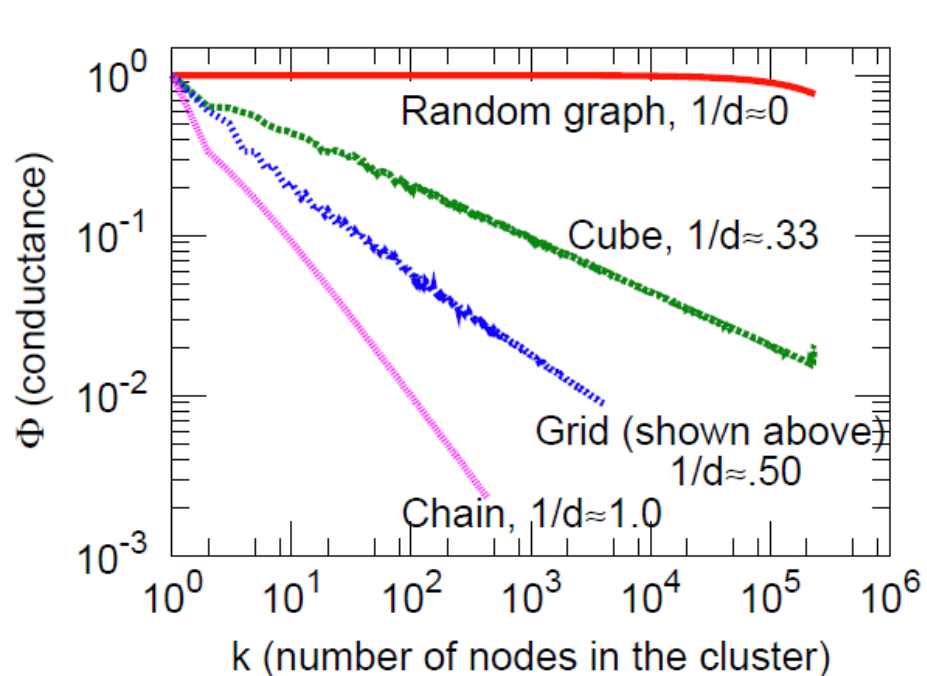


# How to (Really) Compute NCP?

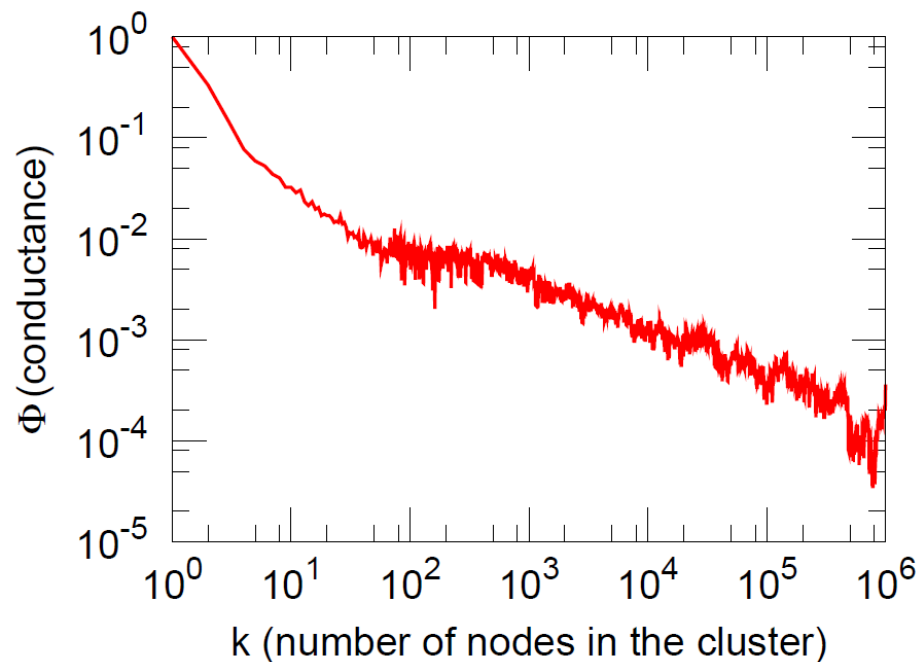


# NCP Plot: Meshes

## ■ Meshes, grids, dense random graphs:



d-dimensional meshes

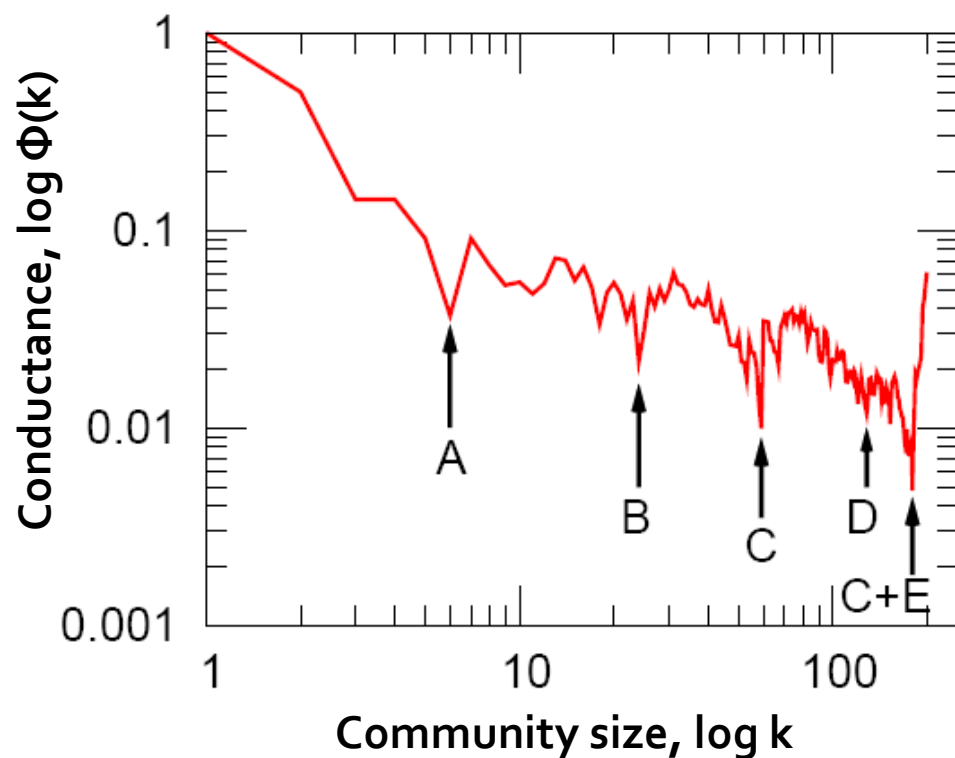
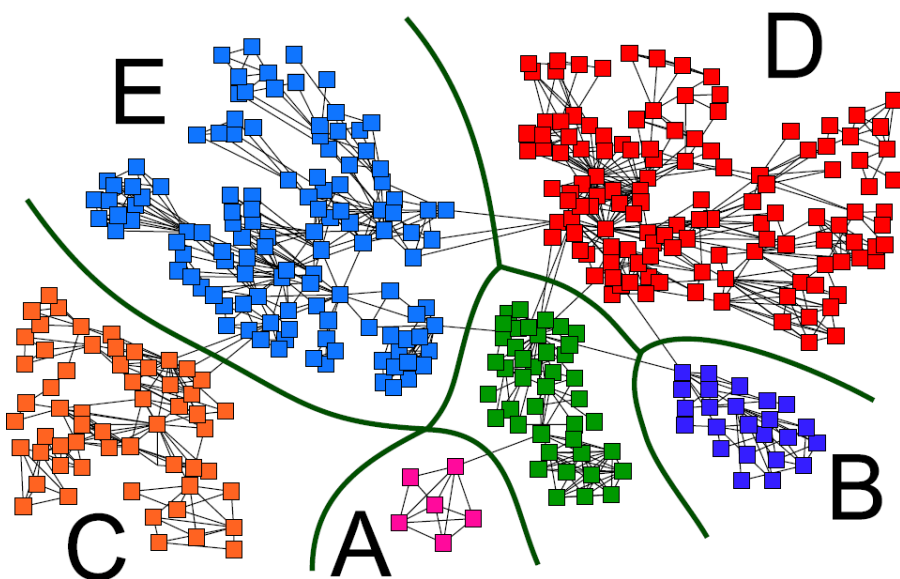


California road network

# NCP plot: Network Science

## ■ Collaborations between scientists in networks

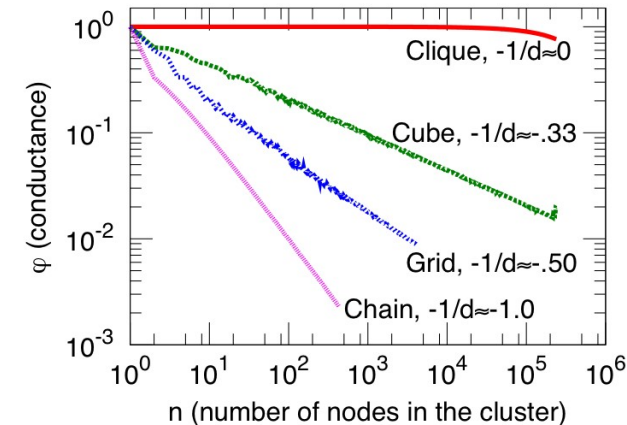
[Newman, 2005]



# Natural Hypothesis

## Natural hypothesis about NCP:

- NCP of real networks slopes downward
- Slope of the NCP corresponds to the “dimensionality” of the network

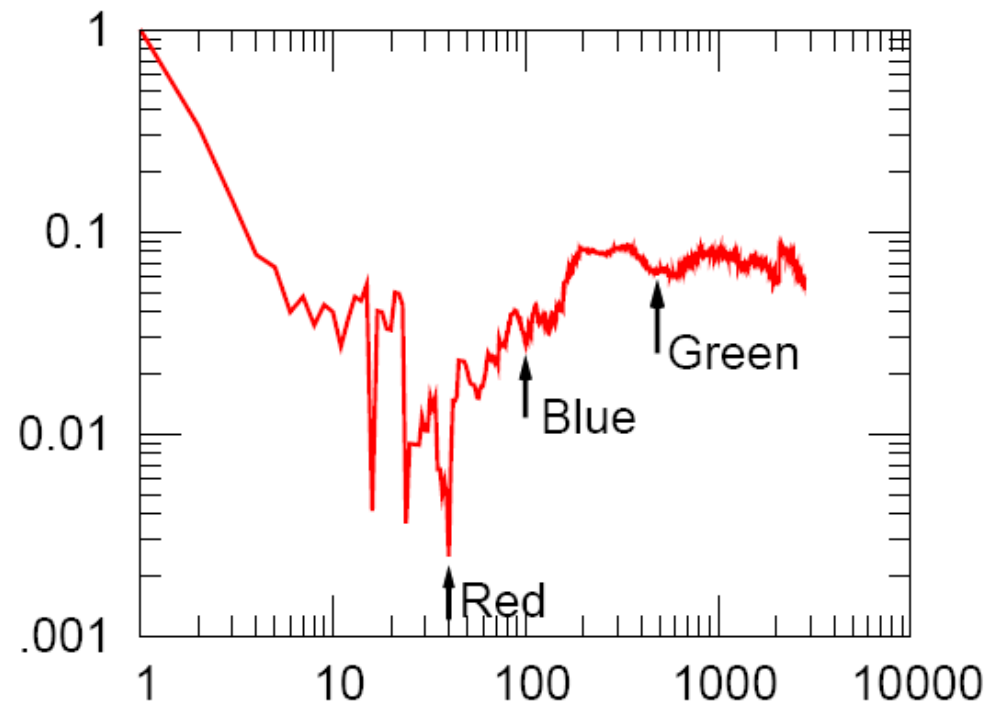
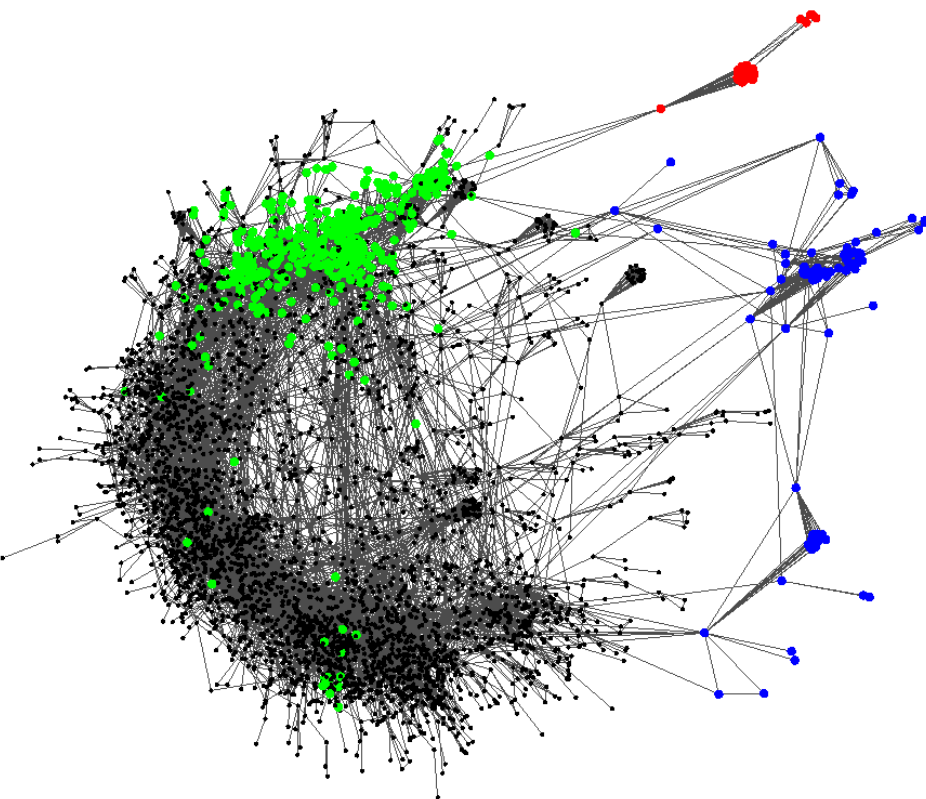


What about  
large networks?

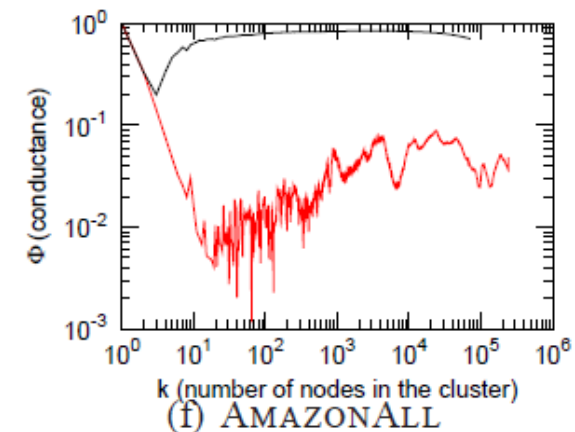
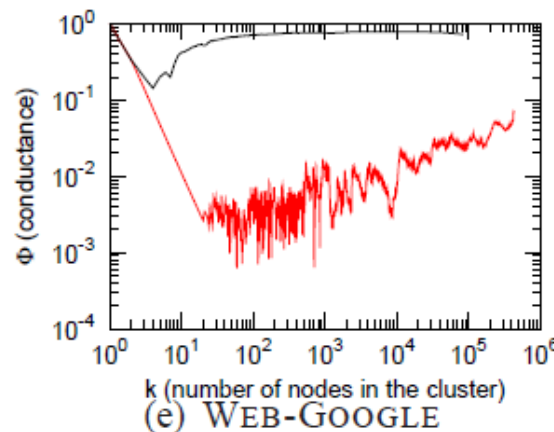
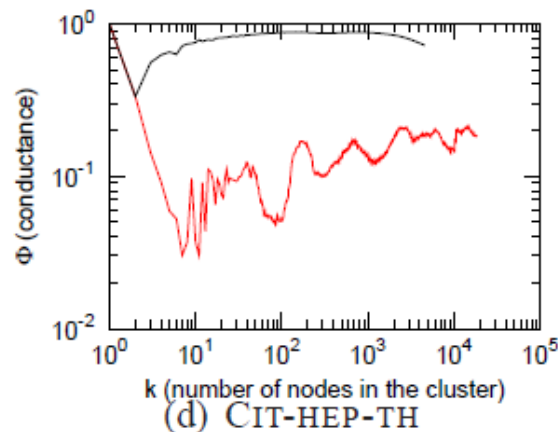
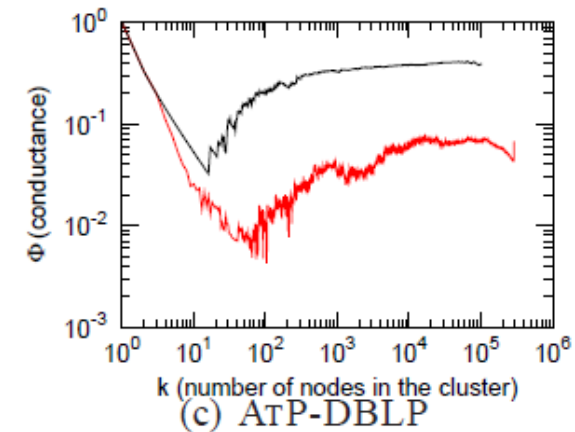
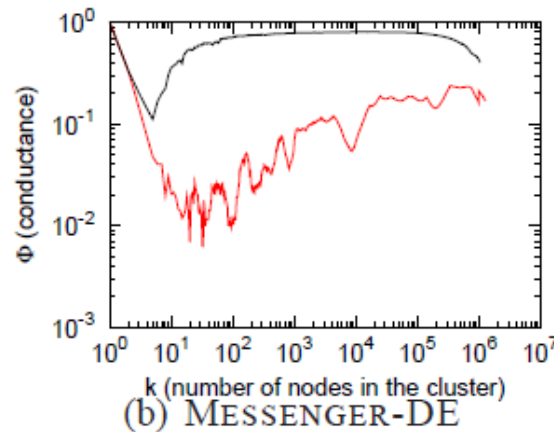
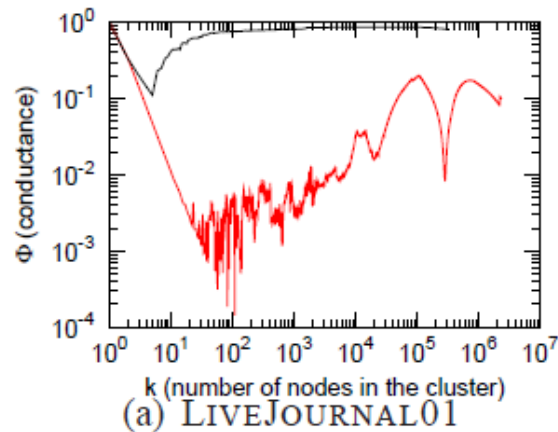
• Social nets	Nodes	Edges	Description
LIVEJOURNAL	4,843,953	42,845,684	Blog friendships [5]
EPINIONS	75,877	405,739	Trust network [28]
CA-DBLP	317,080	1,049,866	Co-authorship [5]
• Information (citation) networks			
CIT-HEP-TH	27,400	352,021	Arxiv hep-th [14]
AMAZONPROD	524,371	1,491,793	Amazon products [8]
• Web graphs			
WEB-GOOGLE	855,802	4,291,352	Google web graph
WEB-WT10G	1,458,316	6,225,033	TREC WT10G
• Bipartite affiliation (authors-to-papers) networks			
ATP-DBLP	615,678	944,456	DBLP [21]
ATM-IMDB	2,076,978	5,847,693	Actors-to-movies
• Internet networks			
ASSKITTER	1,719,037	12,814,089	Autonom. sys.
GNUTELLA	62,561	147,878	P2P network [29]

# Large Networks: Very Different

**Typical example:** General Relativity collaborations  
( $n=4,158$ ,  $m=13,422$ )

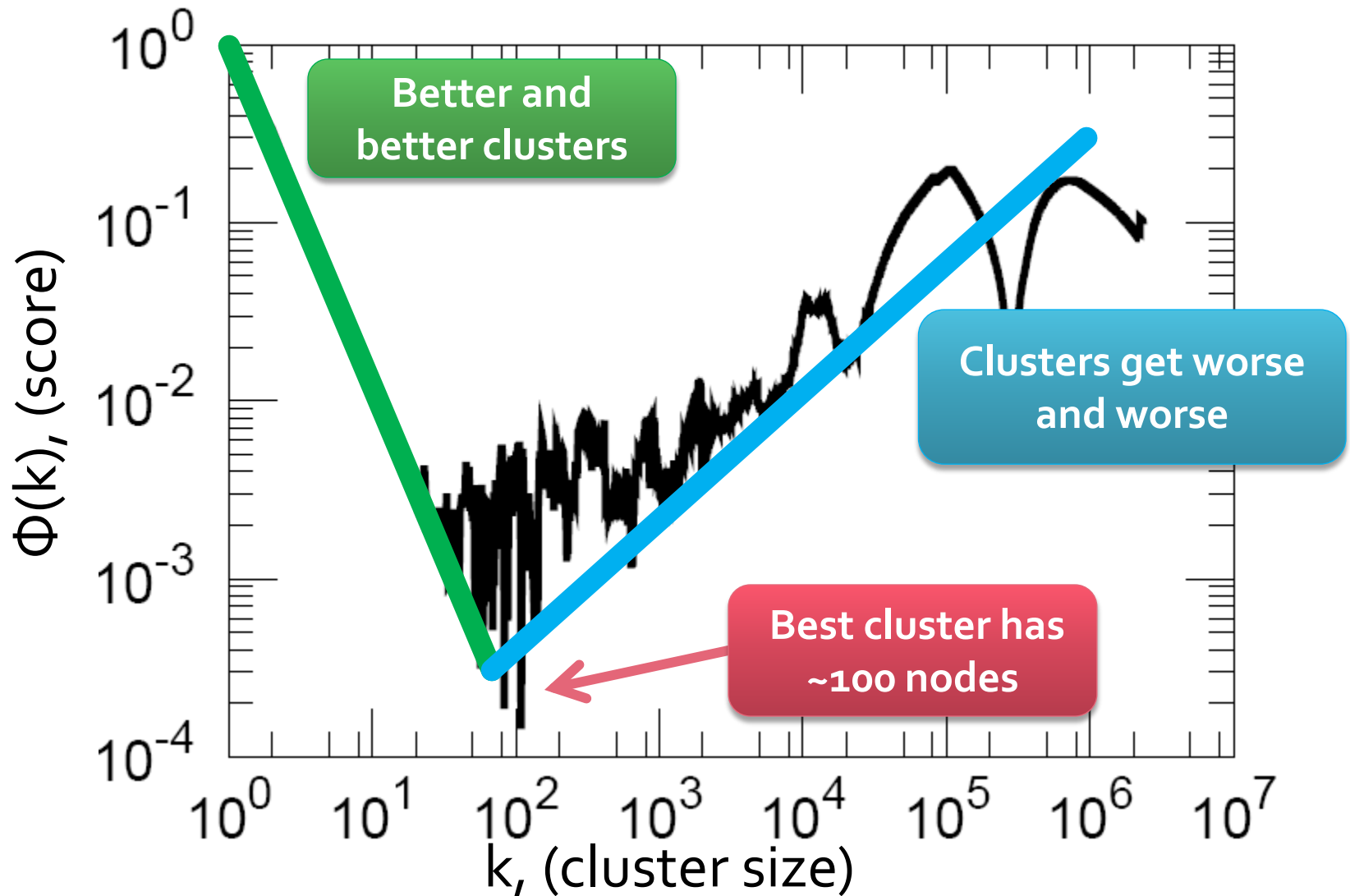


# More NCP Plots of Networks



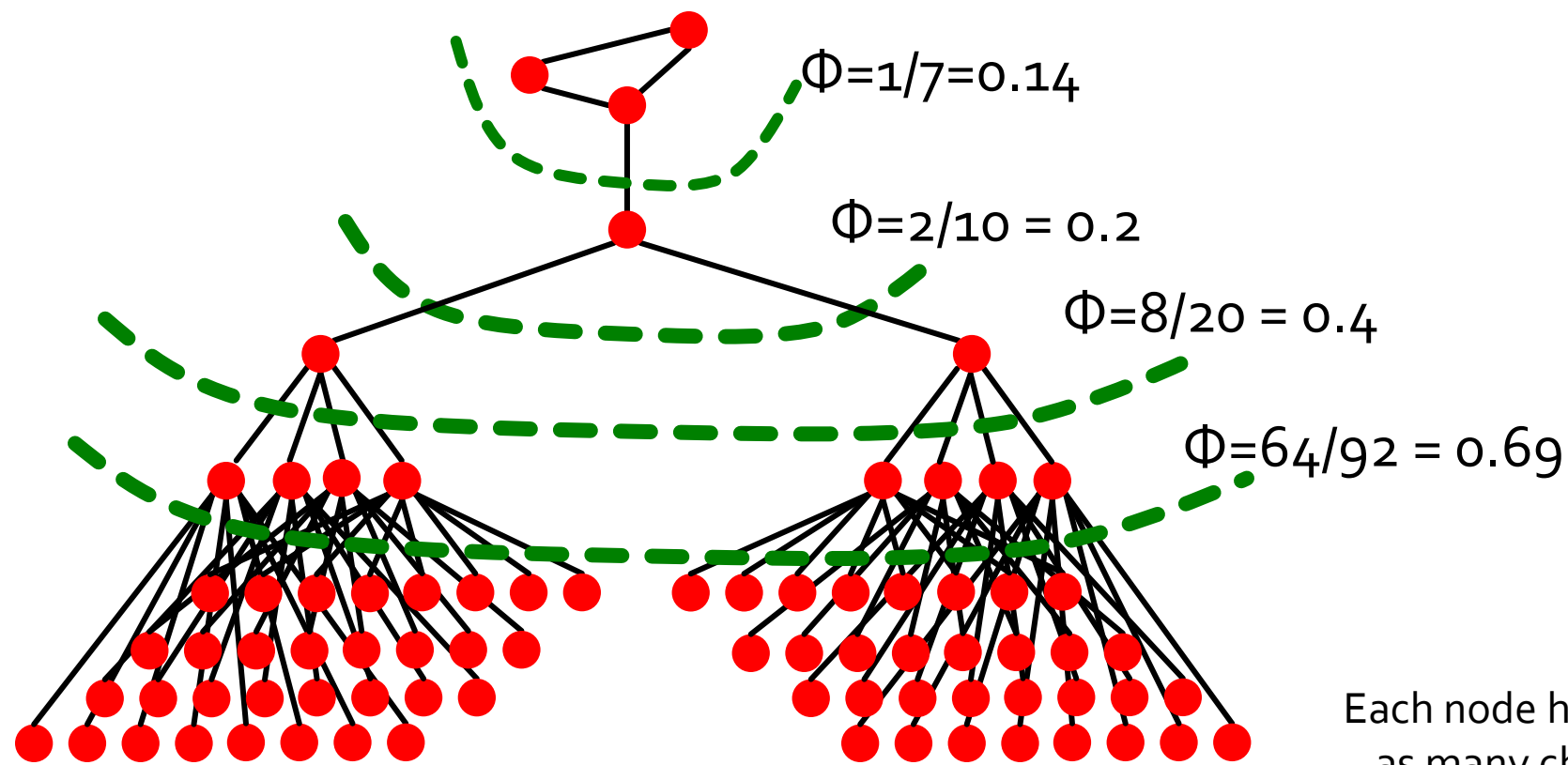


# NCP: LiveJournal (n=5m, m=42m)



# Explanation: The Upward Part

- As clusters grow the number of edges inside grows **slower** than the number crossing

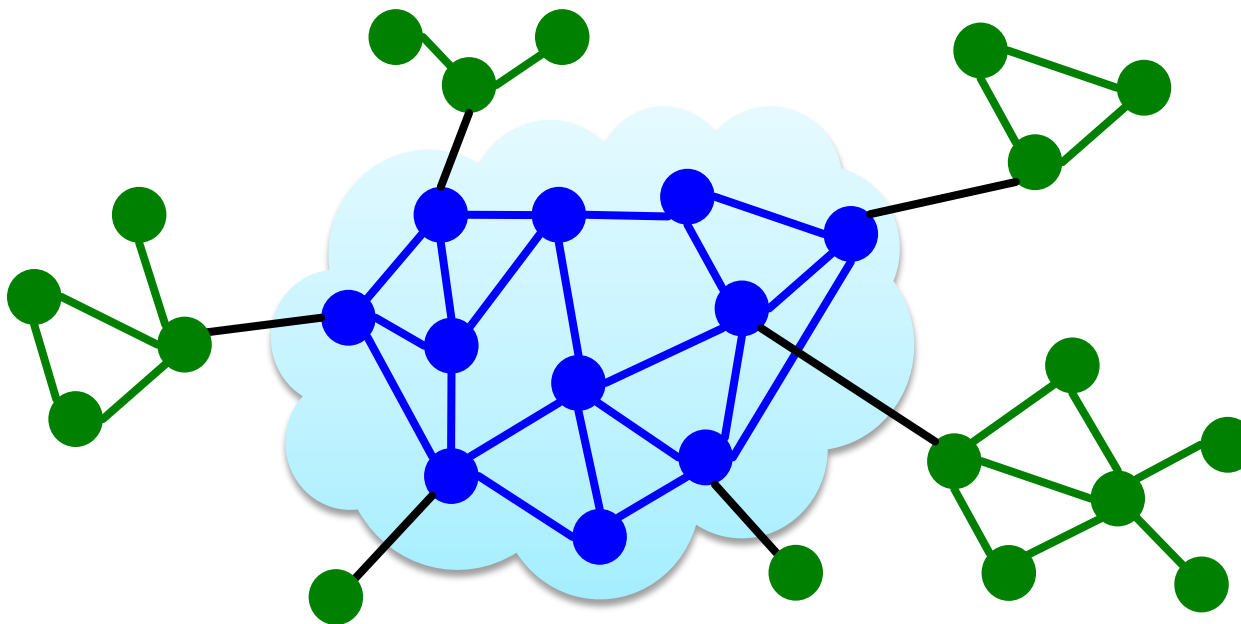


Each node has twice  
as many children

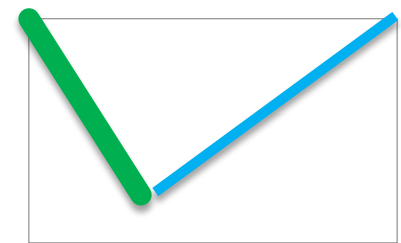
# Explanation: Downward Part



- Empirically we note that **best clusters** are **barely connected** to the network

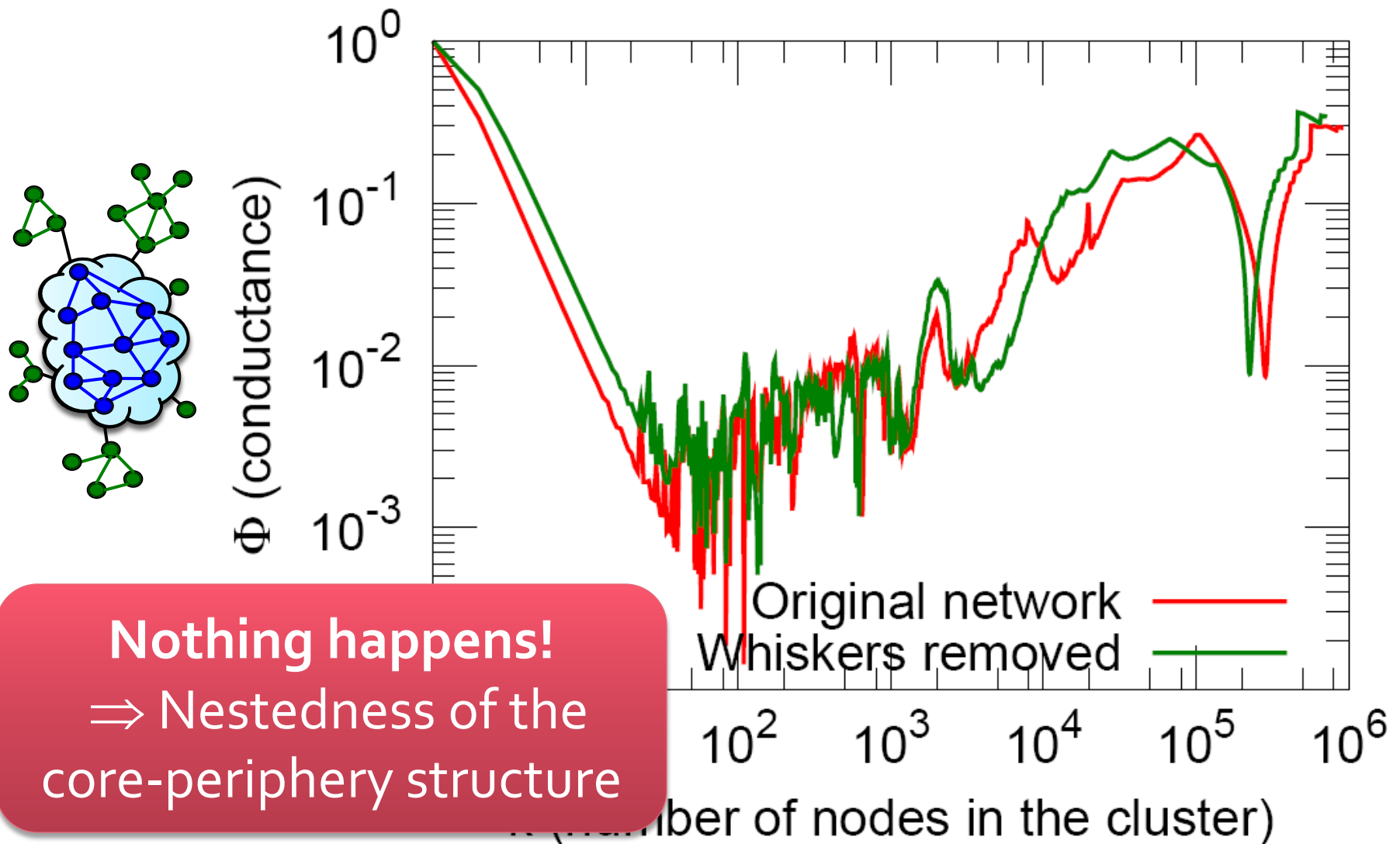


⇒ Core-periphery structure



NCP plot

# What If We Remove Good Clusters?



# Suggested Network Structure



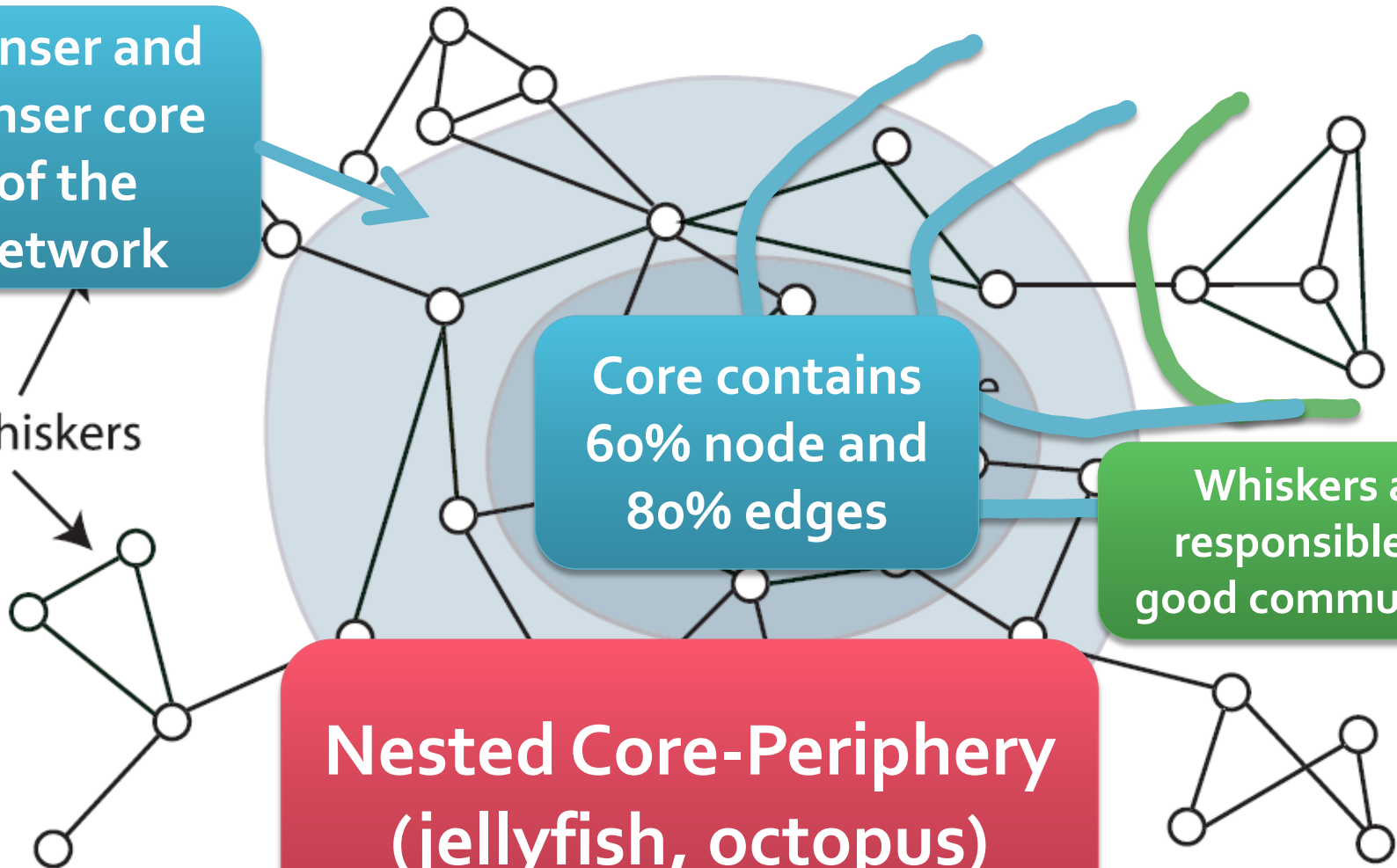
Denser and denser core of the network

Whiskers

Core contains 60% node and 80% edges

Whiskers are responsible for good communities

Nested Core-Periphery  
(jellyfish, octopus)



# Communities: Issues and Questions



# Communities: Issues and Questions

- **Some issues with community detection:**
  - Many different formalizations of clustering objective functions
  - Objectives are NP-hard to optimize exactly
  - Methods can find clusters that are systematically “biased”
- **Questions:**
  - **How well do algorithms optimize objectives?**
  - **What clusters do different methods find?**

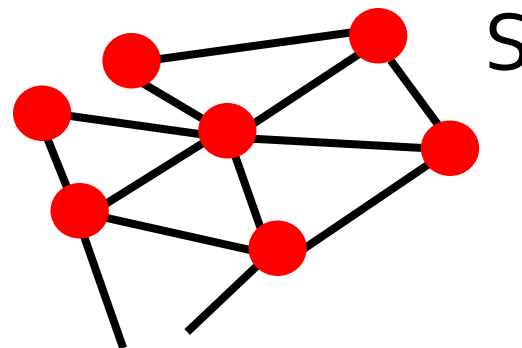
# Many Different Objective Functions

## ■ Single-criterion:

- Modularity:  $m - E(m)$
- Edges cut:  $c$

## ■ Multi-criterion:

- Conductance:  $c/(2m+c)$
- Expansion:  $c/n$
- Density:  $1 - m/n^2$
- CutRatio:  $c/n(N-n)$
- Normalized Cut:  $c/(2m+c) + c/2(M-m)+c$
- Flake-ODF: *frac. of nodes with more than  $1/2$  edges pointing outside  $S$*



$n$ : nodes in  $S$

$m$ : edges in  $S$

$c$ : edges pointing  
outside  $S$

# Many Classes of Algorithms

Many algorithms to that implicitly or explicitly optimize objectives and extract communities:

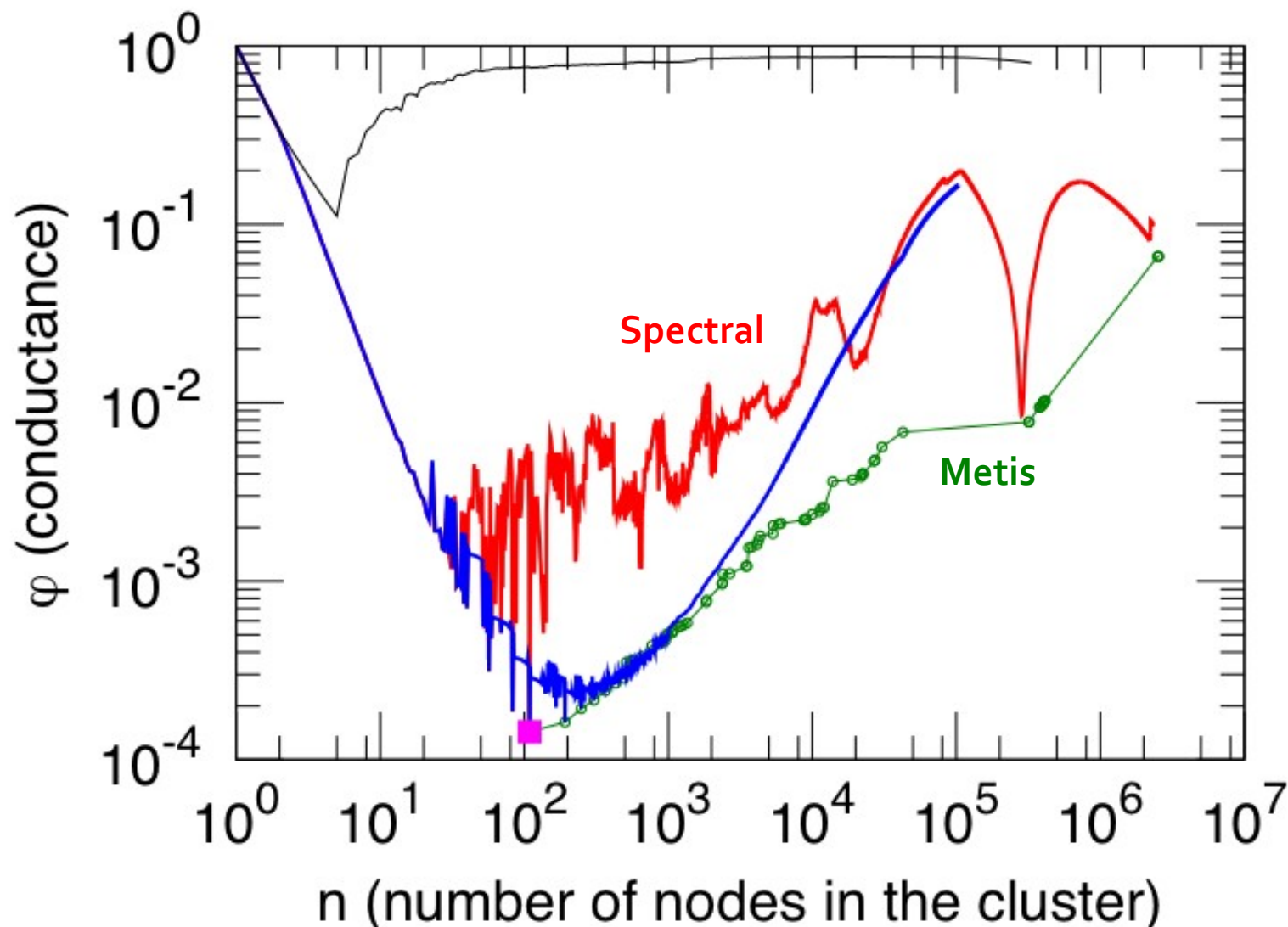
- **Heuristics:**

- Girvan-Newman, Modularity optimization: popular heuristics
- Metis: multi-resolution heuristic [Karypis-Kumar '98]

- **Theoretical approximation algorithms:**

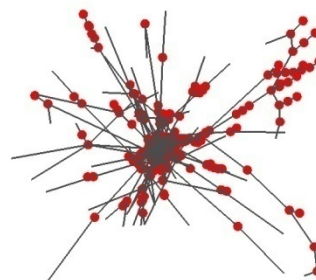
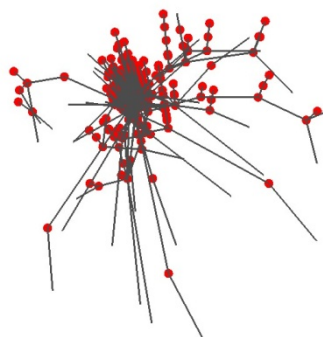
- Spectral partitioning

# NCP: Live Journal

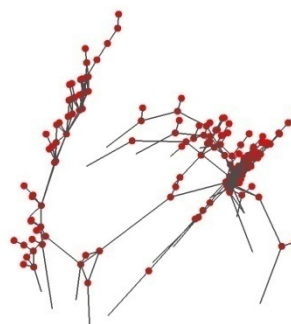
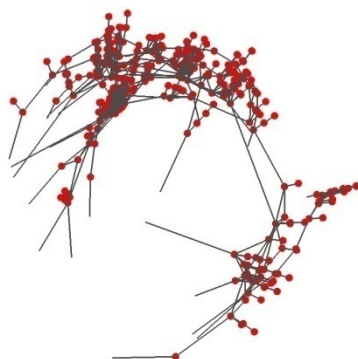


# Properties of Clusters (1)

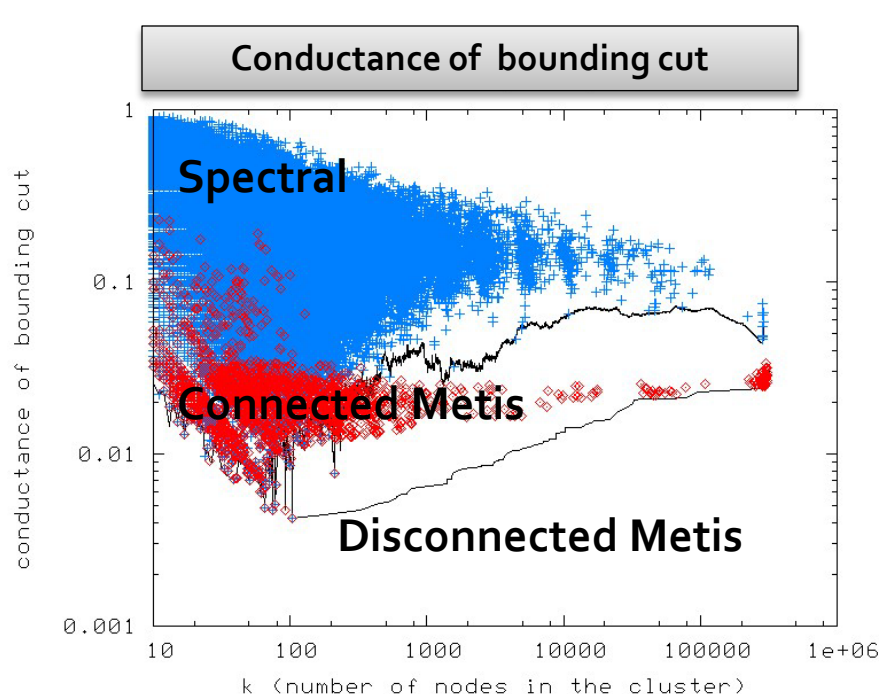
500 node communities from **Spectral**:



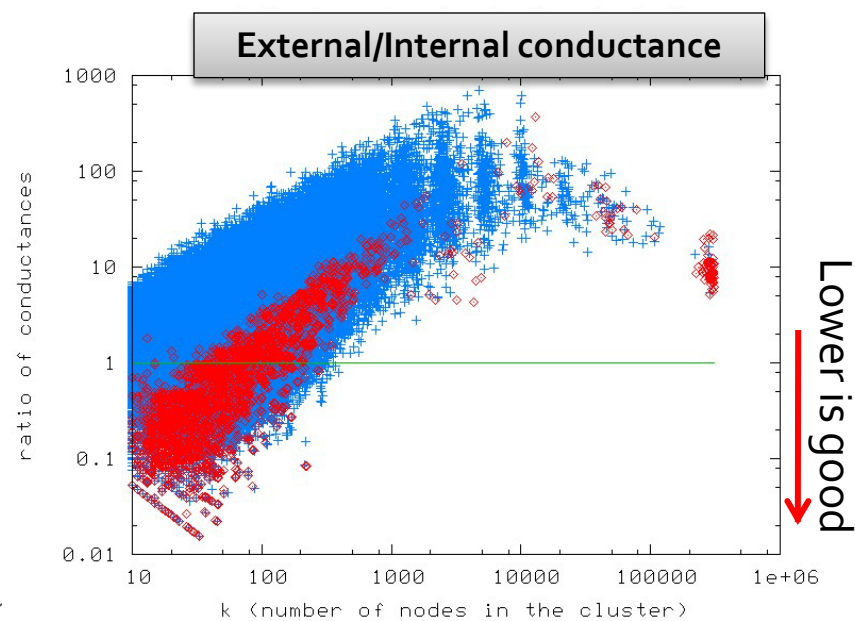
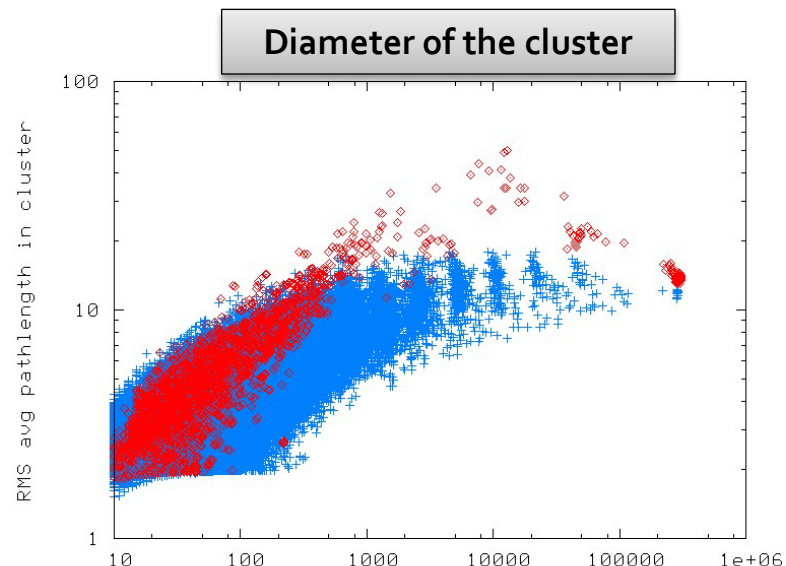
500 node communities from **Metis**:



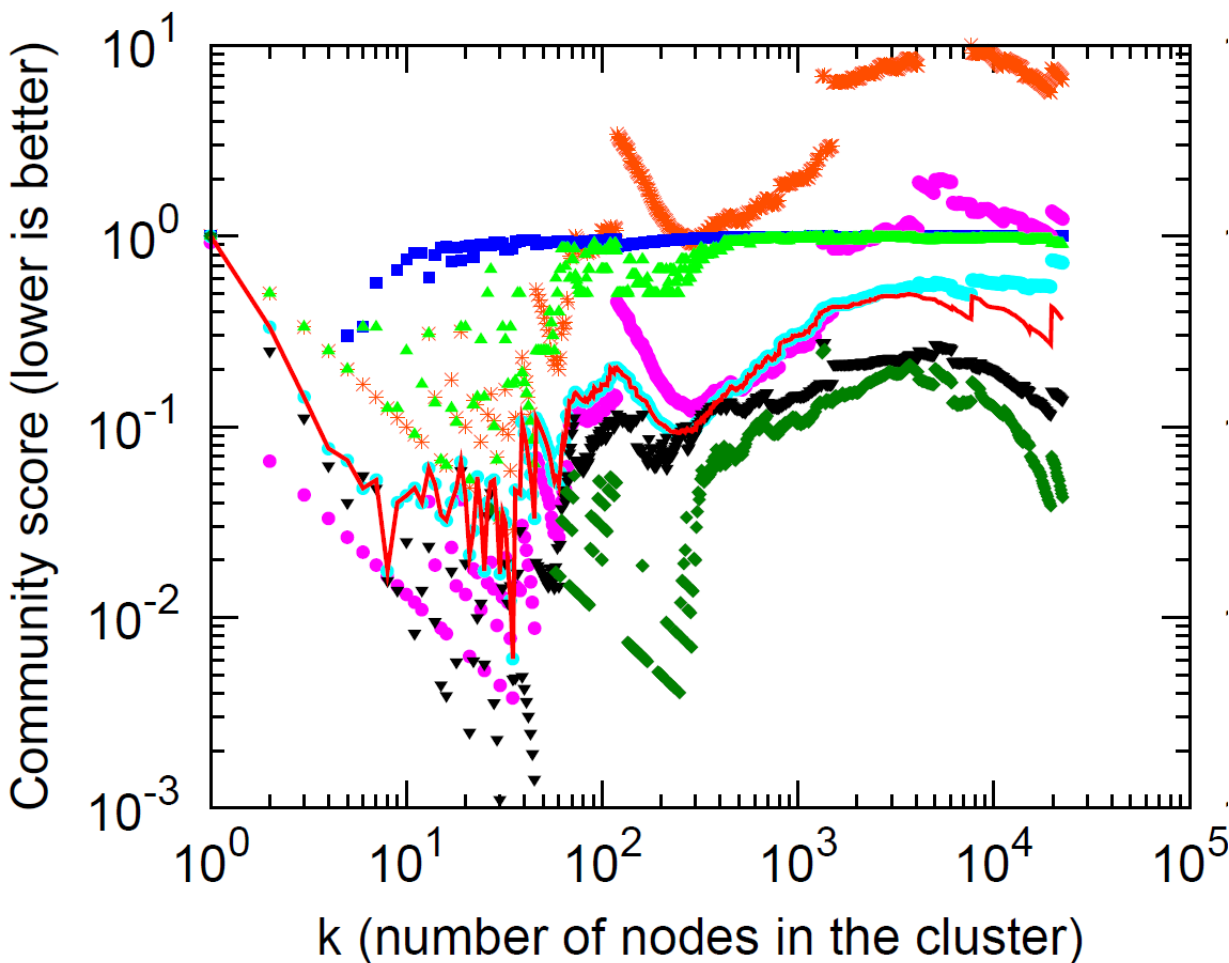
# Properties of Clusters (2)



- **Metis (red)** gives sets with better conductance
- **Spectral (blue)** gives tighter and more well-rounded sets



# Multi-criterion Objectives



■ All qualitatively similar

■ Observations:

- Conductance, Expansion, Norm-cut, Cut-ratio are similar
- Flake-ODF prefers larger clusters
- Density is bad
- Cut-ratio has high variance

Conductance —  
Expansion \*

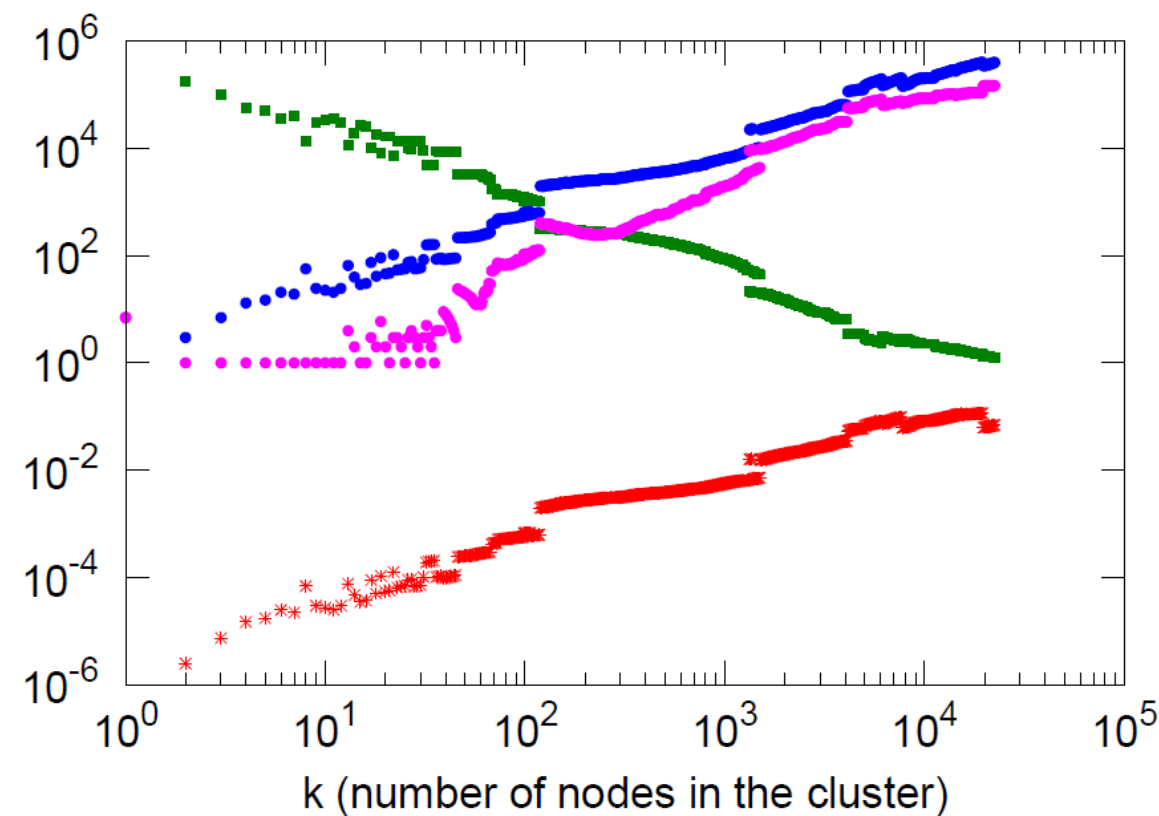
Internal Density ■  
Cut Ratio ●

Normalized Cut ●  
Maximum ODF ▲

Avg ODF ▼  
Flake ODF ◆



# Single-criterion Objectives



## Observations:

- All measures are monotonic
- **Modularity**
  - prefers large clusters
  - Ignores small clusters

Modularity \*

Modularity Ratio ■

Volume ●

Edges cut ●